

Lecture 11

WiFi MAC

CS397/497 – Wireless Protocols for IoT
Branden Ghena – Spring 2022

Materials in collaboration
with Pat Pannuto (UCSD)

Today's Goals

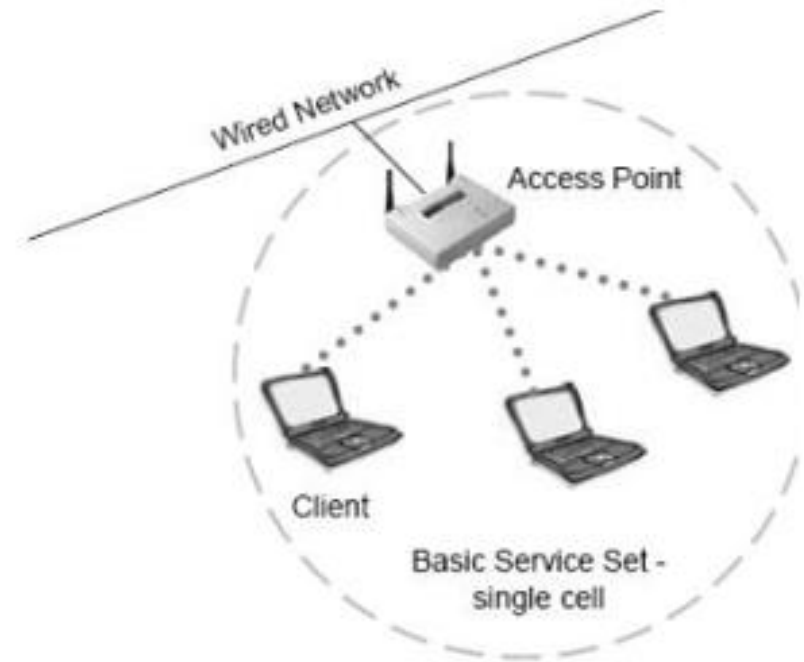
- Introduce MAC layer concepts in 802.11
- Understand what exists, what is actually used, and why
- Explore two additional areas in 802.11
 - Microcontroller use of WiFi
 - Future of WiFi

Outline

- **802.11 Access Control**
- 802.11 Frame format
- 802.11e Improvements to MAC
- Microcontrollers and WiFi

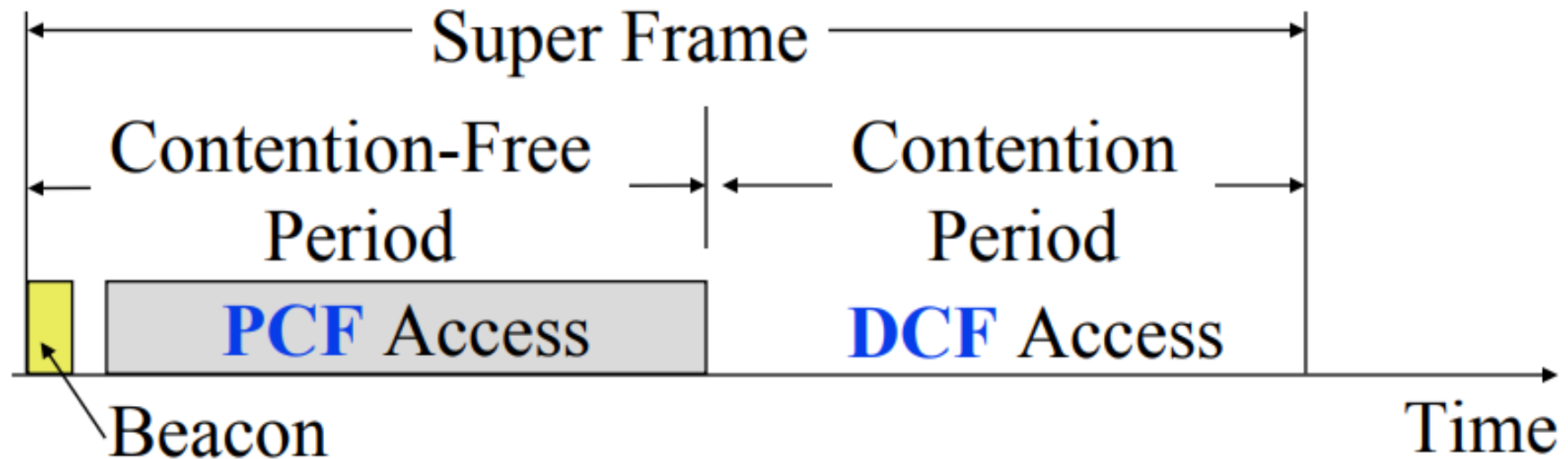
Basic WiFi network

- Star topology network
- Basic Service Set (BSS)
 - Access point(s)
 - Multiple connected clients
- Service Set ID (SSID)
 - Identifies network
 - Broadcast by access point in beacons



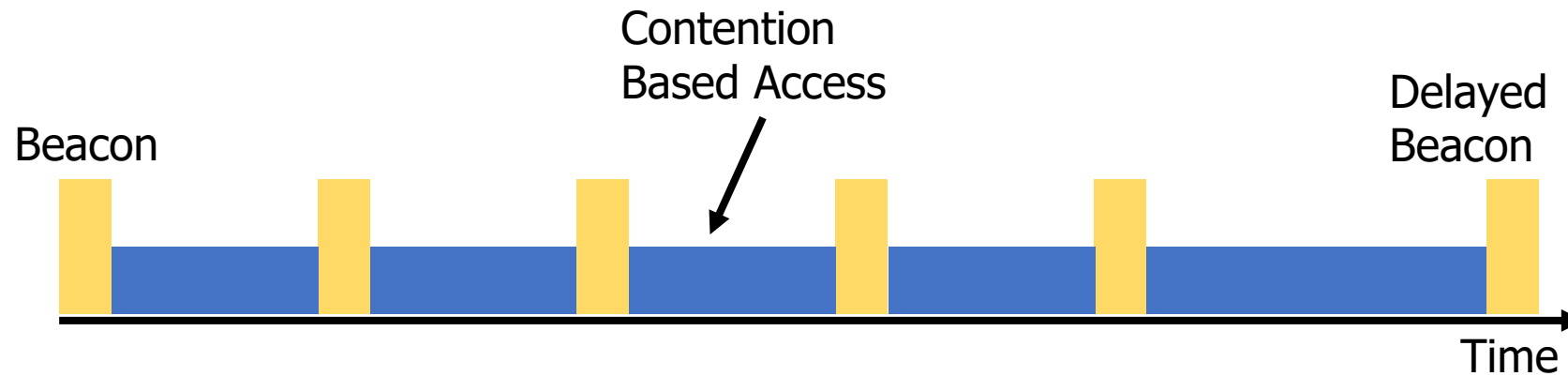
WiFi superframe structure

- Beacon followed by contention-free period followed by contention
 - Repeats periodically (default ~ 100 ms)
 - 802.15.4 adopted a similar superframe
- This is more hypothetical than real



WiFi superframe in practice

- Continuous contention access period
 - Any device may send at any time
 - PCF is unused in practice
- Periodic beacons
 - Which also use CSMA and therefore may be delayed



802.11 beacons

- Transmitted periodically (~ 100 ms by default)
 - Enable discovery of network
 - Contain capabilities and SSID for the network (802.11b/g/n/ac/ax...)
 - Assign contention-free slots if used
 - Notify devices of waiting packets
 - Traffic Indication Map (TIM) has a bitmap specifying which devices data is for
 - Enables devices to sleep, skipping a number of beacons
 - Handles broadcast/multicast messages
 - Every N beacons includes a notation of available broadcast messages
 - Messages are transmitted during next contention access period using normal CSMA
 - Defines maximum sleep period for devices (must listen to these beacons)

Contention-free access

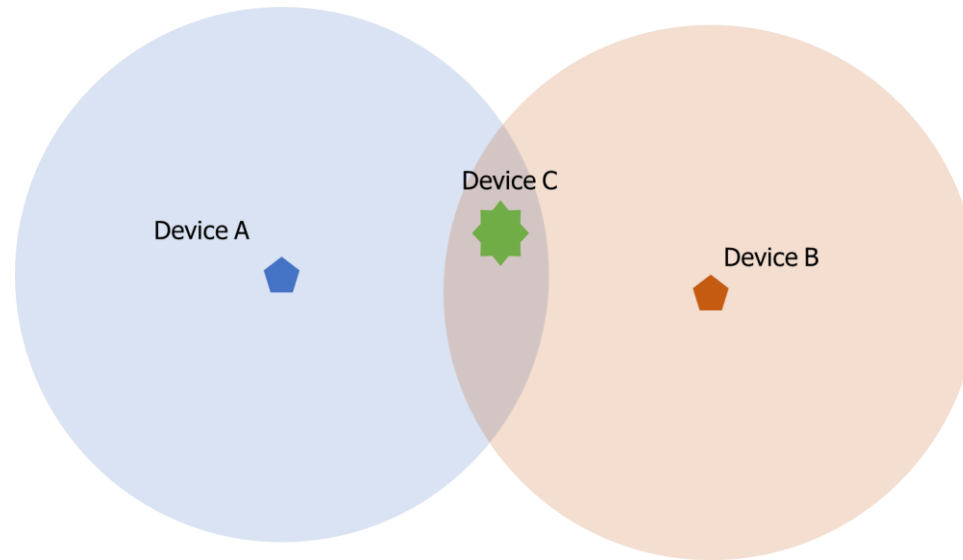
- Known as Point Coordination Function (PCF)
 - Allocates a contention-free period for specific devices
 - Access Point decides when to grant based on requests
- Drawbacks
 - Latency depends on beacon intervals
 - Mechanism for explicit Quality of Service is unclear
- PCF is not used in practice

Contention-based access

- Known as Distributed Coordination Function (DCF)
 - Base communication method for WiFi (essentially always)
 - All packets are immediately ACK'd by receiving device
- Uses CSMA/CA to determine when it can send
 - With random backoff
- Problem: packets can be very long (up to 20 milliseconds)
- Solution: Network Allocation Vector (NAV)
 - Packets include a notation of their duration
 - Sensing the beginning of a packet allows backoff to skip the whole packet duration before continuing

Reminder: hidden terminal problem

- Two devices communicating with Access Point may not be able to hear each other
 - CSMA fails and Access Point losses both messages



- A solution: RTS/CTS (Request/Clear To Send)

Drawbacks of RTS/CTS

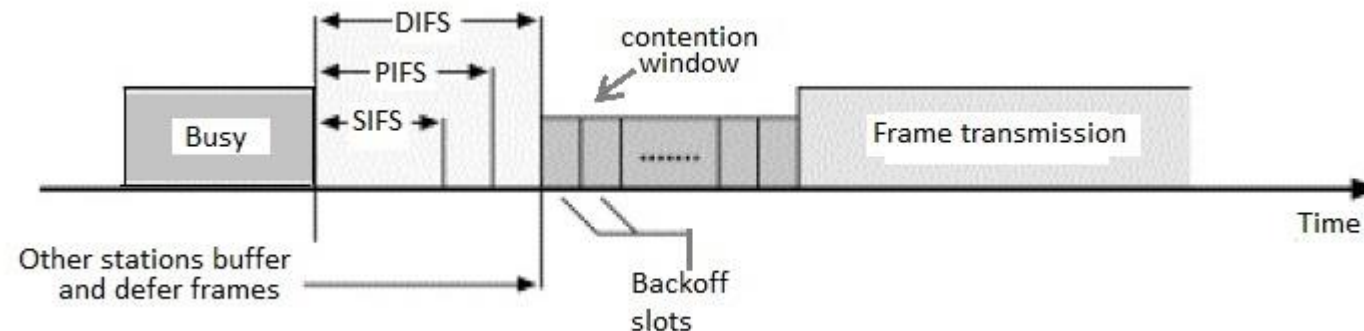
- Four packets per data (RTS, CTS, Data, Ack)
 - Could have just sent data instead of RTS
- Significant portion of traffic are application-layer Acks
 - Probably better to just have it fail and try again later
- RTS/CTS only used for very large packets in practice
 - *It's mentioned still in 802.11n and 802.11ac, so not entirely unused

Backoff in WiFi

- Listen for activity
 - If free
 - Wait for Inter Frame Spacing (IFS)
 - If still free, transmit
 - If busy
 - Randomly select a number of backoff **Slots**
 - Count down slots whenever medium is not busy
 - If busy when backoff completes:
 - Increase maximum backoff Slots
 - Repeat
- Slot time: basic time unit for protocol
 - Total time of: switch from Rx to Tx, plus processing time, plus propagation delay

Prioritizing packets with varying IFS

- Tiered Contention Multiple Access (TCMA)
 - Idea: assign different inter-frame spacing based on traffic class
 - Inherently prioritizes communication
- Acknowledgements sent with Short IFS (SIFS)
 - Will always transmit before new data clears CSMA check
- New data sent with DCF IFS (DIFS)

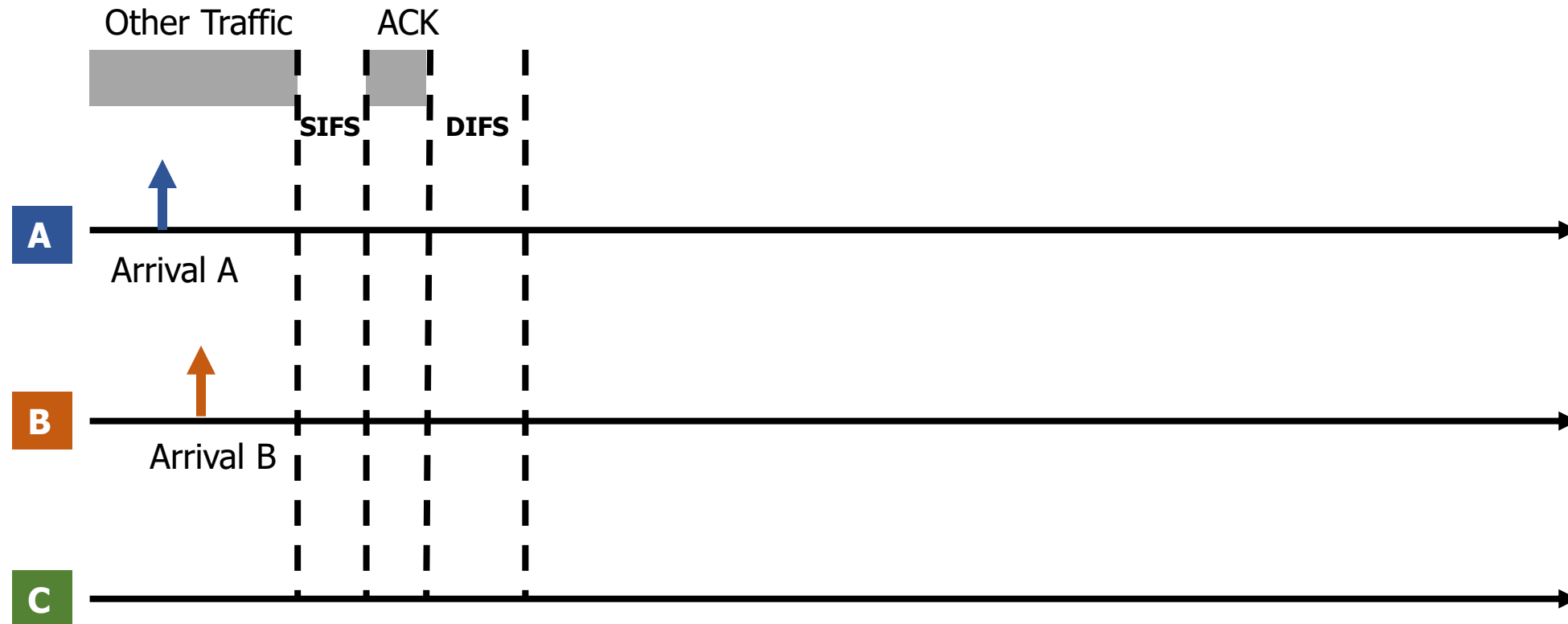


Putting backoff together

- Two variables
 - Contention Window (CW) – maximum backoff amount
 - Backoff Count (BO) – current remaining backoff
- When attempting to send, if busy Backoff selected in $[0, CW]$
 - Countdown Backoff slots whenever medium is not busy
 - At 0, attempt to transmit if not busy
 - If busy, double Window and select Backoff again
- 802.11g values:
 - Slot time= 20 us, CWmin= 15 slots, CWmax= 1023 slots
 - SIFS= 10 us, DIFS= 50 us

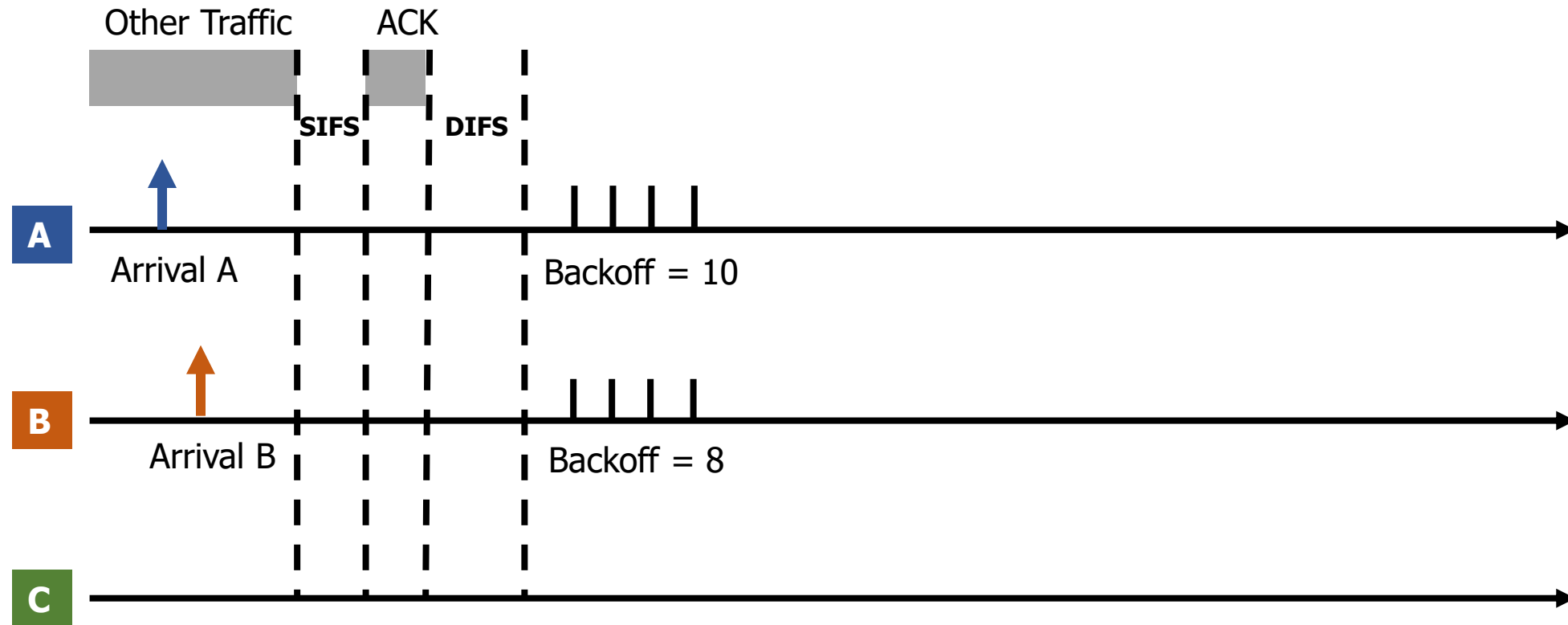
802.11 backoff example

- A and B want to send and see the medium is busy
 - Followed by an Acknowledgement after SIFS



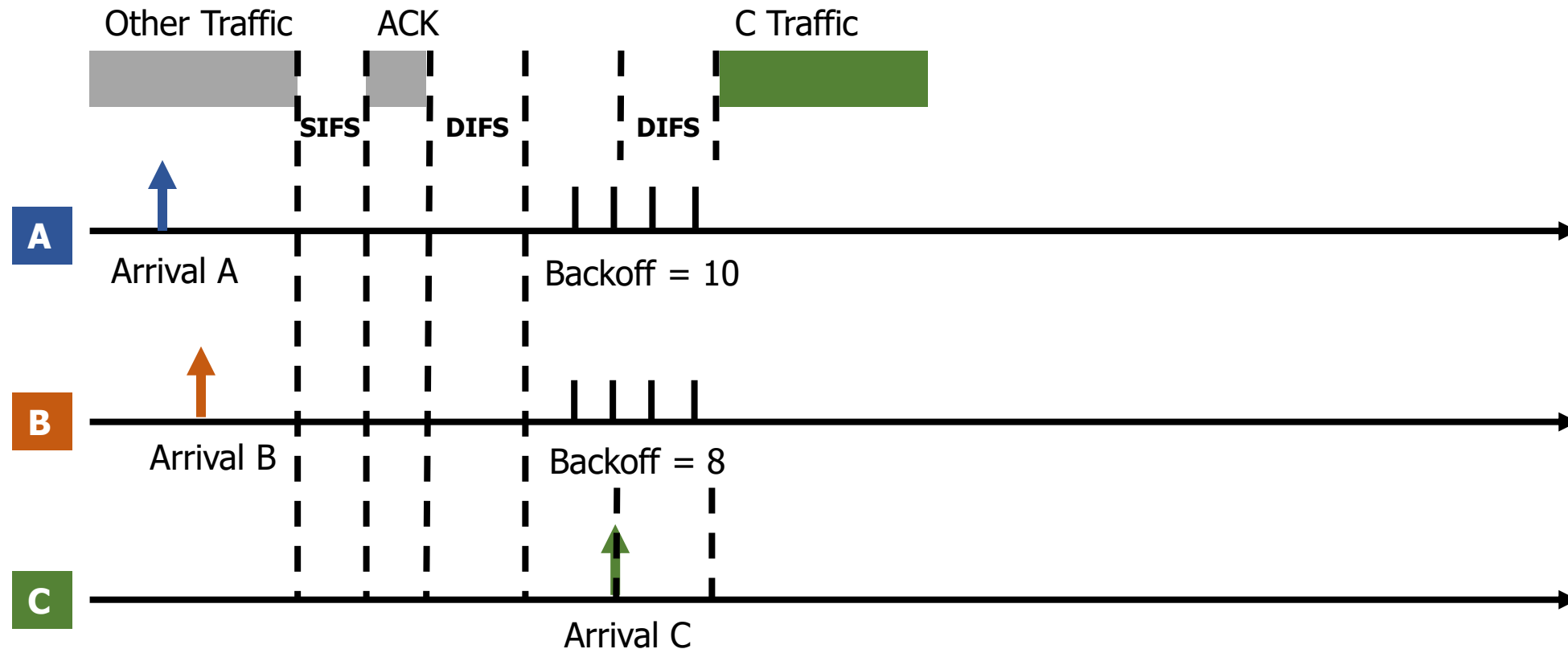
802.11 backoff example

- Each chooses a random backoff [0, CW] (we'll say CW is 32)
 - Start counting down backoff slots



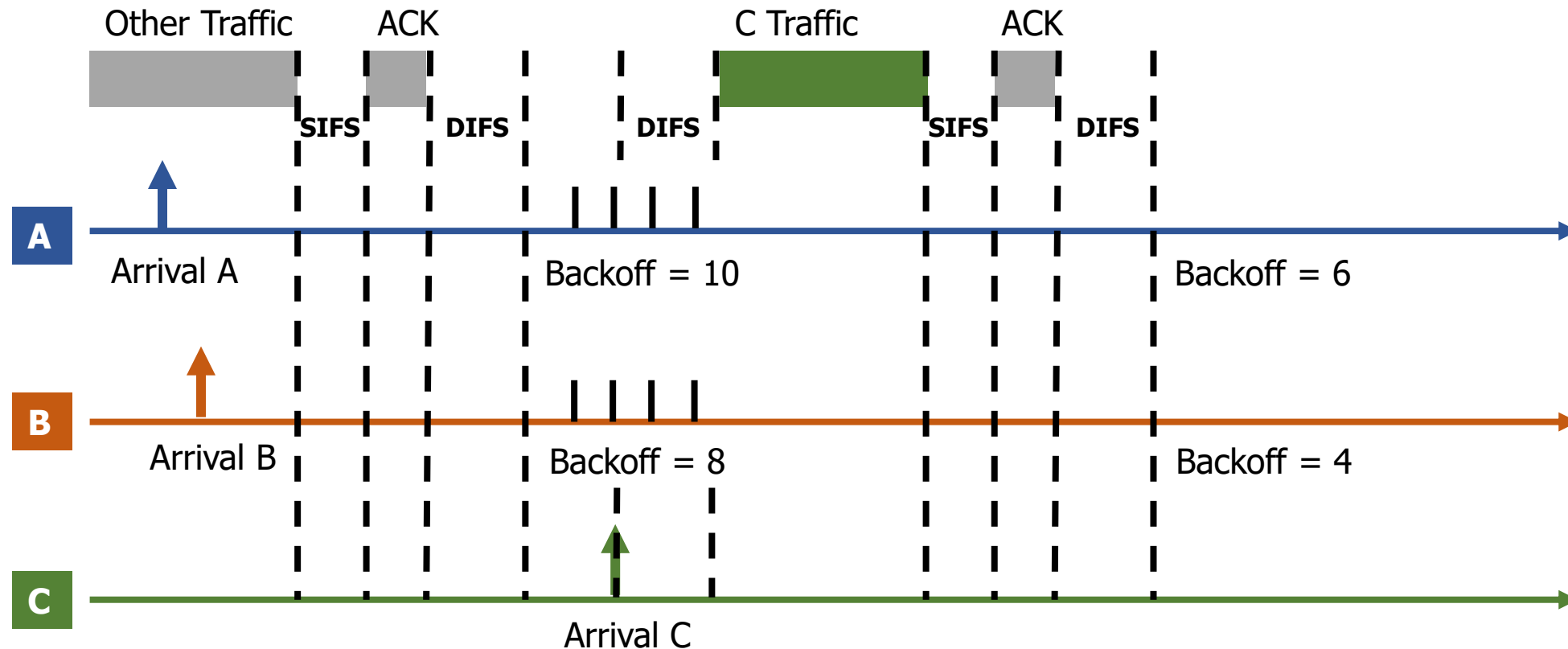
802.11 backoff example

- C wants to send, waits DIFS, and can send immediately
 - No other traffic is going on
 - A and B pause backoff for packet duration



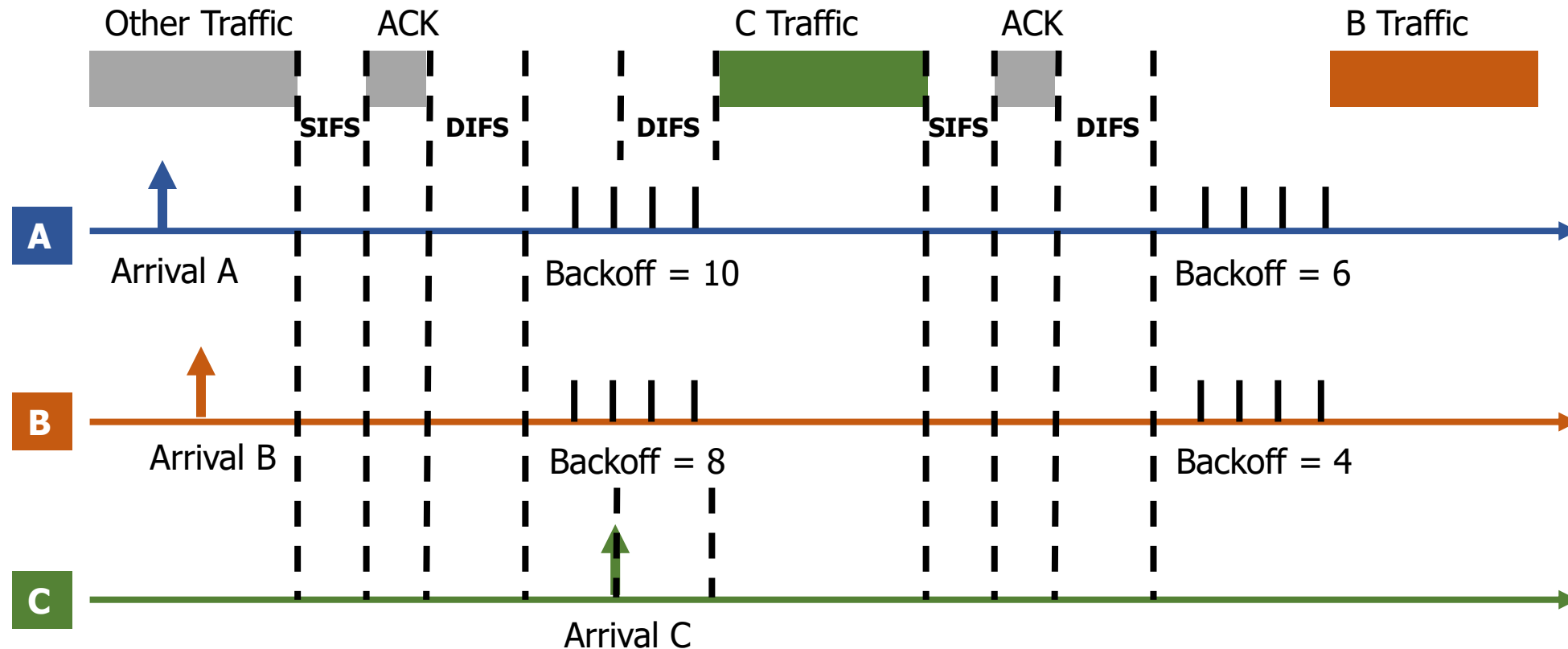
802.11 backoff example

- A and B used NAV to pause backoff for entire traffic plus ACK
 - After DIFS, resume backoff count from its previous value



802.11 backoff example

- B reaches zero backoff, finds channel empty, transmits
 - A pauses its backoff again for duration plus ACK



Break + Hacking

- If you wanted maximum data throughput on a WiFi radio, and you were willing to be non-standards-compliant, what would you do?

Break + Hacking

- If you wanted maximum data throughput on a WiFi radio, and you were willing to be non-standards-compliant, what would you do?
 - Never backoff at all. Just try during the next open period
 - Always be “device C” in our previous example
 - Use a shorter SIFS than other devices
 - If you start transmitting sooner, you get to keep transmitting!
 - Other devices will backoff on your transmission
 - Tragedy of the Commons: this utterly fails if many radios follow it

Outline

- 802.11 Access Control
- **802.11 Frame format**
- 802.11e Improvements to MAC
- Microcontrollers and WiFi

802.11 frame

Field	Frame control	Duration, id.	Address 1	Address 2	Address 3	Sequence control	Address 4	QoS control	HT control	Frame body	Frame check sequence
Length (Bytes)	2	2	6	6	6	0, or 2	6	0, or 2	0, or 4	Variable	4

- Frame control (various bits)
 - Type of packet (Control, Management, Data)
 - Subtype (Association, RTS, CTS, Ack, etc.)
 - Indication of to/from “distribution system” (Internet rather than intranet)

- Duration

- Specifies on-air time of full packet in microseconds
- Note: no actual length field 🤖

Surprising, but smart!

Recall MCS vary — but everyone needs to be able to parse header (for duration, for NAV)

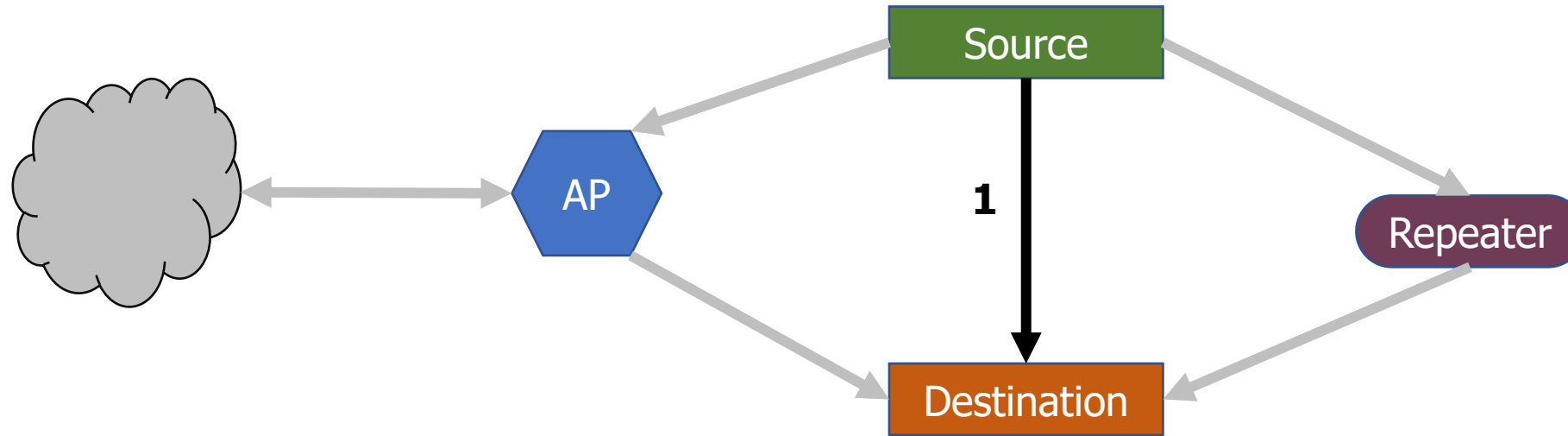
Length can be very large (e.g. in ac: 5.5 ms max duration is 4.5 MB length!); sent at full data rate

802.11 frame

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- Sequence control
 - 4-bit fragment number
 - 12-bit sequence number
- Quality of Service control
 - Identifies traffic category
- High Throughput Control
 - Configurations for selecting best data rate
- Frame body
 - Max size depends on PHY
 - ~2000 for lower rates
 - ~8000 for 802.11n
 - ~11000 for 802.11ac
- Frame check sequence
 - 32-bit CRC

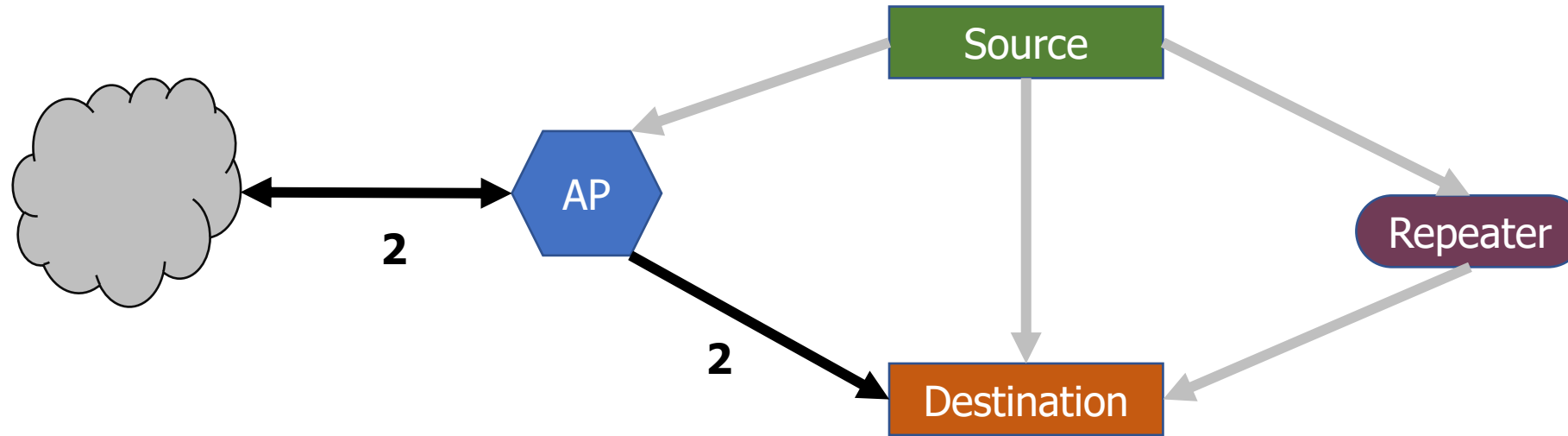
Address field use cases



Devices filter on Address 1

	To DS	From DS	Address 1	Address 2	Address3	Address4	Use Case
1	0	0	Destination Addr	Source Addr	BSS ID	-	Direct communication
2	0	1	Destination Addr	BSS ID	Source Addr	-	Traffic from Internet
3	1	0	BSS ID	Source Addr	Destination Addr	-	Traffic to Internet
4	1	1	Receiver Addr	Transmitter Addr	Destination Addr	Source Addr	Repeater

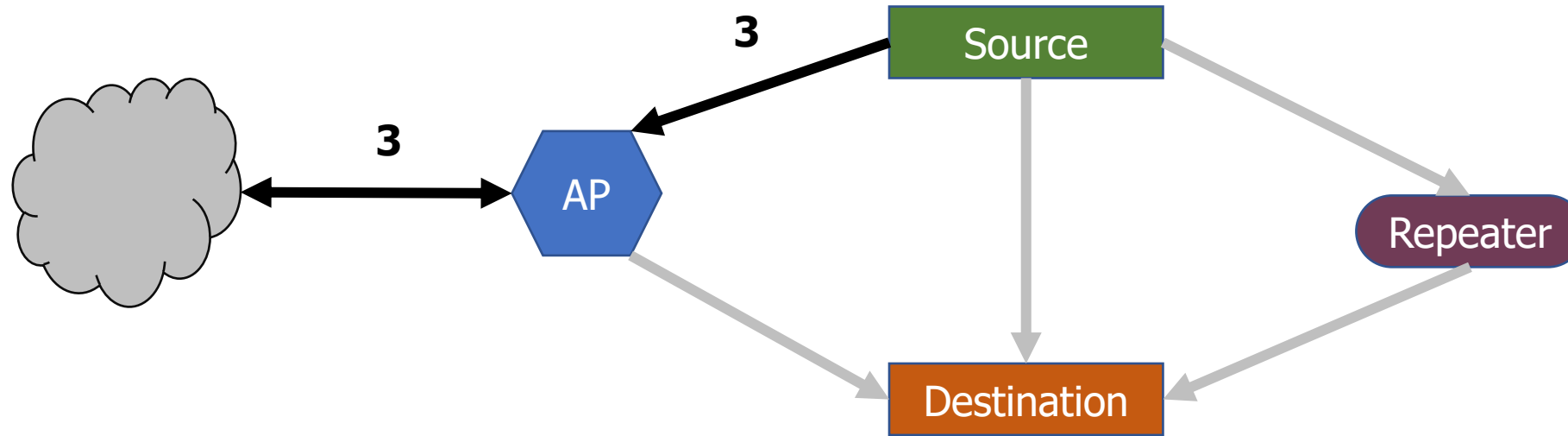
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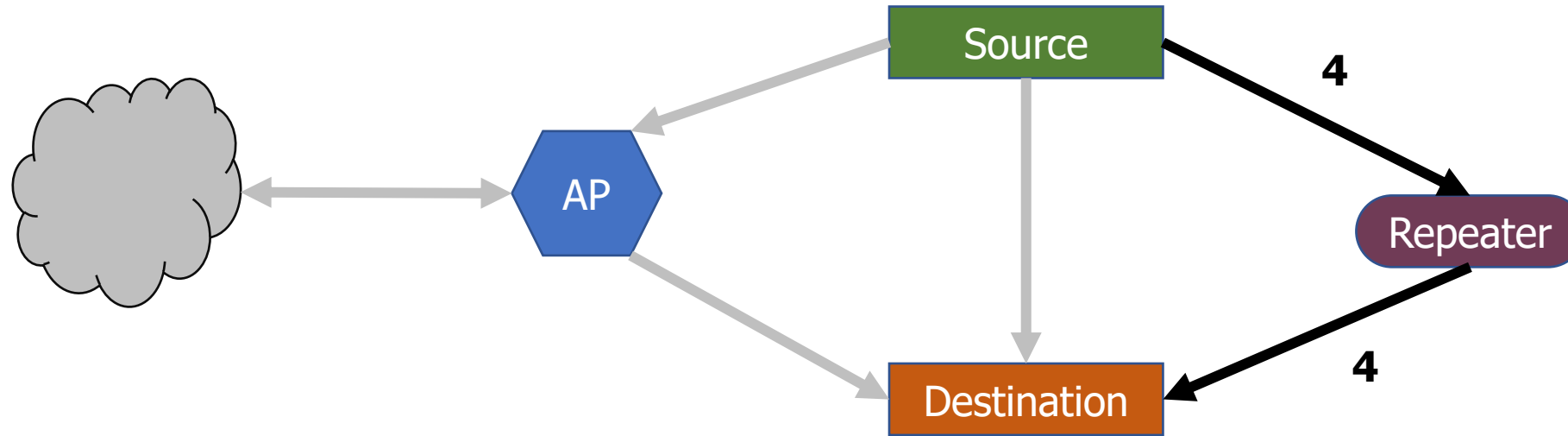
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Address field use cases



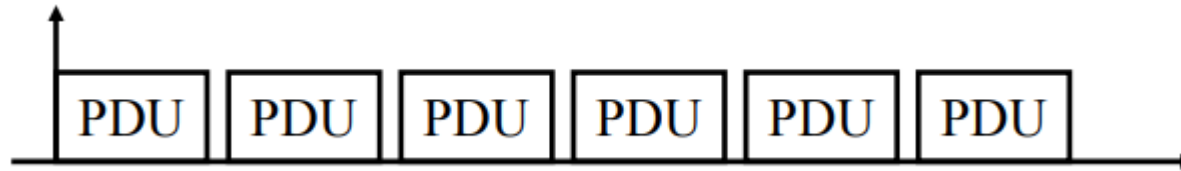
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Sending frames in WiFi

- Frame bursting

- Transmit multiple frames in a row

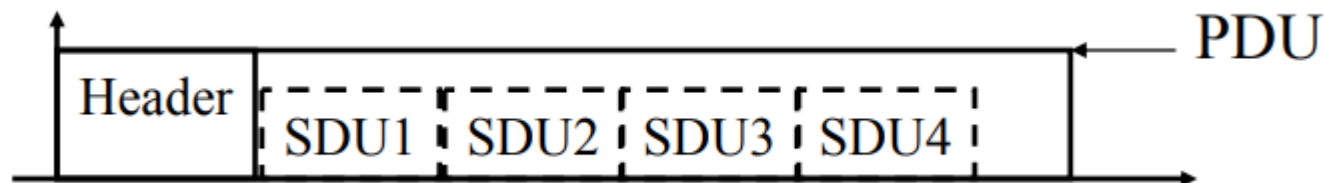


- Frame fragmentation

- Split service data over multiple frames

- Frame aggregation

- Multiple service data in a single frame
- Allows multiple packets to reach Access Point in a single transmission



Calculating packet durations

- Example duration for a 1500 byte 802.11g packet

- 6 Mbps for header
- 24 Mbps for payload
- 566 μ s for total packet
 - Plus 10 μ s for SIFS
 - Plus 34 μ s for ACK

- [https://sarwiki.informatik.hu-berlin.de/Packet transmission time in 802.11](https://sarwiki.informatik.hu-berlin.de/Packet%20transmission%20time%20in%20802.11)

Data transmission bitrate
(802.11g / a*):

	24	
	Mbps	

Bitrate (Mbit/s)	Length (bits)	Time (μ s)
---------------------	------------------	--------------------

			28
D	PHY header: PLCP preamble	-	16
A	PHY header: PLCP header	6	4
T	MAC headers (28 bytes) + MAC body	24	512
A	signal extension time		6

tx time data: 566

			10
A	PHY header: PLCP preamble	-	16
C	PHY header: PLCP header	6	4
K	MAC headers + PHY pad	24	8
	signal extension time		6

tx time ack: 44

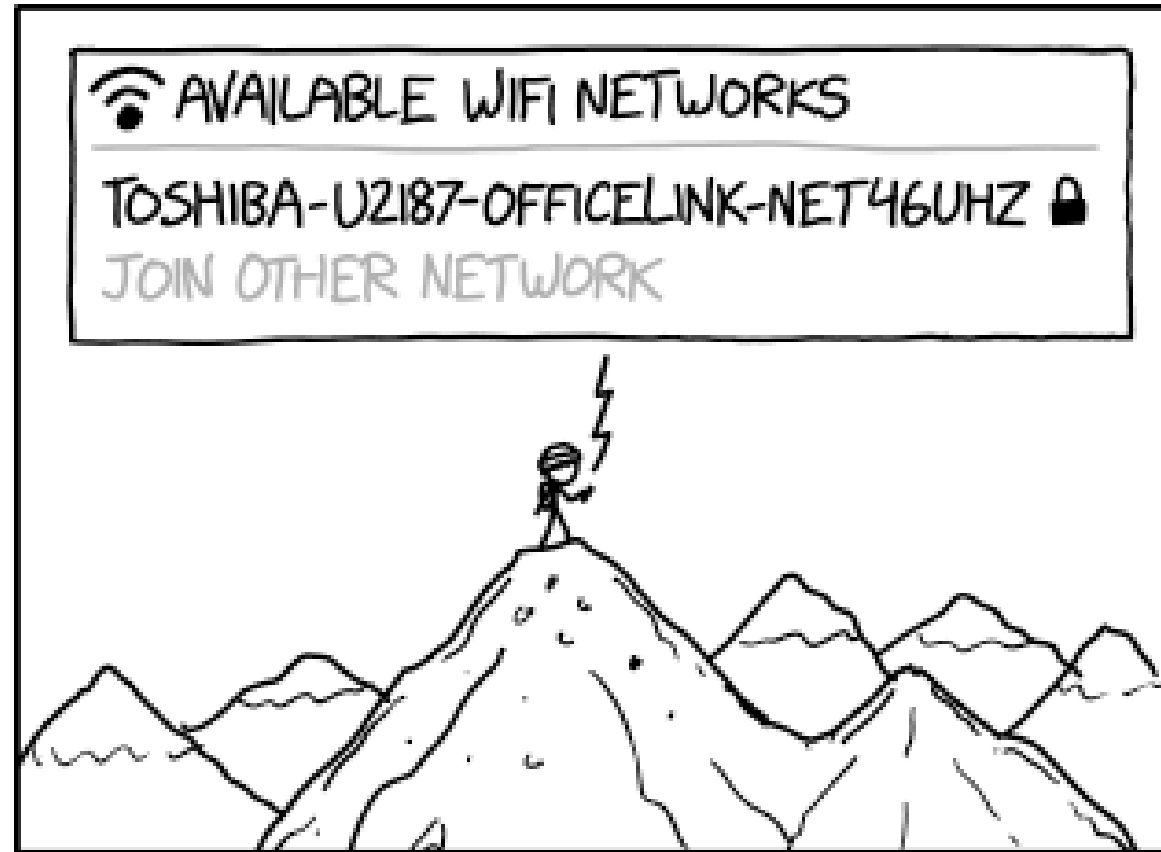
tx time data + ack: 610

Implementation Drives Specification Sometimes

- SIFS nominally defined by processing time
 - Aside: Big challenge for SDRs
- Convolutional decoders need(ed) 16 μ s to finish processing
 - For highest-rate MCS (ERP-OFDM)
 - SIFS is 10 μ s, so extension needed
- Processing must finish before next packet starts
 - To be able to decode NAV in header

		Data transmission bitrate (802.11g / a*):		
		24		
		Mbps		
		Bitrate (Mbit/s)	Length (bits)	Time (μ s)
	DIFS			28
D	PHY header: PLCP preamble	-	-	16
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T	body	24	12246	512
A	signal extension time			6
tx time data:				566
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K	MAC headers + PHY pad	24	134	8
	signal extension time			6
tx time ack:				44
tx time data + ack:				610

Break + xkcd



TECH TRIVIA: NO ONE ACTUALLY KNOWS WHAT DEVICES PRODUCE THOSE CRYPTIC WIFI NETWORKS. THEY JUST APPEAR AT RANDOM ACROSS THE EARTH'S SURFACE.

Outline

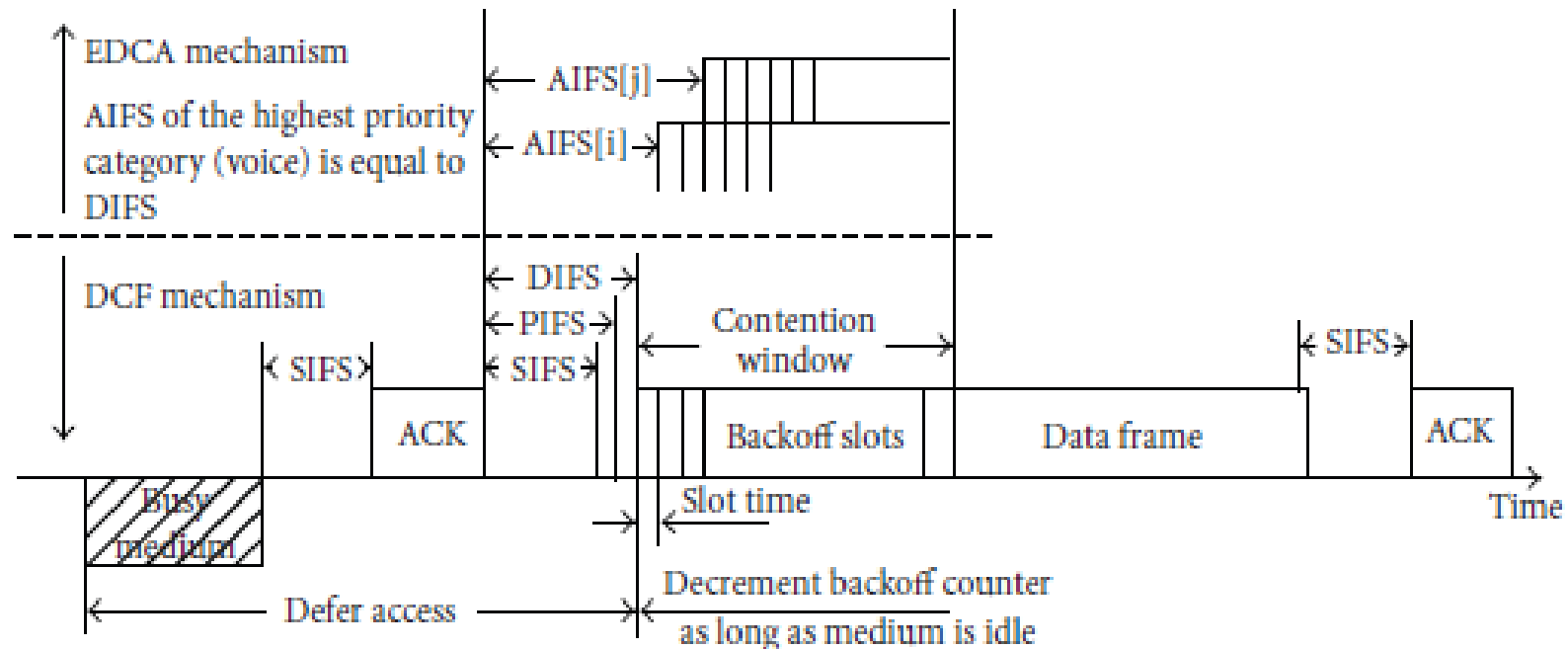
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802.11e improves MAC layer

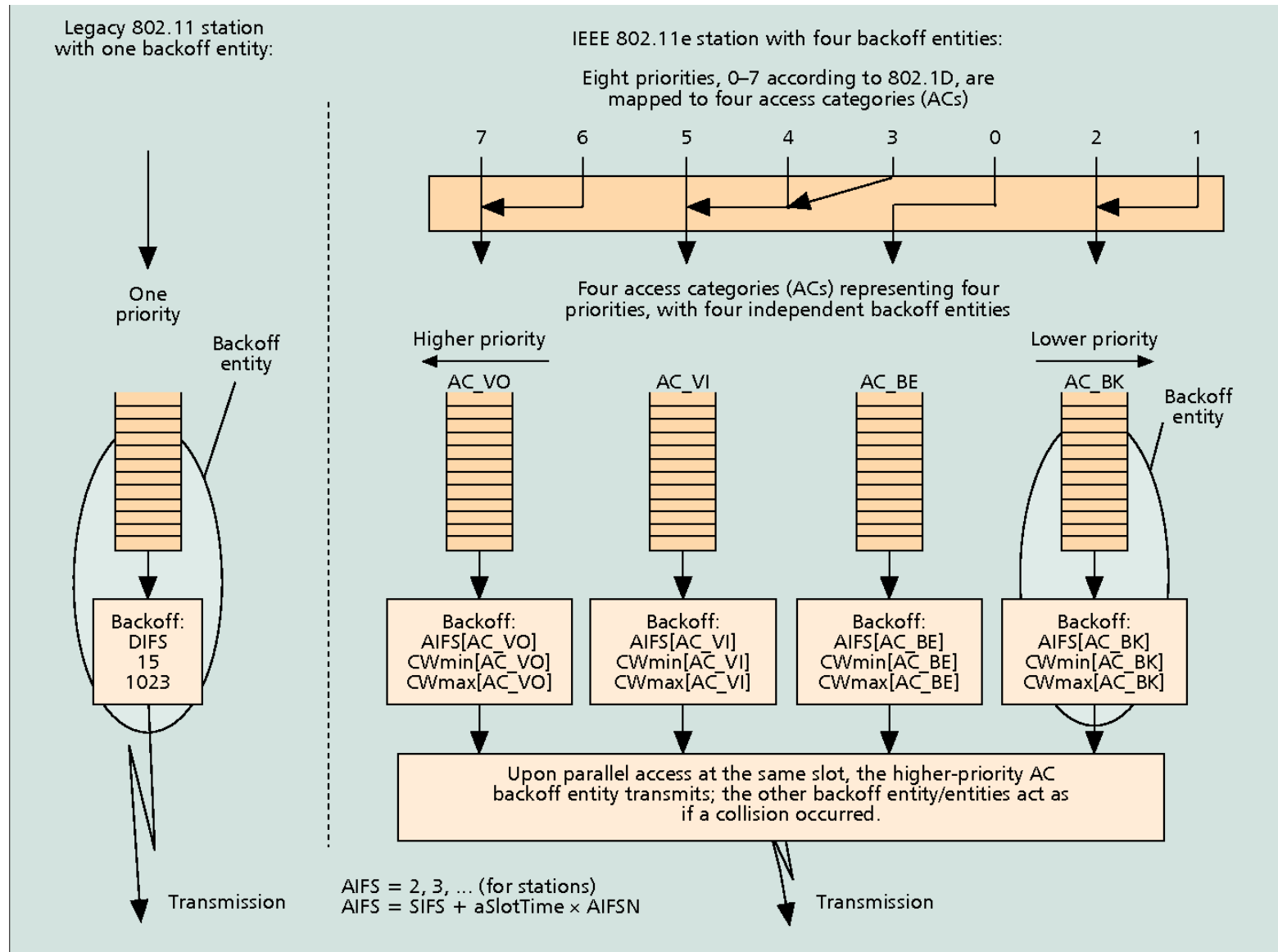
- Hybrid Coordination Function (HCF)
 - Modifies contention-free access (still no one uses it)
 - Modifies contention-based access: Enhanced Distributed Channel Access (EDCA)
- Modifies Quality of Service based on application
 - Example of breaking layering for an optimization
 - Categories (lowest to highest priority):
 - Background
 - Best Effort
 - Video
 - Voice

Different priority for different application category

- Expand to more IFS lengths for different traffic categories
 - Smallest AIFS (equal to DIFS) goes to Voice, Largest to Background
 - Contention Window min and max also change for each category
 - Selects a *probability* that most important category goes first



Multiple queues within a single device



■ Figure 4. [3] Legacy 802.11 station and 802.11e station with four ACs within one station.

802.11e also adds maximum durations

- 802.11e also defines duration a device can transmit for
 - Based on PHY in use and Application category
 - Background/Best Effort: one frame per contention win
 - Example, up to 11 ms for Voice on 802.11ac
 - Could be one really big frame at a low data rate
 - Could be multiple frames in a row separated by SIFS

Outline

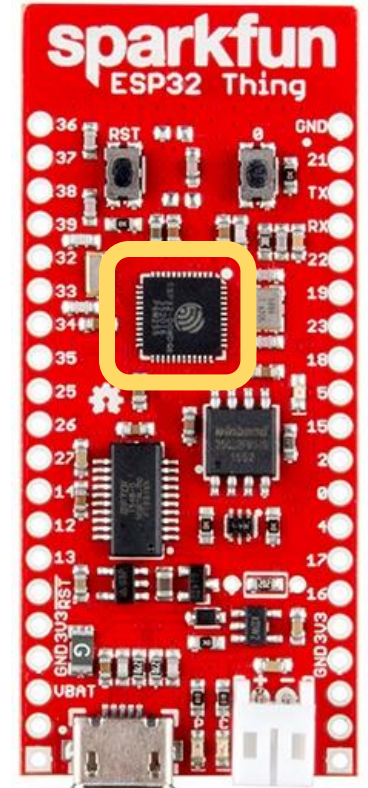
- 802.11 Access Control
- 802.11 Frame format
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- **Microcontrollers and WiFi**

Why, why not, talk WiFi in a wireless for IoT class

- Pros
 - Ubiquitous
 - High-performance
- Cons
 - Complex configuration
 - And security requirements
 - Device-Northwestern anyone?
 - Expensive in energy and money

WiFi capability in microcontrollers

- ESP32
 - Microcontroller plus WiFi radio in single chip
 - (Same idea as nRF52840)
- Capabilities
 - 802.11b/g/n 2.4 GHz only
 - 20 MHz or 40 MHz channels
 - Single antenna only (no MIMO)
 - MCS0-7
 - 7 Mbps – 150 Mbps
 - Tx power up to 20.5 dBm



Low power WiFi

- Question: should a microcontroller stay connected or reconnect?
 - Light sleep: stay connected always, only listening to beacons
 - Deep sleep: reconnect to network each time data is ready
- Answer for ESP32 depends on security and data interval
 - Resecuring during connection takes lots of energy
 - Crossover point is about 60 seconds
 - Insecure transmissions have a crossover of 5-15 seconds

<https://blog.voneicken.com/2018/lp-wifi-esp-comparison/#conclusions>

Wrapup on WiFi

- My takeaway: next time you buy a router, make it WiFi 6E
 - Extra bandwidth with low contention means high speeds
 - Although it won't help until you upgrade devices too
- However: additional WiFi speed won't really help if it's greater than your connection to your ISP
 - 1 Gbps link to router 😁
 - 10 Mbps link to Internet 🤔
- Still useful for local network communication

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