# **Lecture 07 Thread**

# CS397/497 – Wireless Protocols for IoT Branden Ghena – Spring 2024

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Northwestern

#### Administrivia

- Hw: BLE Packets due today
	- Do it. It's good for you. Much like eating spinach.

- Lab: BLE due on Tuesday
	- If you haven't started, do so ASAP
	- Some people had various hardware/software issues getting started
		- Once you've gotten to "blink an LED", everything should be good from there

#### Today's Goals

• Explore 802.15.4 packet structure

• Describe goals and capabilities of Thread networks

• Understand addressing in Thread networks

• Describe runtime behaviors like network joining

# 802.15.4 network topologies

• Only specifies PHY and MAC, but has use cases in mind



#### Modes of operation

- Beacon-enabled PAN
	- Slotted CSMA/CA
	- Structured communication patterns
	- Optionally with some TDMA scheduled slots

- Non-beacon-enabled PAN
	- Unslotted CSMA/CA
	- No particular structure for communication
		- Could be defined by other specifications, like Thread or Zigbee

#### Beacon-enabled superframe structure



- Beacons occur periodically [15 ms 245 seconds]
	- Devices must listen to each beacon
- Contention Access Period
	- Slotted CSMA/CA synchronized by beacon start time
- Inactive Period
	- No communication occurring. Assumes sleepy devices

# Non-beacon-enabled PAN

Contention Access Period …

- Same idea, just no beacons
	- Which removes synchronization benefit (and slotted CSMA/CA)
	- Also removes beacon listening cost
		- Devices only need to check for activity before transmitting
	- Still need an algorithm to determine when it should receive data
		- All the time is a huge energy drain
		- Algorithms can get complicated here
		- **Could BLE mechanism of listen-after-send apply?**
			- Only if sending to a high-power device, not among equals

# 802.15.4 specification versions

- 2011
	- Four PHY options (UWB)
	- MAC capability to support ranging (distance measurements)
- 2015
	- Six PHY options (RFID, Smart Utility, TV White Space)
	- Time Slotted Channel Hopping (TSCH) access control (TDMA+FDMA)
- 2020
	- Several new PHY options (China medical band, alternate modulations)

#### Major 802.15.4 uses

- Zigbee & Thread
- Both build upon the 15.4 PHY and Link primitives
	- Using most but discarding some
- Both consider higher-level application considerations
	- Building a network
	- Communicating between devices
	- Application logic

# **Outline**

- **Thread Overview**
- Thread Addressing
- Runtime Behavior
- Using IP

What's the need for Thread?

- Born from Google's Nest team
- Requirements:
	- IP-based: interoperate with rest of the Internet
	- Scalable: hundreds of devices in a network
	- Low power: years of operation on batteries
	- Secure: authenticated and encrypted communication
	- Reliability mesh networking without single point of failure
- 2014: Thread Group alliance formed
	- Goal: develop, maintain and drive adoption of Thread as an industry networking standard for IoT

Source: <u>openthread.io</u>, [threadgroup.org](http://www.threadgroup.org/) and the set of t



Google

**SIEMENS** 

amazon

888 SmartThings

Qualcomm

#### Thread overview

- Build a networking layer on top of 15.4
	- Reuses most of PHY and MAC
	- Adds IP communication
	- Handles addressing and mesh maintenance

• Industry-focused, but based in academic research



#### References on Thread

- Request for specification: <https://www.threadgroup.org/ThreadSpec>
	- Frustratingly locked down

- Overview on capabilities: <https://openthread.io/guides/thread-primer>
	- Excellent overview
	- Lifting heavily for these slides
- Good overview on Thread Networking:
	- <https://www.threadgroup.org/support#Whitepapers>
	- See the "Thread Network Fundamentals" whitepaper

# Changes to Physical Layer

• Remove all non-2.4 GHz PHY options

- Otherwise the same
	- O-QPSK
	- 16 channels, 5 MHz spacing
	- Typical TX power 0 dBm
	- Typical RX sensitivity -100 dBm



# Changes to Link Layer and MAC

- Non-beacon-enabled PAN only
	- No superframe structure
	- No periodic beacons
	- No Guaranteed Time Slots



- Throw out most existing MAC Commands
	- Network joining will be handled at a higher layer
		- Remove network joining/leaving
		- Remove changing coordinators
		- Remove Guaranteed Time Slot request

# Changes to Link Layer and MAC

- Keep unslotted CSMA/CA algorithm
- Keep packet structure
- Keep Frame Types
	- Beacon
	- MAC Command
		- Beacon Request
		- Data Request
	- Data
	- Acknowledgement



# Thread networks use a mix of star and mesh topologies





**Figure 1: Star Topology** 

Central PAN coordinator in charge of network

• Single point of failure End Devices can be less complicated

PAN Coordinator still in charge But mid-level Routers add route redundancy And End Devices didn't have to change at all

# Combination of star and mesh topology

- Routers (parent)
	- Mesh communication with other routers
	- Radio always on
	- Forwards packets for network devices
	- Enables other devices to join network
	- Up to 32 routers per network
- End devices (child)
	- Communicates with one parent (router)
	- Does not forward packets
	- Can disable transceiver to save power
		- Send packets periodically to avoid timeout
	- Up to 511 end devices per router ( $\sim$ 16k total)



# Other special roles

- Thread leader
	- Device in charge of making decisions
		- Addresses, Joining details
	- Automatically selected from routers
		- One leader at any given time
		- Additional leader is selected if the network partitions
- Border router
	- Router that also has connectivity to another network
		- Commonly WiFi or Ethernet
	- Provides external connectivity
	- Multiple border routers may exist at once



#### Why use Thread instead of basic 802.15.4?

- Full specification of upper layers
	- Clarifies how data is transmitted between devices on a network
		- Not just one single hop away
	- Cleans up a lot of things otherwise left implementation-dependent

- Interaction with the world *outside* of the sensor network!
	- Gateway can be a dumb forwarder of packets
	- Devices can directly talk to NTP servers or POST data to a website!

#### Break + Question

- What kinds of IoT devices are likely to have each thread role?
- End Device

• Router

• Border Router



#### Break + Question

- What kinds of IoT devices are likely to have each thread role?
- End Device
	- Battery-powered sensor
	- Nest Temperature Sensor
- Router
	- Wall-powered device
	- Nest Thermostat
- Border Router
	- Must have Internet connectivity
	- Custom gateway of some type?



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# Thread uses IPv6 for communication

- Why IP?
	- If Wireless Sensor Networks represent a future of billions of connected devices distributed throughout the physical world
		- Shouldn't they run standard protocols wherever possible?
	- Why IPv6?
		- Clean redesign of IPv4 makes it easier to use to support sensor networks
		- Address structure makes device addresses more compressible
- Benefits
	- Interoperability with normal computers and networks
	- Reuse state of the art developed standards instead of remaking them
		- Security, Naming, Discovery, Services
- Costs
	- Packet overhead can be high (will fix)
	- Complexity for supporting protocols

Hui and Culler, "[IP is Dead, Long Live IP for Wireless Sensor Networks](https://patpannuto.com/papers/hui2008ipdead.pdf)". 2008

# Background: IPv6

- Replacement to Internet Protocol v4
	- (Something unrelated used version number 5)
- Extended addressing for devices
	- 32-bits for IPv4 addresses -> 128-bits for IPv6 addresses
	- **340 trillion, trillion, trillion** addresses
		- (> 100 times number of atoms on the surface of the Earth)
	- Example: a39b:239e:ffff:29a2:0021:20f1:aaa2:2112
- Supports multiple transmit models
	- Broadcast: one-to-all
	- Multicast: one-to-many
	- Unicast: one-to-one
- Various other improvements

#### Background: IPv6 address notation rules

- Groups of zeros can be replaced with "::"
	- Can only use "::" in one place in the address
- Leading zeros in a 16-bit group can be omitted

0000:0000:0000:0000:0000:0000:0000:0001 → ::1  $2345:1001:0023:1003:0000:0000:0000:0000 \rightarrow 2345:1001:23:1003:$ aecb:0222:0000:0000:0000:0000:0000:0010 → aecb:222::10

- Special addresses
	- Localhost ::1 (IPv4 version is 127.0.0.1)
	- Link-Local Network fe80:: (bottom 64-bits are ~device MAC address)
	- Local Network fc00:: and fd00::
	- Global Addresses 2000:: (various methods for allocating bottom bits)

# Background: IPv6 datagram format



- Header starts with various data for the router
- **Priority**: like "type of service" in IPv4.
- **Flow label**: marks packets that are part of a transaction for routers
- **Next header**: TCP, UDP, ICMP, etc.
- **Hop limit**: hops to live

# Fitting an IPv6 datagram in an 802.15.4 frame



- How many bytes for the IPv6 header?
- How much for the MAC header + footer?
- How much room do we get for payload?

# Fitting an IPv6 datagram in an 802.15.4 frame



- How many bytes for the IPv6 header?
- How much for the MAC header + footer?
- How much room do we get for payload?
- $-8+16+16 = 40$  bytes
- 5 to 25 bytes depending on options
- $-127-40-N = 82$  to 62 bytes

#### IPv6 tricky to use directly with 802.15.4

- Header overhead very high
	- around half of bytes in PSDU, even more when security enabled
- IP packets may be large, compared to 802.15.4 max frame size
	- IPv6 requires all links support 1280 byte packets [RFC 2460]

• Need some way to compress and fragment!

# 6LoWPAN

- Method for running IPv6 over 802.15.4 links
	- IPv**6** over **Lo**w-Power **W**ireless **P**ersonal **A**rea **N**etworks
	- IETF Standard [\(RFC4944](https://tools.ietf.org/html/rfc4944) + updates in [RFC6282\)](https://tools.ietf.org/html/rfc6282)
- Directly out of the research world (Jonathan Hui + David Culler)
	- Research Paper: [IP is Dead, Long Live IP for Wireless Sensor Networks](https://patpannuto.com/papers/hui2008ipdead.pdf)
	- Thesis of work: sensor networks can and should use IPv6
- Important goals
	- Compress IPv6 headers
	- Handle fragmentation of packets
	- Enable sending packets through mesh

#### 6LoWPAN header compression

- 40 bytes of IPv6 header are a lot for a 127-byte payload
- Most important goals
	- Communication with devices **inside** the 15.4 network should be low-overhead
	- Communication **outside** of the 15.4 network must be possible
		- Still minimize overhead where possible, but might be larger
- Assume a bunch of common parameters to save space
	- A bunch of options are set to default values
	- Payload length can be re-determined from packet length
	- Source/Destination addresses can often be reassembled from link layer data
		- Plus information about network address assignment known by routers
- Border router "inflates" the packet before sending externally



- Example of HC1 header compression
	- Note: Thread actually uses IPHC from rfc6282 (not HC1), but similar idea<br>• Good overview: https://www.youtube.com/watch?v=wRdrCsqJekM
	-

#### 6LoWPAN fragmentation

- Only the first packet of the fragments will hold the IPv6 header
	- Tag, offset, and size are used to reconstruct



Figure 14. Fragmenting and Reassembling an IPv6 Packet

# 6LoWPAN mesh forwarding

• Additional header with originator and final addresses



**Figure 17. Mesh Header Format** 

• Which of these headers are used depends on the packet



Figure 4. 6LoWPAN stacked headers

#### Sidebar: IPv6 over BLE

• [RFC7668](https://tools.ietf.org/html/rfc7668) defines 6LoWPAN techniques for BLE connections



Figure 3: IPv6 and IPSS on the Bluetooth LE Stack

Benefit to IPv6: multiple address spaces per Thread device

- Each device gets an IPv6 address for each way to contact it
	- Global IP address
	- Mesh-local IP address
	- Link-local IP address
	- Topology-based IP address
		- Send to parent
		- Send to child
	- Role-based IP address(es)
		- Send to all Routers
		- Send to Border Router



# Traditional addresses in Thread

- Link-Local Addresses
	- FE80::/16
	- Bottommost 64-bits are EUI-64 (MAC address with 0xFFFE in the middle)
	- Permanent for a given device (no matter the network)
	- Used for low-layer interactions with neighbors (discovery, routing info)
- Mesh-Local Addresses
	- FD00::/8 (FD00:: and FC00:: are for local networks in IPv6)
	- Remaining bits are randomly chosen as part of joining the network
	- Permanent while connection is maintained to a network
	- Used for application-layer interactions
- Global Addresses
	- 2000::/3 (2000:: are for global, unicast IP addresses in IPv6)
	- Public address for communicating with broader internet through Border Router
	- Various methods for allocation (SLAAC, DHCP, Manual)

Topology-based addresses in Thread

- FD00::00FF:FE00:RLOC16
	- Same top bits as mesh-local
- Routing Locator (RLOC)
	- Router ID concatenated with Child ID

Router ID R **Child ID**  $00000100000000001 \rightarrow 0x401$ 

- Changes with network topology
	- Used for routing packets



#### Role-based addresses in Thread

- Multicast
	- FF02::1  $-$  link-local, all listening devices
	- FF02::2 link-local, all routers/router-eligible
	- FF03::1 mesh-local, all listening devices
	- FF03::2 mesh-local, all routers/router-eligible
- Anycast
	- FD00::00FF:FE00:FC**xx**
		- 00 Thread Leader
		- 01-0F DHCPv6 Agent
		- 30-37 Commissioner
		- etc.

#### Break + Question

Which type of address(es) could you use for communication?

- 1. Global
- 2. Mesh-local
- 3. Link-local
- 4. Topology
- 5. Role-based
- Communicate with each circled target



# Break + Question

- 1. Green node
	- Mesh-local
	- Topology (parent)
	- Link-local
- 2. Blue node
	- Mesh-local
	- Topology (other child of parent)
	- Maybe link-local
- 3. Orange node
	- Mesh-local
	- Role-based
- 4. Purple "cloud" • Global



- 1. Global
- 2. Mesh -local
- 3. Link-local
- 4. Topology
- 5. Role -based



# **Outline**

- Thread Overview
- Thread Addressing
- **Runtime Behavior**
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# Thread device types

- Full Thread Devices
	- MAC: Always-on
	- Participate in routing
	- Examples:
		- Routers
		- Thread Leader
		- Border Router
		- Router-Eligible End Device (REED)



#### Router promotion

- Router Eligible End Device (REED)
	- A router without any children
	- Can operate as an end device with one connection (lower power)
	- Promotes to a router when a joining end device relies on it
		- If there is room for an additional router (max 32, typical 16-23)



# Thread device types

- Full Thread Devices
	- MAC: Always-on
	- Participate in routing
	- Examples:
		- Routers
		- Thread Leader
		- Border Router
		- Router-Eligible End Device (REED)
- Minimal Thread Devices
	- MAC: Can be low power
	- Only send/receive to/from parent
	- Examples:
		- Minimal End Device (MED)
		- Sleepy End Device (SED)
		- Synchronized Sleepy End Device (SSED)



# Sleepy End Devices

- These are the low power devices
- Periodically must exchange messages with Router
- Sleepy End Device (SED)
	- MAC:
		- TX: Router always-on, so send as needed (ALOHA)
		- RX: Periodically check in with router if new data pending
- Synchronized Sleep End Device (SSED)
	- MAC:
		- TX: ALOHA
		- RX: Listen during pre-configured windows



Comparison of Thread SED and Thread SSED radio activity

# Discovering Thread networks

- "Beacon Request" MAC command
	- Routers/Router-eligible devices respond
	- Payload contains information about network
- Thread network specification
	- PAN ID 16-bit ID
	- XPAN ID extended 64-bit ID
	- Network Name human-readable
- Active scanning across channels can quickly find all existing nearby networks



# Creating a new network

• Select a channel (possibly by scanning for availability)

- Become a router
	- Elect yourself as Thread Leader
	- Respond to Beacon Requests from other devices
- Further organization occurs through Mesh-Level Establishment protocol

# Mesh-Level Establishment (MLE)

- Creating and configuring mesh links
	- Payloads placed in UDP packets within IPv6 payloads
- Commands for mesh
	- Establish link
	- Advertise link quality
	- Connect to parent



OR (secure version)



- TLVs (Type-Length-Value)
	- Various data types that may be helpful within those packets
	- Addresses, Link Quality, Routing Data, Timestamps

# Joining an existing network

- All devices join as a child of some existing router
- 1. Send a Parent Request (to all routers/router-eligible)
	- Using the multicast, link-local address
- 2. Receive a Parent Response (from all routers/router-eligible separately) • Contains information on link quality
- 3. Send a Child ID Request (to router with best link)
	- Contains parameters about the new child device
- 4. Receive a Child ID Response (from that router)
	- Contains address configurations

**SED** 





# Thread Commissioning

• As part of joining a network, a thread device is also "commissioned"

• Bonus step in the joining process for real-world networks

• Authenticates the joining device allowing it to join and encrypt communications



#### Updates to Thread specification

- Thread version 1.2
	- Simpler, lower-power end devices
		- Any packet sent is enough to keep the device attached to network
		- All acknowledgements identify if packets are available for it
	- Power control
		- Reduce transmission power to minimum needed for "good connection"
		- Reduces potential for "hidden terminal problem"
- Thread version 1.3
	- Better IPv6 support
		- Coordinate address selection with IP network boarder router connects to
		- Shuttle DNS Service Discovery messages between networks
	- Support for TCP

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# Communicating with IP

- Any communication that layers on top of IP is now possible
	- If there is a library to support it
- Common choices
	- UDP
		- DNS translate hostnames into IP addresses
		- SNTP get real-world time, accuracy better than 1 second
		- CoAP send and receive data

# Constrained Application Protocol - CoAP

- HTTP, but over UDP targeting less-capable devices
	- Same REST architecture
	- Adds capability for automatic retransmissions



#### • CoAP Requests

- Have a type: GET, POST, PUT, DELETE
- Have a URL: /file/etc
- Have data up to 65 KB

Sensor networks don't use TCP (yet?)

- Uncommon choice: TCP
	- Concerns: Too large, too slow, poorly suited to lossy networks
	- Also concerning: We're just replicating TCP poorly
	- Active research:
		- Sam Kumar, Michael Anderson, Hyung-Sin Kim, David Culler. "[Performant TCP for Low-Power Wireless Networks](https://people.eecs.berkeley.edu/~samkumar/papers/tcplp_nsdi2020.pdf)". 2020.
		- The debate is still very much open

- 2023 update: Thread now supports TCP!!
	- Primarily for large data payloads, like firmware updates

# A problem: the siloed internet of things

- Problem: companies are more interested in selling you the whole stack
	- Which then makes it harder for devices to be interoperable
- This is not Thread or IP-specific, but a problem all IoT devices are facing
- Concerns
	- What IP address do you send data to?
	- Manufacturer's server is an obvious choice
- We'll talk about the Matter standard as an approach to this issue



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