

# Lecture 07

# Thread

CS397/497 – Wireless Protocols for IoT  
Branden Gena – Winter 2023

With some advice from Neal Jackson (UC Berkeley)

Materials in collaboration  
with Pat Pannuto (UCSD)

# Administrivia

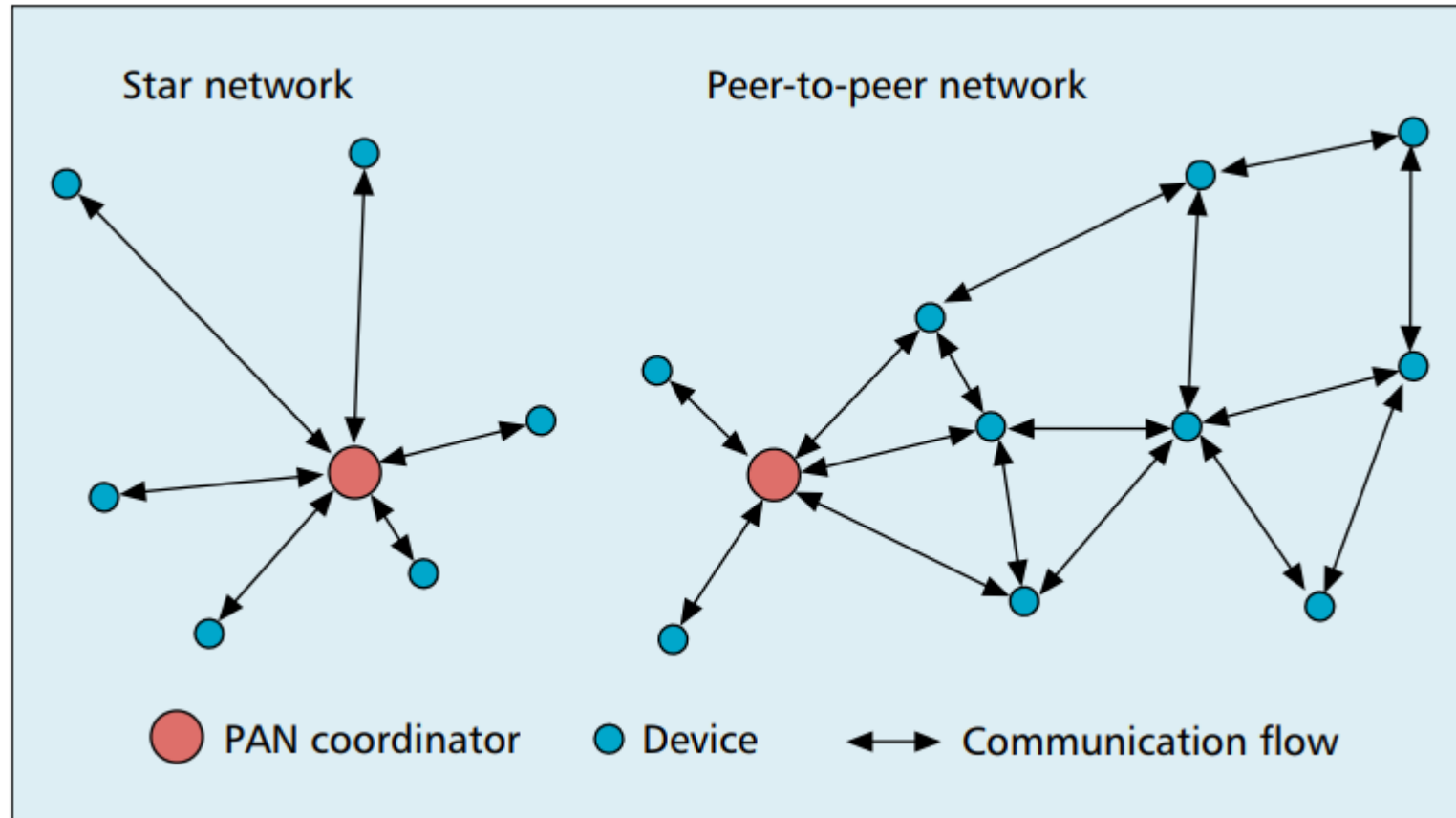
- Hw: BLE Packets due today (19/33)
  - Do it. It's good for you. Much like eating spinach.
- Lab: BLE due on Monday (6/33)
  - If you haven't started, do so ASAP
  - Some people had various hardware/software issues getting started
    - Once you've gotten to "blink an LED", everything should be good from there
- Lab: Thread (-1/33)
  - Should be out Friday or Saturday
  - Will require multiple of you to be together in one space to work on it

# Today's Goals

- Explore 802.15.4 packet structure
- Describe goals and capabilities of Thread networks
- Understand addressing in Thread networks
- Describe runtime behaviors like network joining

# 802.15.4 network topologies

- Only specifies PHY and MAC, but has use cases in mind



# Modes of operation

- Beacon-enabled PAN
  - Slotted CSMA/CA
  - Structured communication patterns
  - Optionally with some TDMA scheduled slots
  
- Non-beacon-enabled PAN
  - Unslotted CSMA/CA
  - No particular structure for communication
    - Could be defined by other specifications, like Thread or Zigbee

# Beacon-enabled superframe structure



- Beacons occur periodically [15 ms – 245 seconds]
  - Devices must listen to each beacon
- Contention Access Period
  - Slotted CSMA/CA synchronized by beacon start time
- Inactive Period
  - No communication occurring. Assumes sleepy devices

# Non-beacon-enabled PAN

Contention Access Period ...

- Same idea, just no beacons
  - Which removes synchronization benefit (and slotted CSMA/CA)
  - Also removes beacon listening cost
    - Devices only need to check for activity before transmitting
  - Still need an algorithm to determine when it should receive data
    - All the time is a huge energy drain
    - Algorithms can get complicated here
    - **Could BLE mechanism of listen-after-send apply?**
      - Only if sending to a high-power device, not among equals

# 802.15.4 specification versions

- 2011
  - Four PHY options (UWB)
  - MAC capability to support ranging (distance measurements)
- 2015
  - Six PHY options (RFID, Smart Utility, TV White Space)
  - Time Slotted Channel Hopping (TSCH) access control (TDMA+FDMA)
- 2020
  - Several new PHY options (China medical band, alternate modulations)



# Major 802.15.4 uses

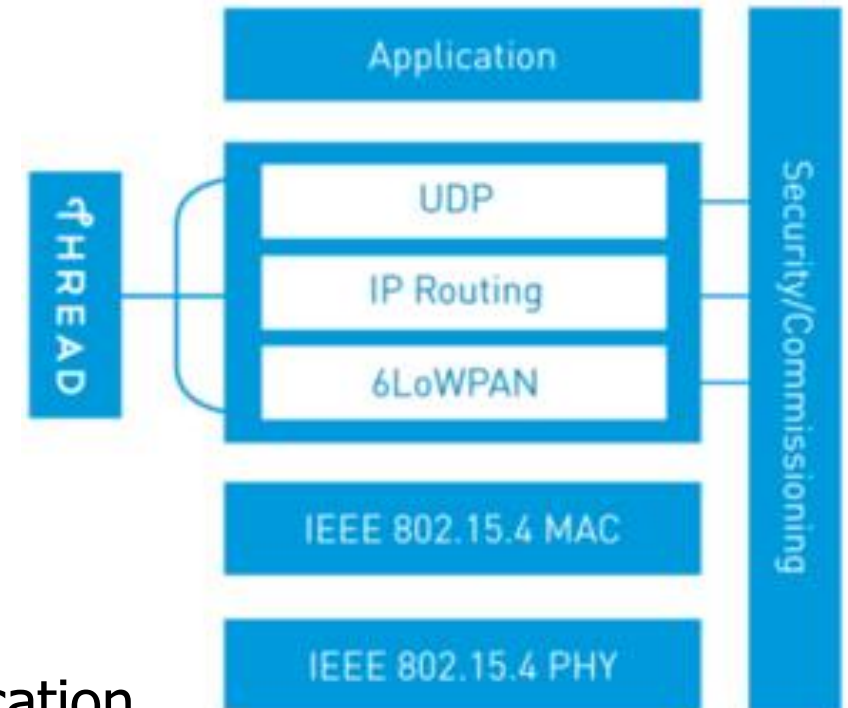
- Zigbee & Thread
- Both build upon the 15.4 PHY and Link primitives
  - Using most but discarding some
- Both consider higher-level application considerations
  - Building a network
  - Communicating between devices
  - Application logic

# Outline

- **Thread Overview**
- Thread Addressing
- Runtime Behavior
- Using IP

# Thread overview

- Build a networking layer on top of 15.4
  - Reuses most of PHY and MAC
  - Adds IP communication
  - Handles addressing and mesh maintenance
- Goals
  - Simplicity – easy to install and operate
  - Efficiency – years of operation on batteries
  - Scalability – hundreds of devices in a network
  - Security – authenticated and encrypted communication
  - Reliability – mesh networking without single point of failure
- Industry-focused, but based in academic research

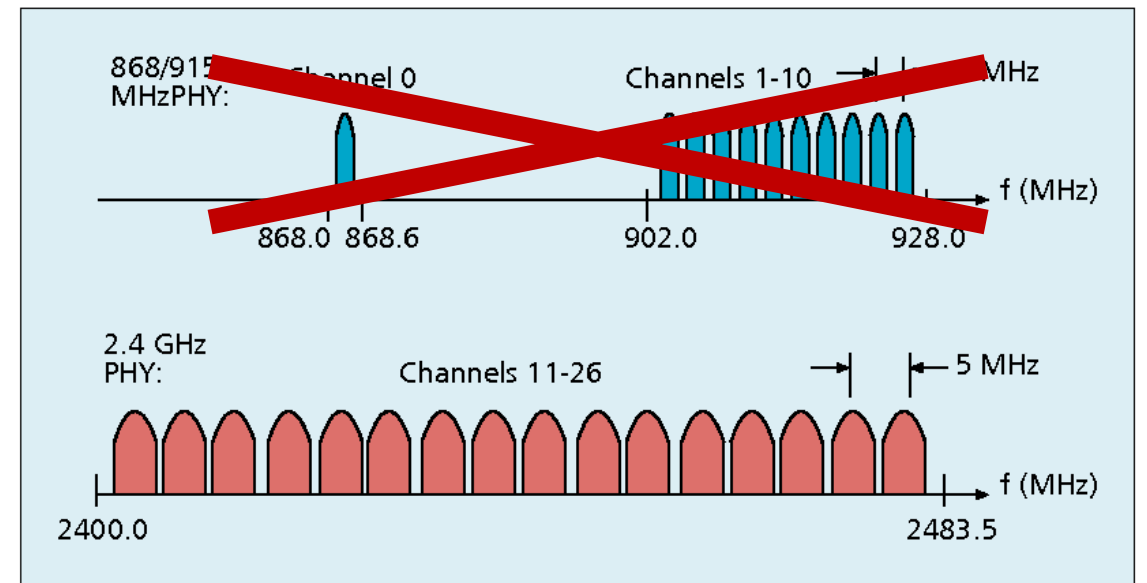


# References on Thread

- Request for specification: <https://www.threadgroup.org/ThreadSpec>
  - Frustratingly locked down 😞
- Overview on capabilities: <https://openthread.io/guides/thread-primer>
  - Excellent overview
  - Lifting heavily for these slides

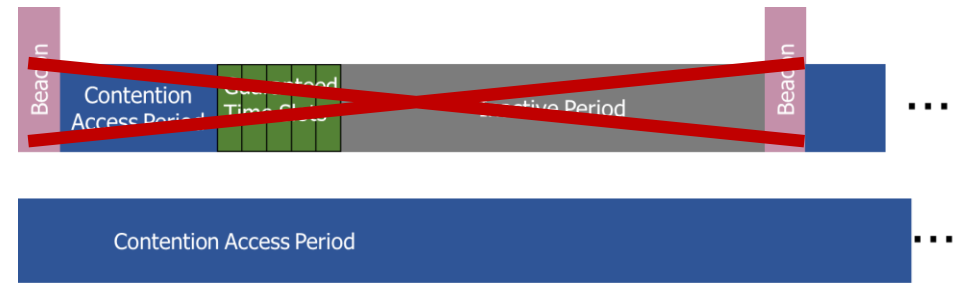
# Changes to Physical Layer

- Remove all non-2.4 GHz PHY options
- Otherwise the same
  - O-QPSK
  - 16 channels, 5 MHz spacing
  - Typical TX power 0 dBm
  - Typical RX sensitivity -100 dBm



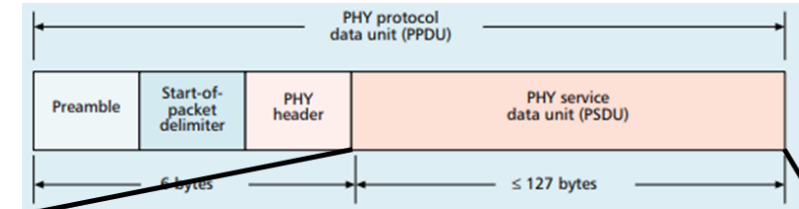
# Changes to Link Layer and MAC

- Non-beacon-enabled PAN only
  - No superframe structure
  - No periodic beacons
  - No Guaranteed Time Slots
- Throw out most existing MAC Commands
  - Remove network joining/leaving
  - Remove changing coordinators
  - Remove Guaranteed Time Slot request
  - Network joining will be handled at a higher layer



# Changes to Link Layer and MAC

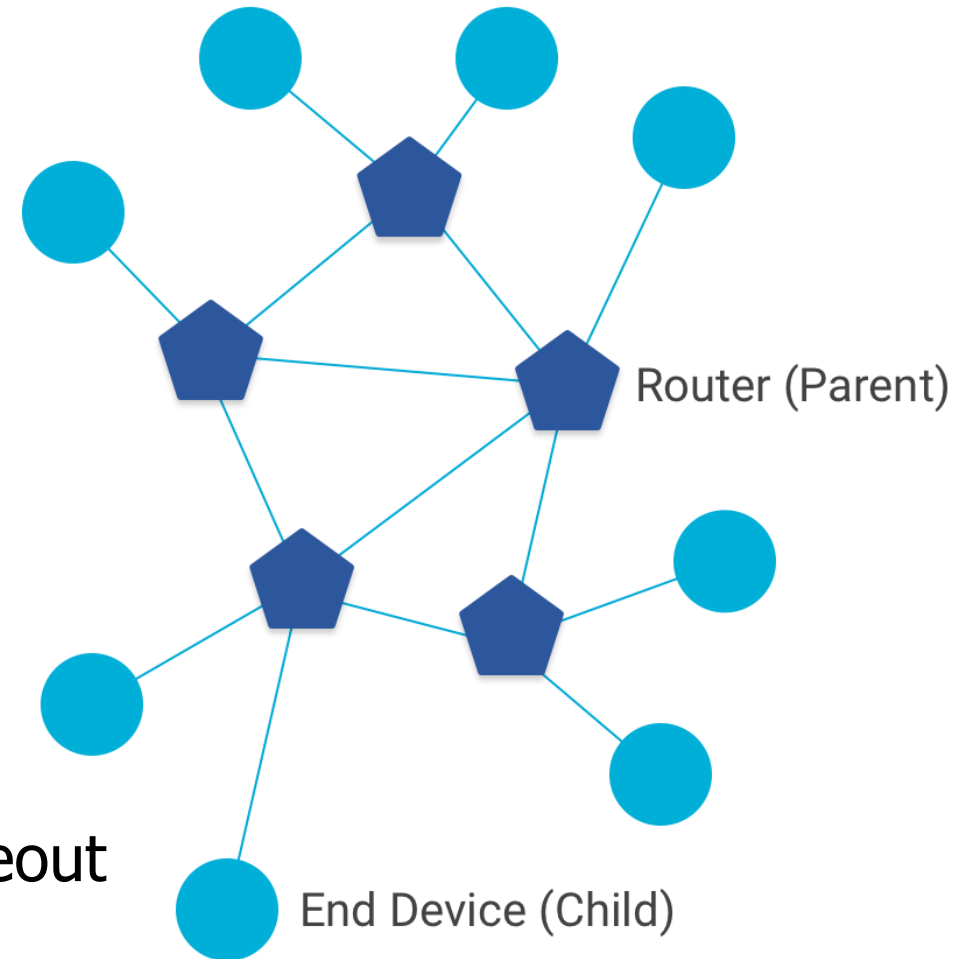
- Keep unslotted CSMA/CA algorithm
- Keep packet structure
- Keep Frame Types
  - Beacon
  - MAC Command
    - Beacon Request
    - Data Request
  - Data
  - Acknowledgement



Octets:2	1	0/2	0/2/8	0/2	0/2/8	variable	2
Frame control	Sequence number	Destination PAN identifier	Destination address	Source PAN identifier	Source address	Frame payload	Frame check sequence
Addressing fields							
<b>MAC header</b>						<b>MAC payload</b>	<b>MAC footer</b>

# Combination of star and mesh topology

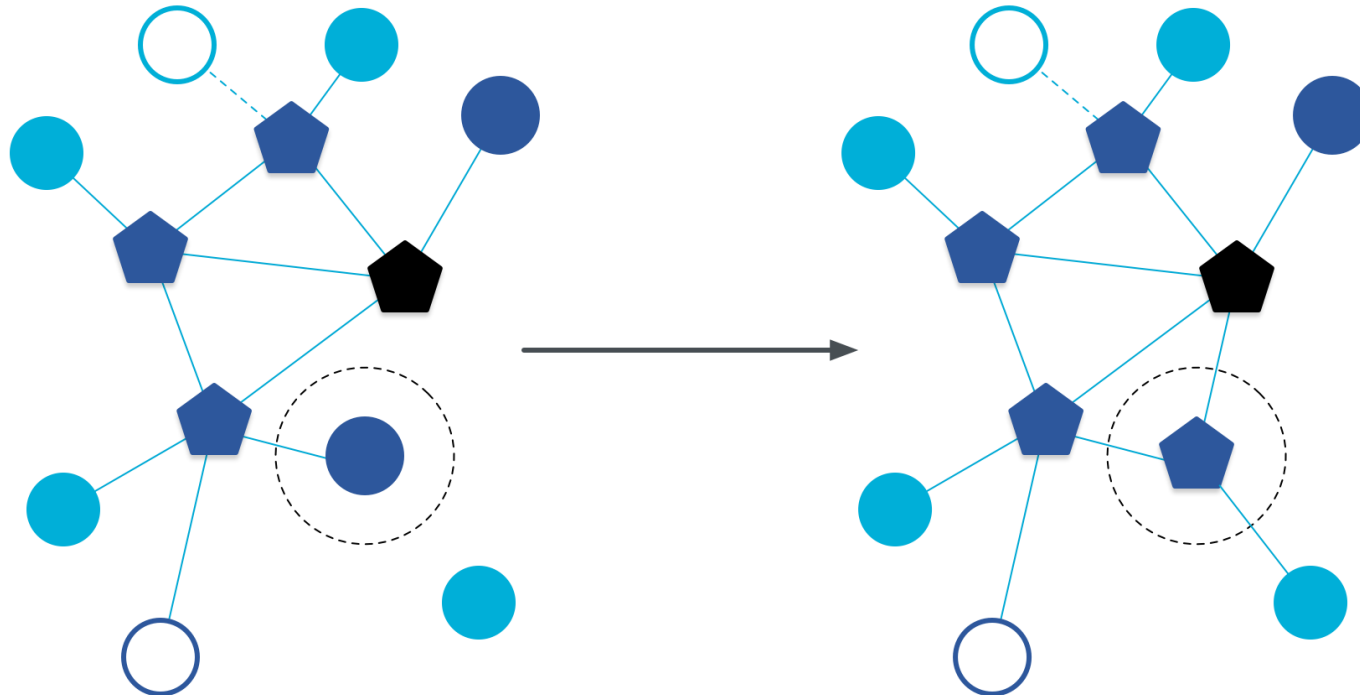
- Routers (parent)
  - Mesh communication with other routers
  - Radio always on
  - Forwards packets for network devices
  - Enables other devices to join network
  - 32 routers per network
- End devices (child)
  - Communicates with one parent (router)
  - Does not forward packets
  - Can disable transceiver to save power
    - Send packets periodically to avoid timeout
  - 511 end devices per router





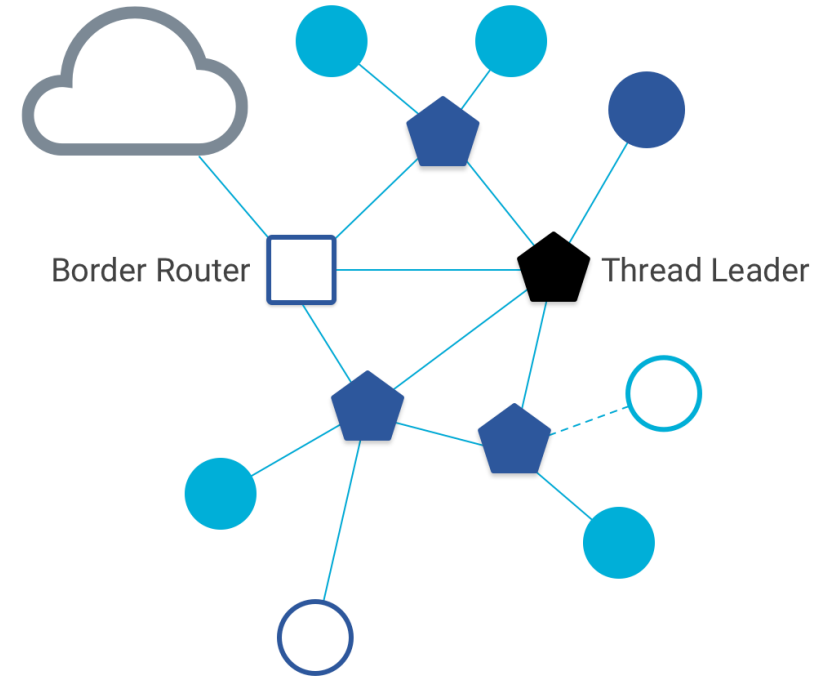
# Router promotion

- “Router Eligible End Device”
  - A router without any children
  - Can operate as an end device with one connection (lower power)
  - Promotes to a router when a joining end device relies on it
    - If there is room for an additional router (max 32, typical 16-23)



# Other special roles

- Thread leader
  - Device in charge of making decisions
    - Addresses, Joining details
  - Automatically selected from routers
    - One leader at any given time
    - Additional leader is selected if the network partitions
- Border router
  - Router that also has connectivity to another network
    - Commonly WiFi or Ethernet
  - Provides external connectivity
  - Multiple border routers may exist at once



# Why use Thread instead of basic 802.15.4?

- Full specification of upper layers
  - Clarifies how data is transmitted between devices on a network
  - Cleans up a lot of things otherwise left implementation-dependent
- Interaction with the world *outside* of the sensor network!
  - Gateway can be a dumb forwarder of packets
  - Devices can directly talk to NTP servers or POST data to a website!

# Outline

- Thread Overview
- **Thread Addressing**
- Runtime Behavior
- Using IP

# Thread uses IPv6 for communication

- Why IP?

- If Wireless Sensor Networks represent a future of billions of connected devices distributed throughout the physical world
- Why shouldn't they run standard protocols wherever possible?
- Why IPv6?
  - Generalized, Flexible, Capable

- Benefits

- Interoperability with normal computers and networks
- Reuse state of the art developed standards instead of remaking them
  - Security, Naming, Discovery, Services

- Costs

- Packet overhead can be high (will fix)
- Complexity for supporting protocols

Hui and Culler, "[IP is Dead, Long Live IP for Wireless Sensor Networks](#)". 2008

# Background: IPv6

- Replacement to Internet Protocol v4
  - (Something unrelated used version number 5)
- Extended addressing for devices
  - 32-bits for IPv4 addresses -> 128-bits for IPv6 addresses
  - Example: a39b:239e:ffff:29a2:0021:20f1:aaa2:2112
- Supports multiple transmit models
  - Broadcast: one-to-all
  - Multicast: one-to-many
  - Unicast: one-to-one
- Various other improvements

# Background: IPv6 address notation rules

- Groups of zeros can be replaced with “::”
  - Can only use “::” in one place in the address
- Leading zeros in a 16-bit group can be omitted

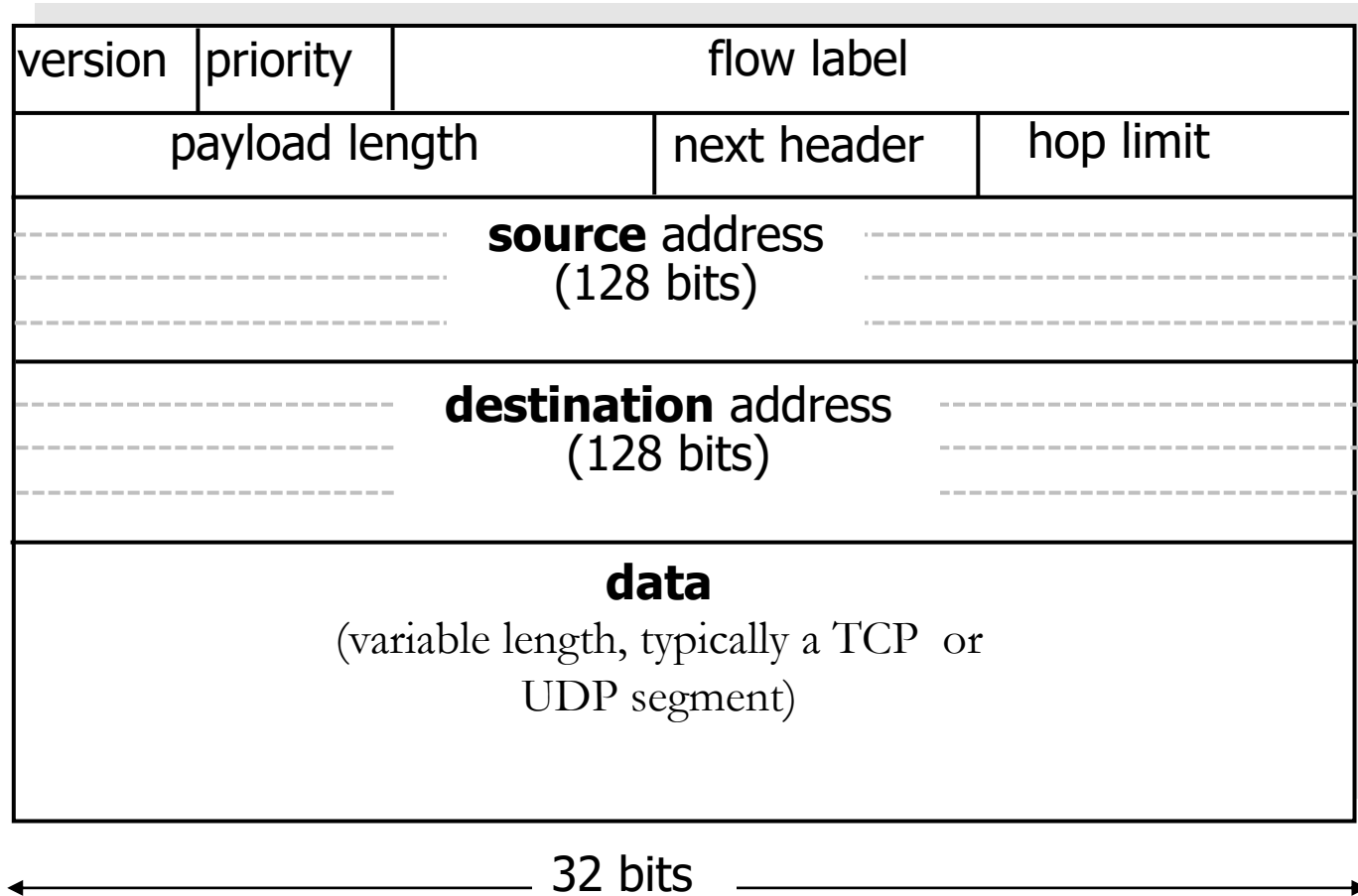
0000:0000:0000:0000:0000:0000:0000:0001 → ::1

2345:1001:0023:1003:0000:0000:0000:0000 → 2345:1001:23:1003::

aecb:0222:0000:0000:0000:0000:0000:0010 → aecb:222::10

- Special addresses
  - Localhost - ::1 (IPv4 version is 127.0.0.1)
  - Link-Local Network - fe80:: (bottom 64-bits are ~device MAC address)
  - Local Network – fc00:: and fd00::
  - Global Addresses – 2000:: (various methods for allocating bottom bits)

# Background: IPv6 datagram format



- **Priority**: like “type of service” in IPv4.
- **Flow label**: ambiguous
- **Next header**: TCP, UDP
- **Hop limit** = TTL

*how much overhead?*

- **40 bytes** of IPv6
- 20 more than IPv4



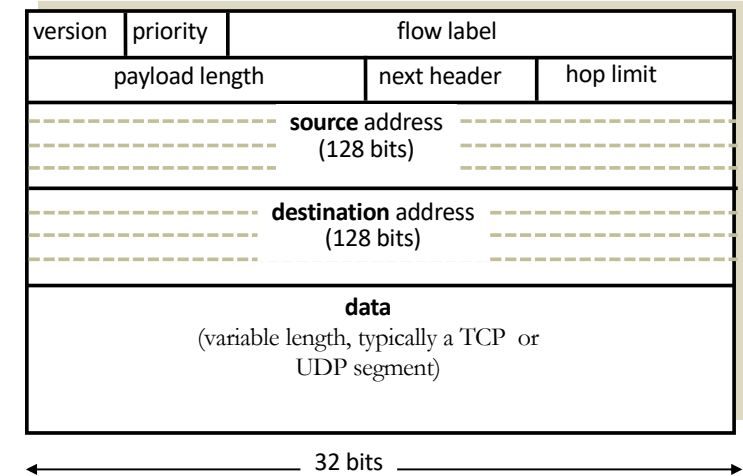
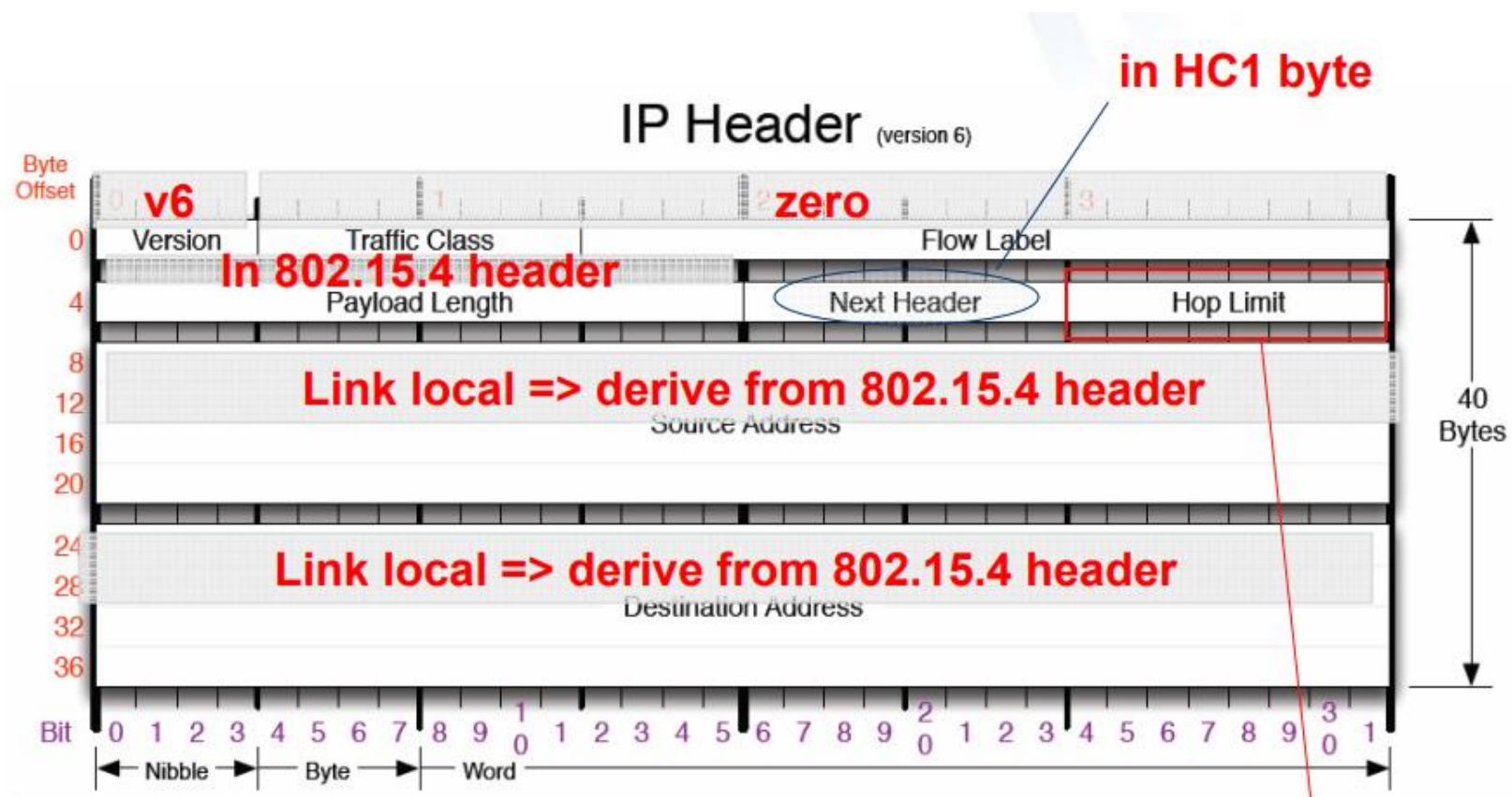
# 6LoWPAN

- Method for running IPv6 over 802.15.4 links
  - IPv6 over **L**ow-Power **W**ireless **P**ersonal **A**rea **N**etworks
  - IETF Standard ([RFC4944](#) + updates in [RFC6282](#))
- Directly out of the research world (Jonathan Hui + David Culler)
  - Research Paper: [IP is Dead, Long Live IP for Wireless Sensor Networks](#)
  - Thesis of work: sensor networks can and should use IPv6
- Important goals
  - Compress IPv6 headers
  - Handle fragmentation of packets
  - Enable sending packets through mesh

# 6LoWPAN header compression

- 40 bytes of IPv6 header are a lot for a 127-byte payload
- Most important goals
  - Communication with devices in the 15.4 network should be low-overhead
  - Communication outside of the 15.4 network should still minimize overhead where possible
- Assume a bunch of common parameters to save space
  - A bunch of options are set to default values
  - Payload length can be re-determined from packet length
  - Source/Destination addresses can often be reassembled from link layer data
    - Plus information about network address assignment known by routers
- Border router “inflates” the packet before sending externally

# Example of compression



• [http://www.visi.com/~mjb/Drawings/IP\\_Header\\_v6.pdf](http://www.visi.com/~mjb/Drawings/IP_Header_v6.pdf)

## • Example of HC1 header compression

- Note: Thread actually uses IPHC from rfc6282 (not HC1), but similar idea

# 6LoWPAN fragmentation

- Only the first packet of the fragments will hold the IPv6 header
  - Tag, offset, and size are used to reconstruct

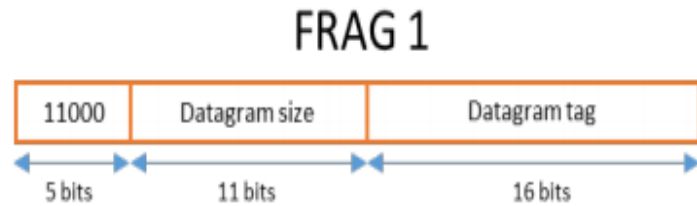


Figure 15. First Fragment Header

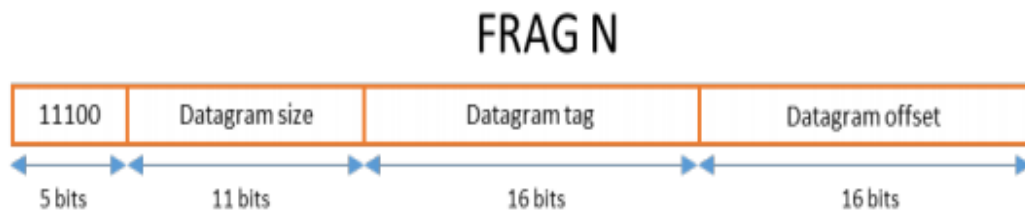


Figure 16. Subsequent Fragment Header

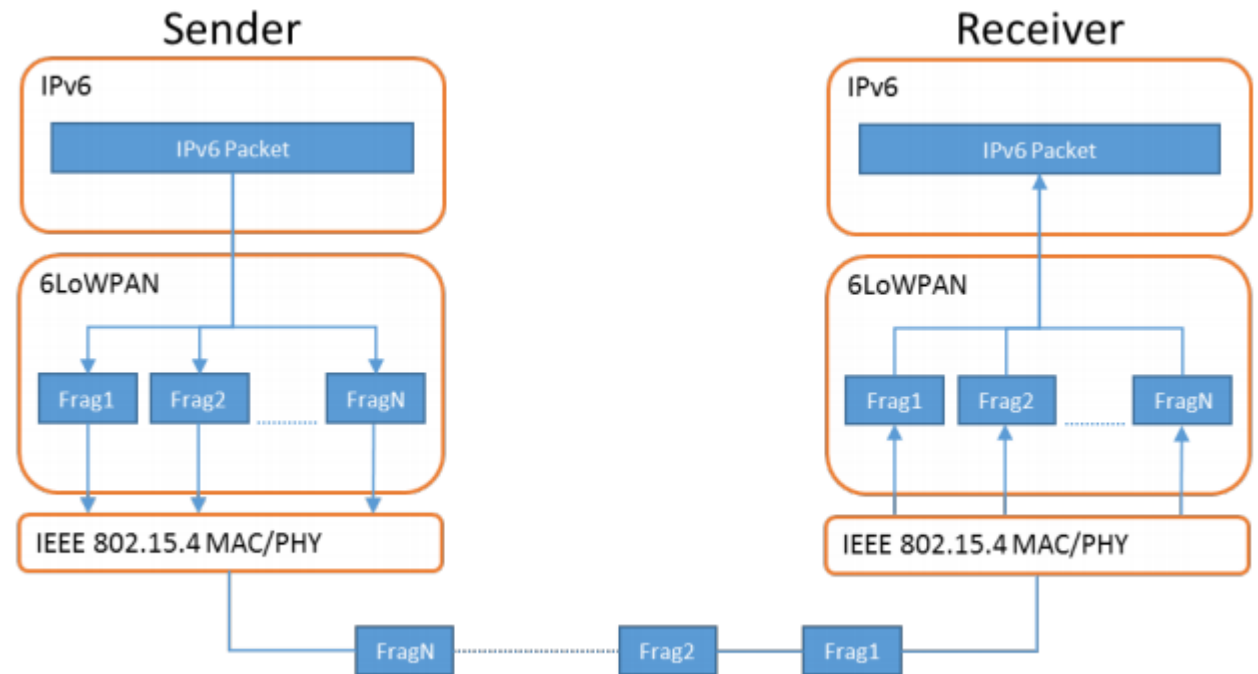
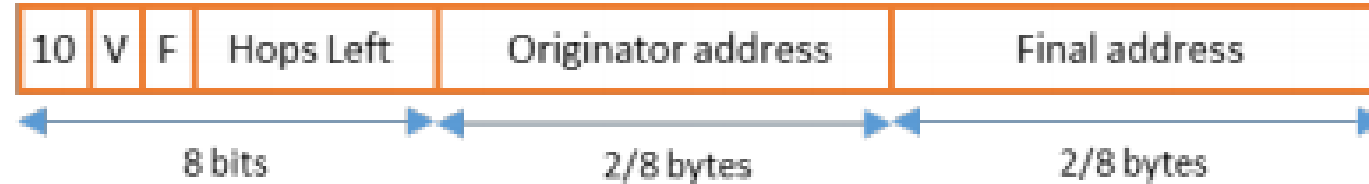


Figure 14. Fragmenting and Reassembling an IPv6 Packet

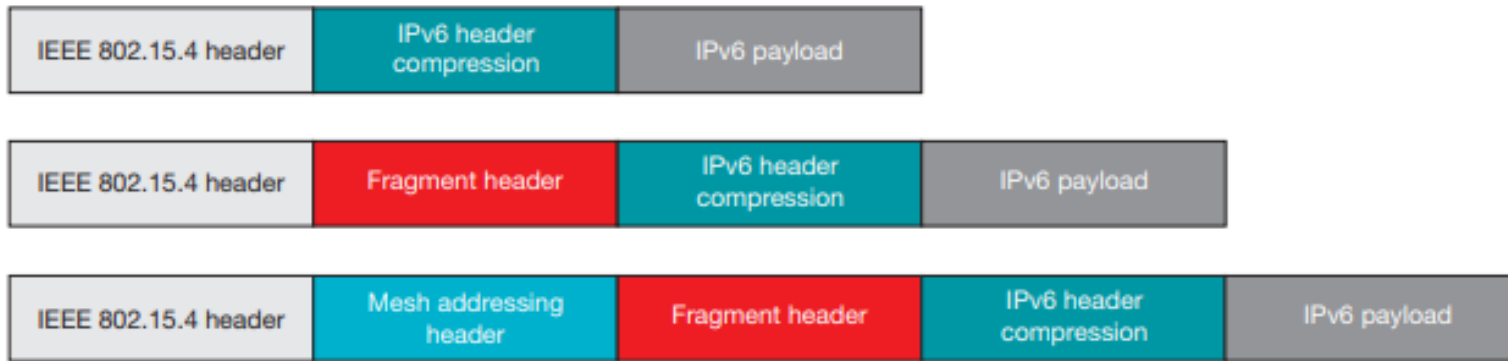
# 6LoWPAN mesh forwarding

- Additional header with originator and final addresses



**Figure 17. Mesh Header Format**

- Which of these headers are used depends on the packet



*Figure 4. 6LoWPAN stacked headers*

# Sidebar: IPv6 over BLE

- [RFC7668](#) defines 6LoWPAN techniques for BLE connections

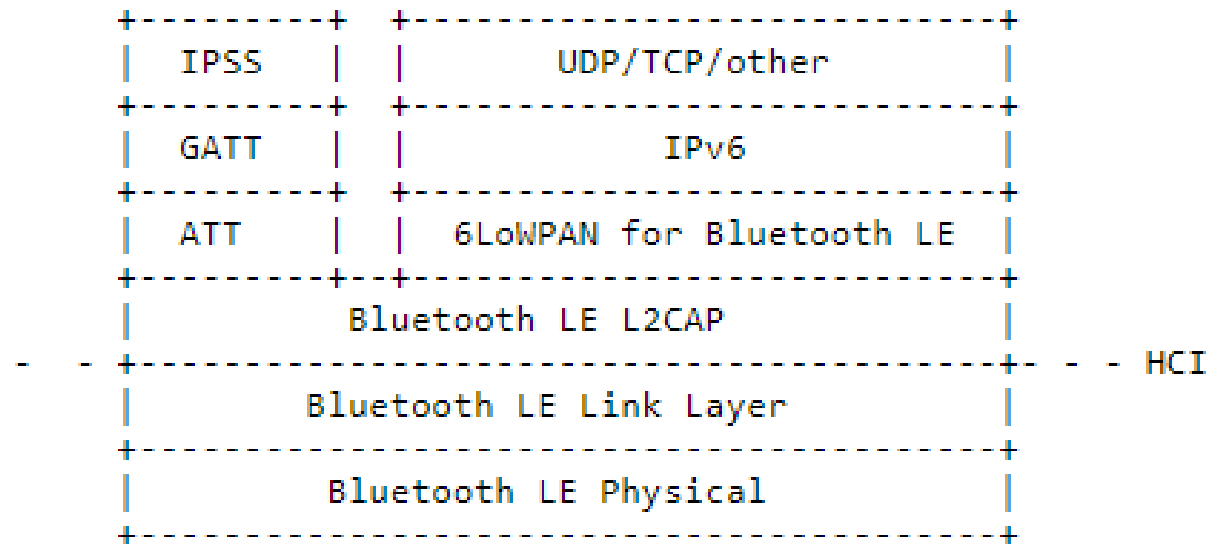
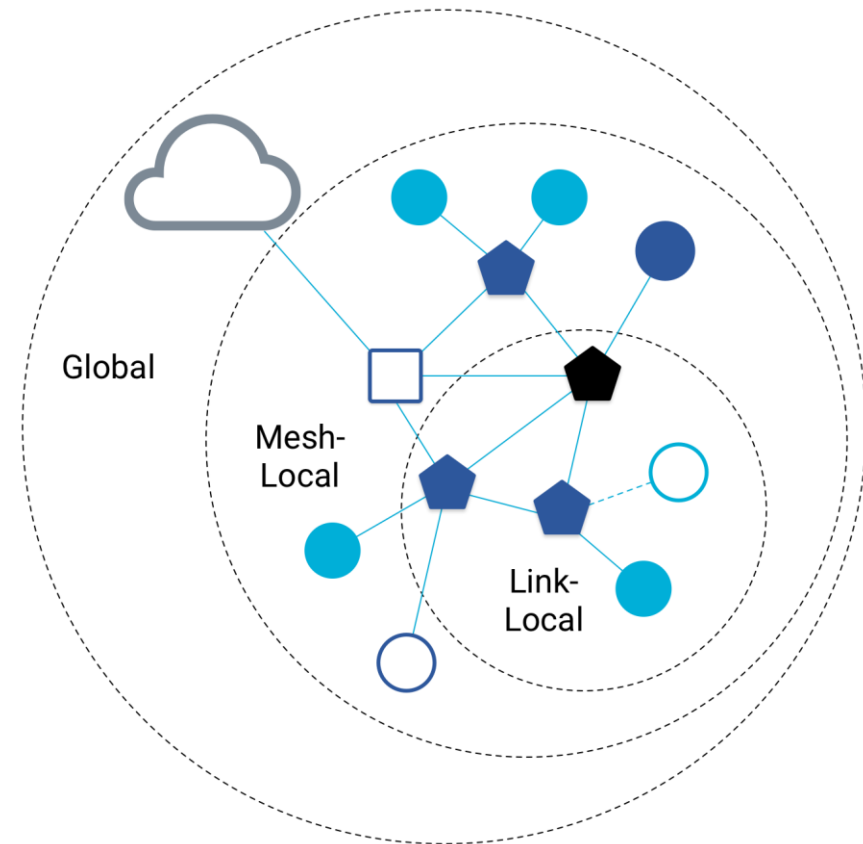


Figure 3: IPv6 and IPSS on the Bluetooth LE Stack

# Benefit to IPv6: multiple address spaces per Thread device

- Each device gets an IPv6 address for each way to contact it
  - Global IP address
  - Mesh-local IP address
  - Link-local IP address
- Topology-based IP address
  - Send to parent
  - Send to child
- Role-based IP address(es)
  - Send to all Routers
  - Send to Border Router



# Traditional addresses in Thread

- Link-Local Addresses
  - FE80::/16
  - Bottommost 64-bits are EUI-64 (MAC address with 0xFFFE in the middle)
  - Permanent for a given device (no matter the network)
  - Used for low-layer interactions with neighbors (discovery, routing info)
- Mesh-Local Addresses
  - FD00::/8 (FD00:: and FC00:: are for local networks in IPv6)
  - Remaining bits are randomly chosen as part of joining the network
  - Permanent while connection is maintained to a network
  - Used for application-layer interactions
- Global Addresses
  - 2000::/3 (2000:: are for global, unicast IP addresses in IPv6)
  - Public address for communicating with broader internet through Border Router
  - Various methods for allocation (SLAAC, DHCP, Manual)

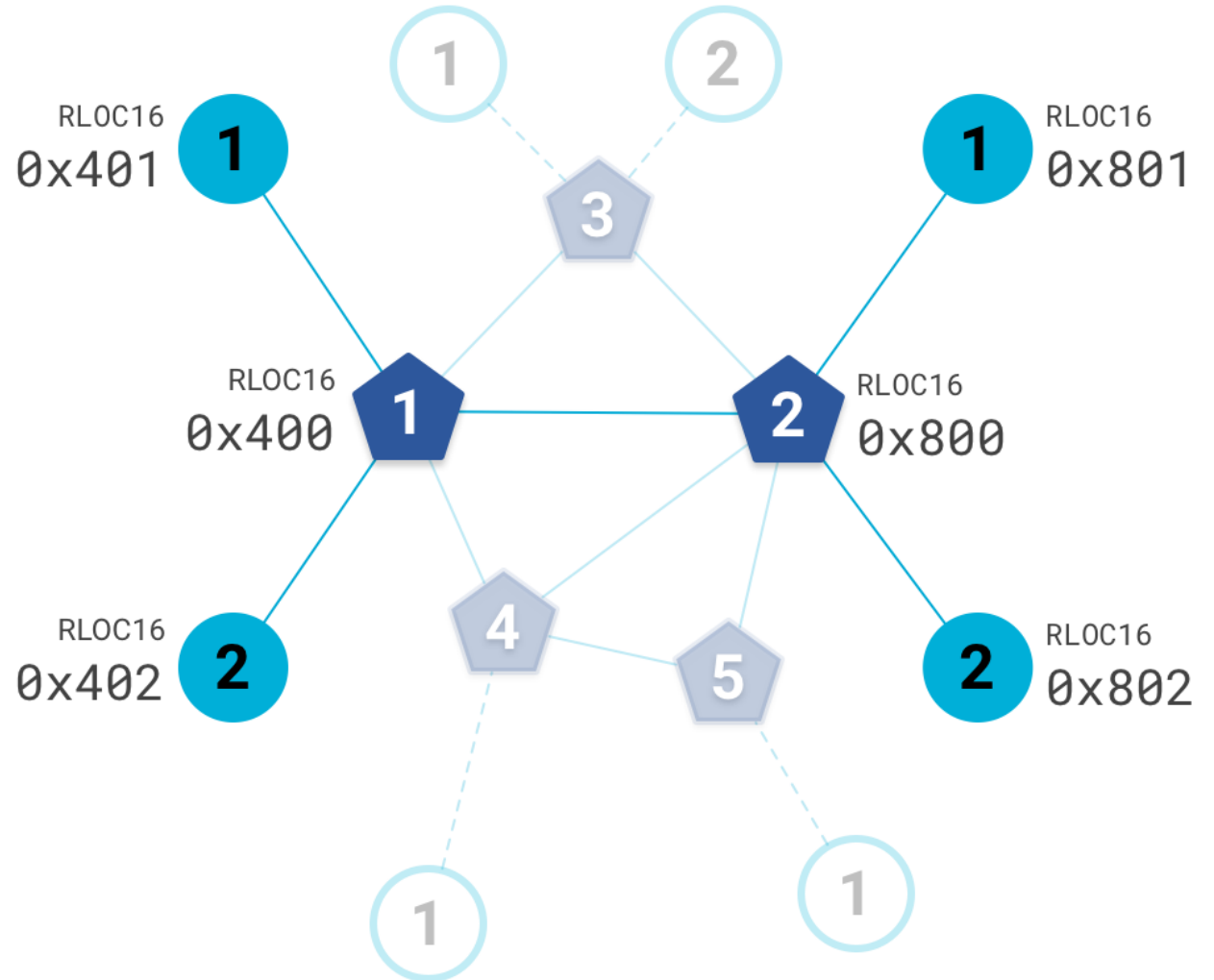


# Topology-based addresses in Thread

- FD00::00FF:FE00:RLOC16
  - Same top bits as mesh-local
- Routing Locator (RLOC)
  - Router ID concatenated with Child ID



- Changes with network topology
  - Used for routing packets



# Role-based addresses in Thread

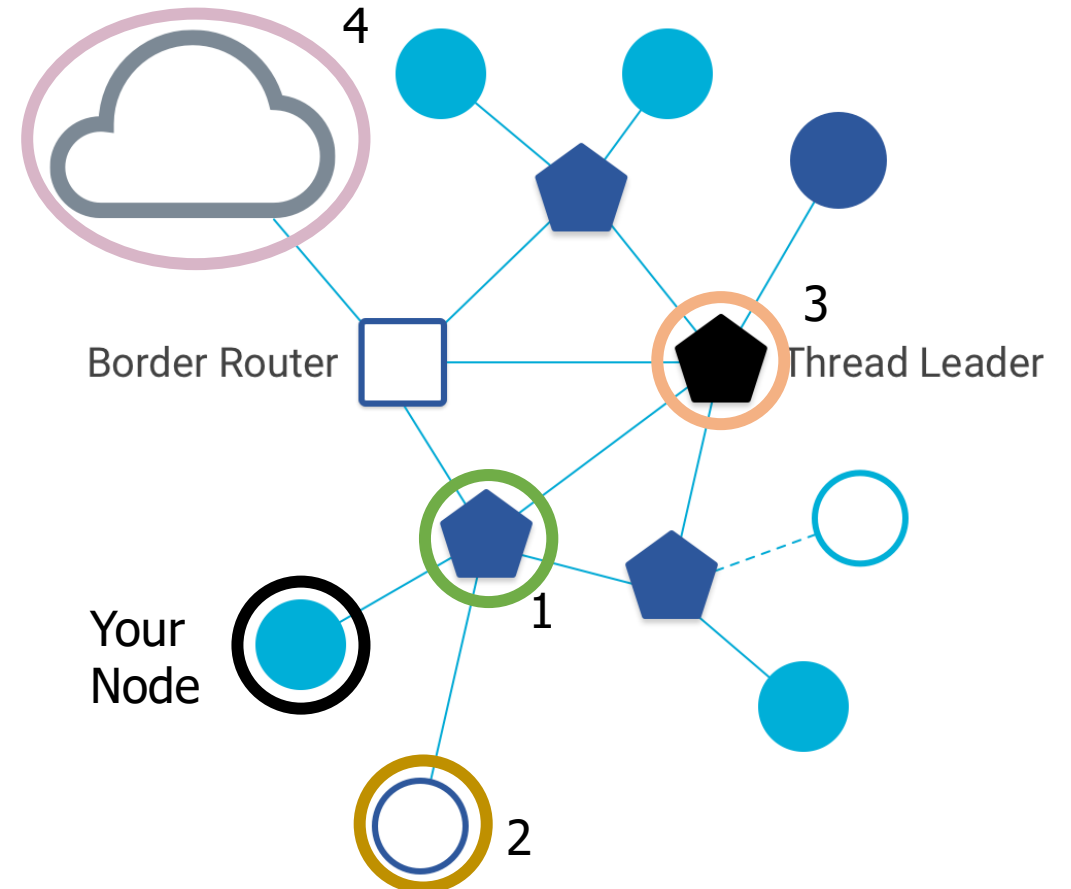
- Multicast
  - FF02::1 – link-local, all listening devices
  - FF02::2 – link-local, all routers/router-eligible
  - FF03::1 – mesh-local, all listening devices
  - FF03::2 – mesh-local, all routers/router-eligible
  
- Anycast
  - FD00::00FF:FE00:FC**xx**
    - 00 – Thread Leader
    - 01-0F – DHCPv6 Agent
    - 30-37 – Commissioner
    - etc.

# Break + Question

Which type of address to use for communication

1. Global
2. Mesh-local
3. Link-local
4. Topology
5. Role-based

- Communicate with each circled target



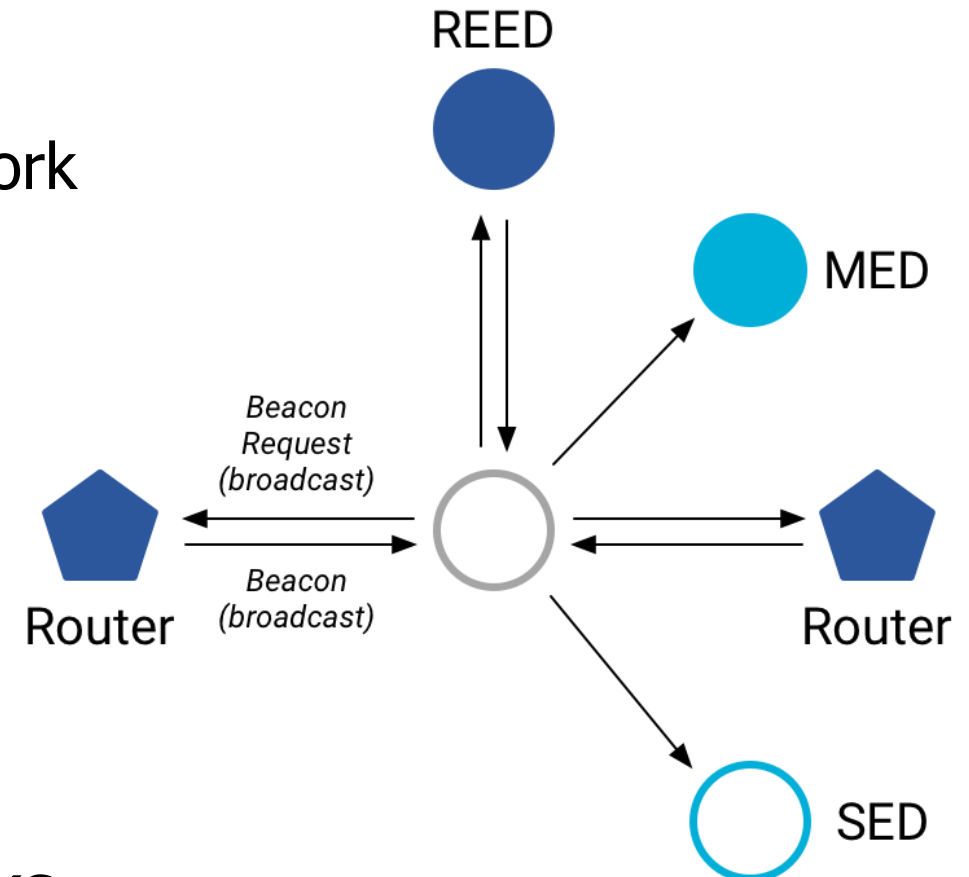


# Outline

- Thread Overview
- Thread Addressing
- **Runtime Behavior**
- Using IP

# Discovering Thread networks

- “Beacon Request” MAC command
  - Routers/Router-eligible devices respond
  - Payload contains information about network
- Thread network specification
  - PAN ID – 16-bit ID
  - XPAN ID – extended 64-bit ID
  - Network Name – human-readable
- Active scanning across channels can quickly find all existing nearby networks



# Creating a new network

- Select a channel (possibly by scanning for availability)
- Become a router
  - Elect yourself as Thread Leader
  - Respond to Beacon Requests from other devices
- Further organization occurs through Mesh-Level Establishment protocol

# Mesh-Level Establishment

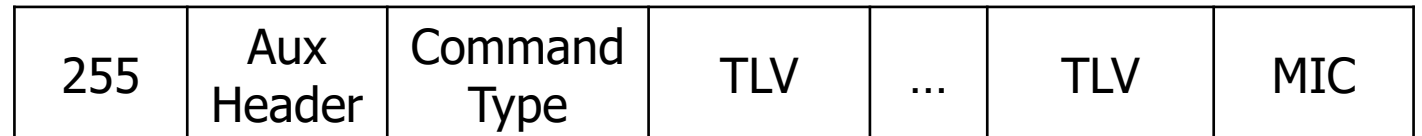
- Creating and configuring mesh links
  - Payloads placed in UDP packets within IPv6 payloads

- Commands for mesh

- Establish link
- Advertise link quality
- Connect to parent



OR (secure version)



- TLVs (Type-Length-Value)

- Various data types that may be helpful within those packets
- Addresses, Link Quality, Routing Data, Timestamps



# Joining an existing network

- All devices join as a child of some existing router

## 1. Send a Parent Request (to all routers/router-eligible)

- Using the multicast, link-local address

## 2. Receive a Parent Response (from all routers/router-eligible separately)

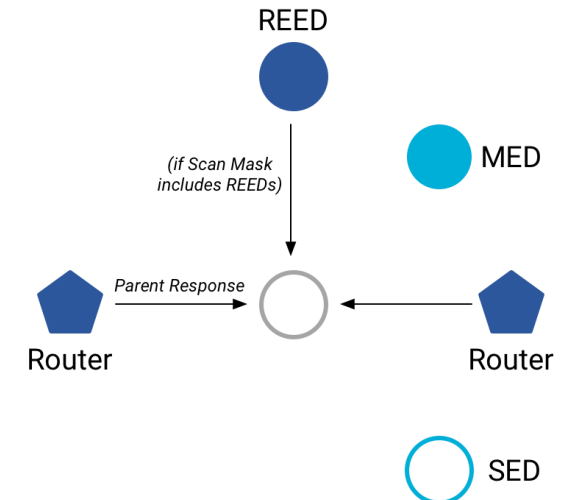
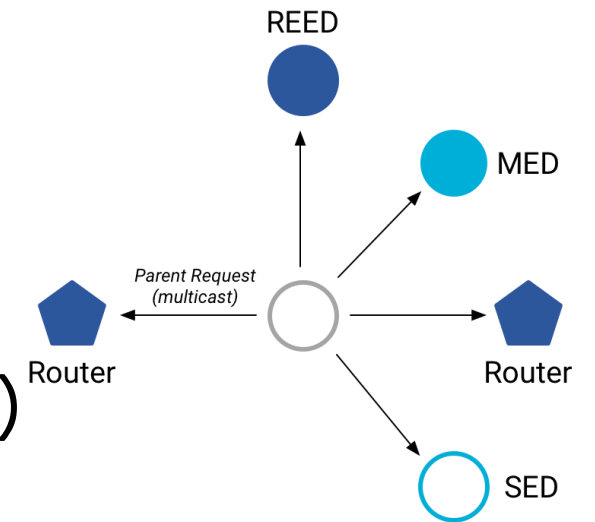
- Contains information on link quality

## 3. Send a Child ID Request (to router with best link)

- Contains parameters about the new child device

## 4. Receive a Child ID Response (from that router)

- Contains address configurations

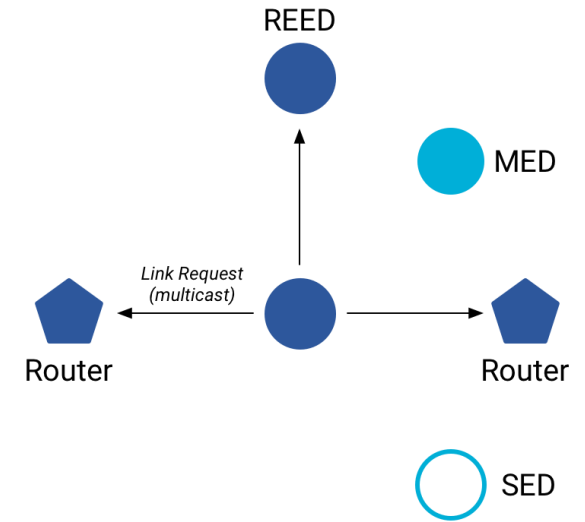


# Becoming a router

- Thread tries to maintain 16-23 routers (max 32)
  - Goals: path diversity, extend connectivity

## 1. Send a Link Request (to all routers/router-eligible)

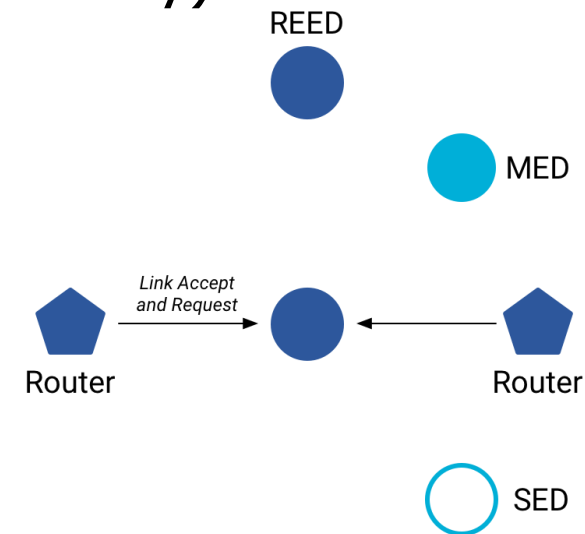
- Using the multicast, link-local address



## 2. Receive Link Accept and Request (from each router separately)

- Forms bi-directional link

## 3. Send a Link Accept (to each router individually)



# Outline

- Thread Overview
- Thread Addressing
- Runtime Behavior
- **Using IP**

# Communicating with IP

- Any communication that layers on top of IP is now possible
  - If there is a library to support it
- Common choices
  - UDP
    - DNS – translate hostnames into IP addresses
    - SNTP – get real-world time, accuracy better than 1 second
    - CoAP – send and receive data

# Constrained Application Protocol - CoAP

- HTTP, but over UDP targeting less-capable devices
  - Same REST architecture
  - Adds capability for automatic retransmissions



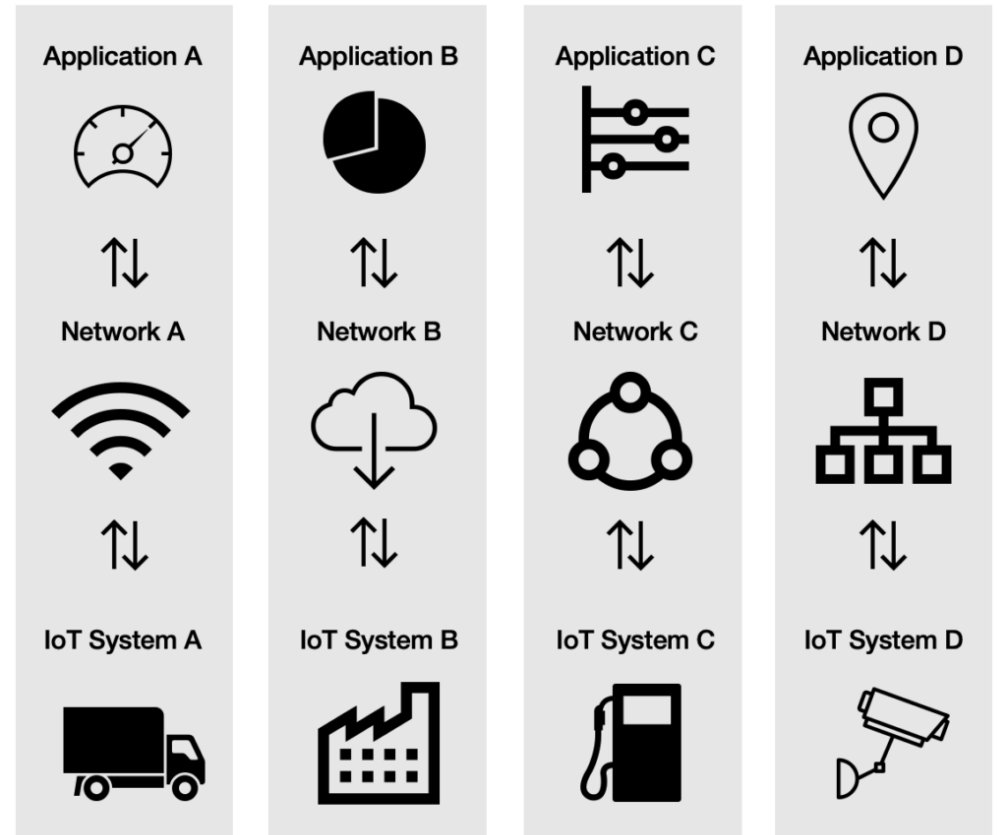
- CoAP Requests
  - Have a type: GET, POST, PUT, DELETE
  - Have a URL: /file/etc
  - Have data up to 65 KB

# Sensor networks don't use TCP (yet?)

- Uncommon choice: TCP
  - Concerns: Too large, too slow, poorly suited to lossy networks
  - Also concerning: We're just replicating TCP poorly
- Work in progress:
  - Sam Kumar, Michael Anderson, Hyung-Sin Kim, David Culler. ["Performant TCP for Low-Power Wireless Networks"](#). 2020.
  - The debate is still very much open

# A problem: the siloed internet of things

- Problem: companies are more interested in selling you the whole stack
  - Which then makes it harder for devices to be interoperable
- This is not Thread or IP-specific, but a problem all IoT devices are facing
- Branden's concern:
  - What IP address do you send data to?
  - Manufacturer's server is an obvious choice



# Outline

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- Using IP