Lecture 01 Introduction

CS213 – Intro to Computer Systems Branden Ghena – Winter 2023

Slides adapted from: St-Amour, Hardavellas, Bustamente (Northwestern), Bryant, O'Hallaron (CMU), Garcia, Weaver (UC Berkeley)

Northwestern

Welcome to CS213!

• In brief: how *does* a computer work anyway?

- We will explore that question across four major sections:
 - **Representations** of information on a computer
 - How the **machine** executes software
 - How **memory** is organized
 - How the **operating system** manages this all for efficiency and security

Branden Ghena (he/him)

- Assistant Faculty of Instruction
- Education
 - Undergrad: Michigan Tech
 - Master's: University of Michigan
 - PhD: University of California, Berkeley
- Research
 - Resource-constrained sensing systems
 - Low-energy wireless networks
 - Embedded operating systems
- Teaching
 - Computer Systems
 - CS211: Fundamentals of Programming II
 - CS213: Intro to Computer Systems
 - CS343: Operating Systems
 - CE346: Microprocessor System Design
 - CS397: Wireless Protocols for the IoT















Today's Goals

• Introduce the theme and goals of the course

• Describe how this class is going to function

• Discuss how a computer system works at a high level

 Begin exploring how computers represent information with bits and bytes

Outline

- Course Themes
- Logistics
- Running a program
- Representing numbers with binary

Convenient computing

- Computers operate on integers, reals, structs, arrays, etc.
- Computers operate on variables and functions
- Computers execute conditionals, loops, etc.
- Memory is an infinite bag of objects my program can allocate
- Memory doesn't have to be shared with any other program
- Memory is always equivalently fast to access
- Etc.

Convenient **illusions** in computing

- Computers operate on integers, reals, structs, arrays, etc.
- Computers operate on variables and functions
- Computers execute conditionals, loops, etc.
- Memory is an infinite bag of objects my program can allocate
- Memory doesn't have to be shared with any other program
- Memory is always equivalently fast to access
- Etc.
- None of these are actually true!
 - But we usually program as if they were, and we get away with it!
 - What's going on?

The power of abstraction

- These illusions are called *abstractions*
- They approximate reality, but leave out details
 - Instead, they provide an *interface* that we can work and think with
- We can forget about those details, and be more productive
- Abstractions we love
 - Abstract data types
 - Asymptotic analysis
 - High-level programming languages
 - Operating systems
 - Etc.

The Limits of Abstraction

- Sometimes, abstractions break down
 - Their implementation is buggy
 - Mismatch between expected interface and implementation
 - Their performance is inadequate
 - We need control over the details they hide
 - Security concerns make these details important
- At that point, details come rushing back
 - Can't pretend they don't exist anymore
 - We must know how to deal with them
- This class prepares you to be ready when that happens

Complicated designs fail in unexpected ways

- Some software engineers at Microsoft came up with a cute way of storing dates
 - Two-digit year, month, date, hour, minute concatenated into a 10-digit number
 - Example: 2005230710 -> May 23, 2020 at 7:10 AM
- Stored as a 32-bit signed number
 - Maximum value: 2147483647
- Result: Starting January 1st, 2022, Microsoft Exchange email servers could no longer send email
 - 2201010001 is greater than the largest 32-bit number
 - Microsoft had to issue an emergency patch

Expectation mismatches lead to real-world problems

- Ariane 5 explosion (1996) double short
 - Inertial reference system converted a 64-bit float to a 16-bit integer
 - Expectation: converting from decimal to whole numbers is safe
 - Had worked in the past in Ariane 4, but Ariane 5 was faster
 - Speed too large to fit in a 16-bit integer -> software fault
 - Reality: rocket explodes

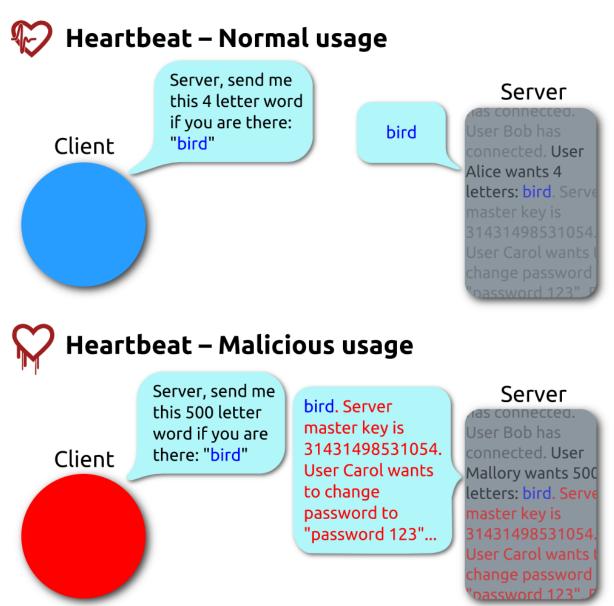


Hardware realities impact software performance

- Abstracted lower-level details can affect performance a lot!
- Question: does the order of iterating through an array matter?
 - Each column in a row OR each row in a column?
- Answer: right code is 10-32 times slower on Intel systems
 - Due to caches

Simple bugs can result in massive vulnerabilities

- 2014 vulnerability in OpenSSL
- Clients can check if server is active by sending a message and listening for echoed response
- C library forgot to check bounds of array and could be abused to return important memory



CS213 goals

1. Break through abstractions to understand how computer processors and memories affect software design and performance

- 2. Introduce concepts of "computer systems" areas:
 - Architecture, Compilers, Security, Embedded, Operating Systems, etc.

Course design goal

- Most systems courses are builder-centric
 - Computer Architecture: design a pipelined processor in Verilog
 - **Operating Systems**: implement portions of an operating system
 - **Compilers**: write a compiler for a simple language
 - **Networking**: implement and simulate network protocols
 - Fun, for sure
 - But ultimately, many more of you will *build on* systems
 - Rather than *build systems* directly
- This course is programmer-centric
 - Purpose is to show that by knowing more about the underlying system, one can be more effective as a programmer
 - Not just a course for dedicated hackers
 - We want to bring out the hacker in everyone!

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Course Staff

- TA (1)
 - Mohammad Kavousi
 - PhD student in Computer Science
- PMs (13):
 - Francis Brenner
 - Adam Chen
 - Elena Fabian
 - Dimitri Hatzisavas
 - Dilan Nair
 - Sean Rhee
 - Evan Waite

Kellen Bryant Huaxuan Chen Joseph Grantham Alex Kang Danny Pineda Santi Roches

Their role: support student questions via office hours and piazza

Course details - how to learn stuff

- Lectures: here in class, Mondays and Wednesdays
 - Please attend and ask questions!
 - Panopto tab on Canvas should have best-effort recordings (a few hours later)
- Textbook:
 - Computer Systems: A Programmer's Perspective 3rd Edition
 - A very useful reference
- Office hours: (starting next week)
 - Planning a mix of in-person and online (gather.town)
 - More info will be posted to Piazza when schedule is ready
 - Can reach out on Piazza to schedule a meeting too

Asking questions

- Class and office hours are always an option!
 - We can do extra questions right after class too
- Piazza: (similar to Campuswire)
 - Post questions
 - Answer each other's questions
 - Find lab partners
 - Find posts from the course staff
 - Post private info just to course staff
- Please do not email me! Post to Piazza instead!
 - I'll be updating roster again a few times

Programming Labs

- Four labs
 - 1. Pack Lab (new) manipulate bits and bytes of a file
 - 2. Bomb Lab deconstruct software to understand it
 - 3. Attack Lab exploit security vulnerabilities in software
 - 4. SETI Lab make software faster with concurrency
- Work on these preferably as a group of two
 - Work together and don't split up assignments (otherwise you won't learn)
 - Individual is acceptable but less good
 - We'll do a pairing survey if you don't already have a partner in mind
- Very different from CS211 style projects
 - Emphasis on the thinking rather than the programming

Homeworks

• Worksheet-style practice problems to help you actually understand what's going on and practice for exams

- Four homeworks that cover class topics
 - 1. Bits and Bytes (releases on Wednesday)
 - 2. Floating Point
 - 3. Assembly
 - 4. Caches

Midterm Exams

- First midterm exam will be during class time
 - Should be back in person well before then
- Second midterm exam will be during exam week
 - **Important:** Wednesday of exam week is our scheduled slot

- Not cumulative, second midterm is second half of class
 - But material in this class builds on itself...

Grades

- Grade breakdown
 - 50% Programming Labs
 - 20% Homeworks
 - 15% Midterm Exam 1
 - 15% Midterm Exam 2

- (4 labs at 12.5% each)
- (4 homeworks at 5% each)

- Exact number to letter mapping is a little flexible
 - But this course is NOT curved

Late Policy

- You can submit homeworks and labs late
- 20% penalty to maximum grade per day late
 - Example: three days late means maximum grade is 40%
- There are exceptions to this:
- 1. We will be flexible with deadlines for problems outside of your control
 - Sick, family emergency, broken computer
 - Contact me (via Piazza)

Slip Days

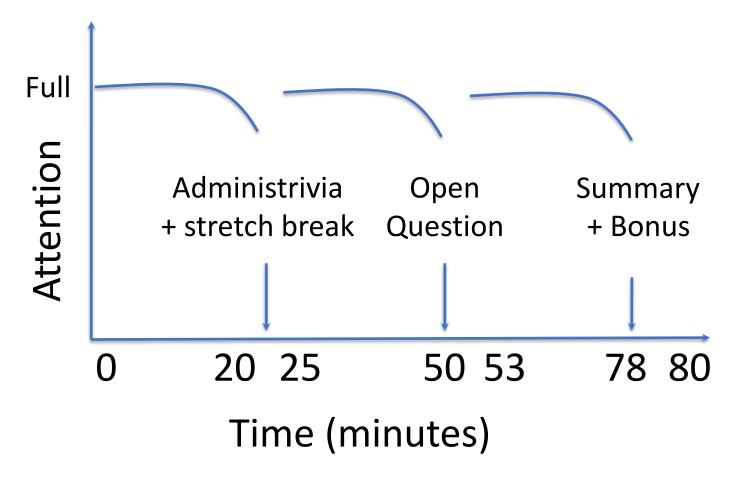
2. Slip days let you turn in a homework late and receive no penalty

- Each student gets **3 slip days**
 - Apply to homeworks and labs
 - You don't need to tell us you're using them, we'll just automatically apply them at the end of the year
 - Be sure to coordinate about them on partner assignments
- Examples:
 - Turn in homework 1 three days late
 - Turn in homework 4 two days late and SETI lab one day late
 - Turn in homework 2 four days late with only a one-day penalty

Academic Integrity

- This is something I take very seriously
- Collaboration good; plagiarism bad
 - You should know where that line is, and be nowhere near it
 - When in doubt, ask the instructor *before* you do something you're not sure about
- At no point should you see someone else's solutions
 - Not your colleagues', not your friends', not your cousin's, not something you found online
- I report everything suspicious to the dean

Break + Architecture of a lecture



Expectations

• This class is **hard**

- And it's hard in a different way. Lots of new material that builds on itself
- You have an opportunity to learn a lot from it
- I'm confident that you can all succeed
 - Labs, Homeworks, Lecture, Office Hours are all designed to support you
- You'll gain a much deeper understanding of how computers operate
 Maybe it's not for you, maybe you'll love it

How to succeed in this class

- Come to lecture
- Ask questions
- Consult the textbook for clarity and practice
- Start assignments early
- Stay on top of the material

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Hello World

- What happens when you run "hello" on your system?
 - And *why* does it happen?



• Goal: introduce key concepts, terminology, and components

Compiling hello

• Compiling hello

unix> gcc -o hello hello.c

- GCC is our compiler
- 1. It takes our source code (hello.c)
 - A text file containing characters
 - Text file = readable by humans
- 2. And translates (compiles) it into **assembly code**
 - A text representation of x86 instructions
 - Here, not explicitly stored in a file
 - We'll be working with assembly a lot this quarter
- 3. Then translates (assembles) that into an executable (hello)
 - A binary file containing x86 machine code
 - Binary file = not meant to be read by humans (but sometimes we have to)

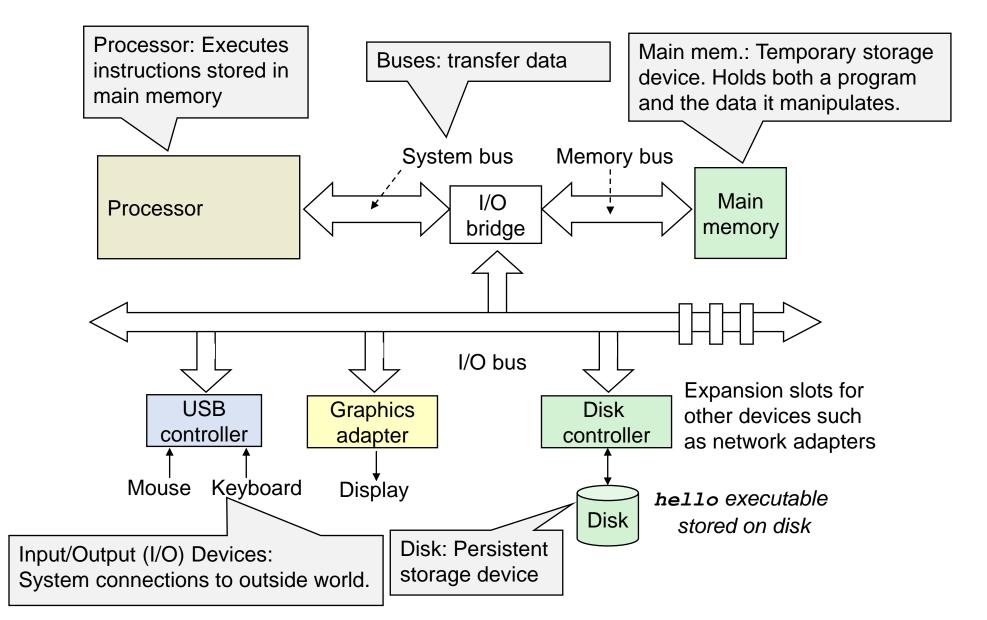
Running hello

• Running hello

| unix> ./hello | |
|---------------|--|
| hello, world | |
| unix> | |

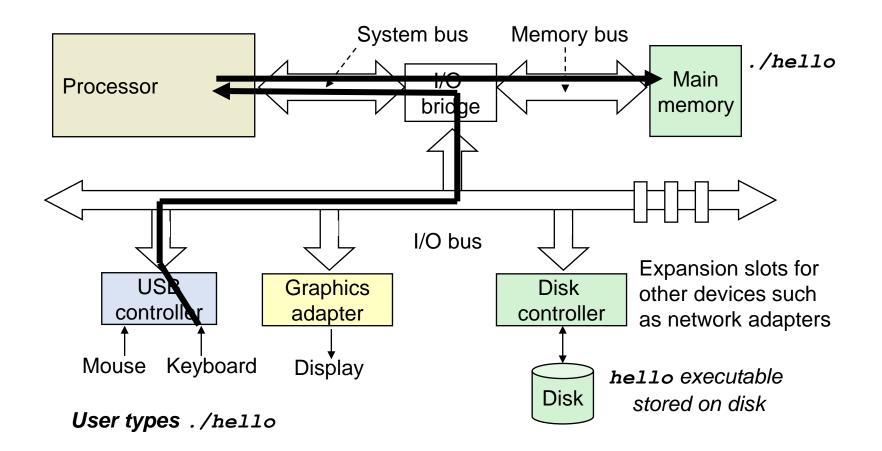
- What does the shell do?
 - Prints a prompt
 - Waits for you to type a command
 - Interpret the command
 - Then loads and runs the hello program
- What happens at the hardware level?

Hardware organization



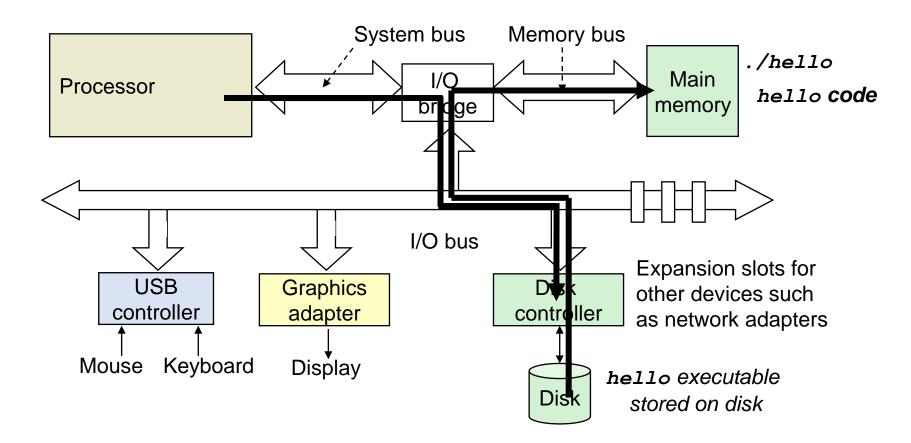
Running hello

Reading the ./hello command from the keyboard



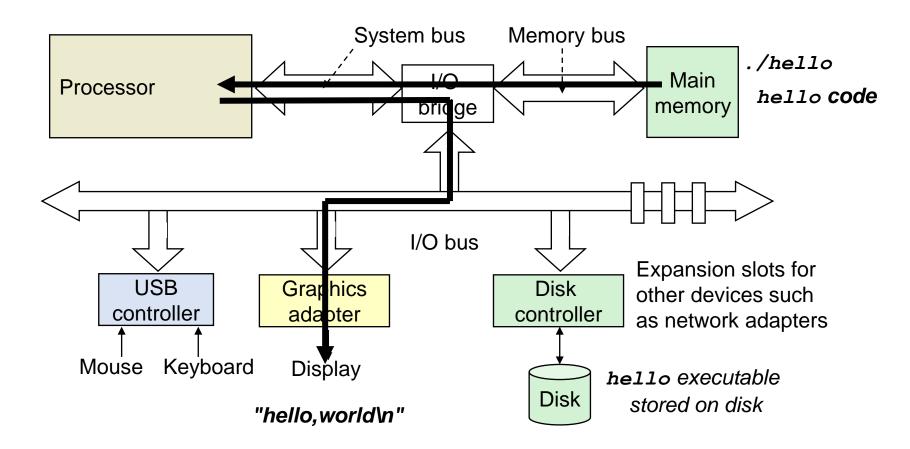
Running hello

Shell program loads the hello executable into main memory



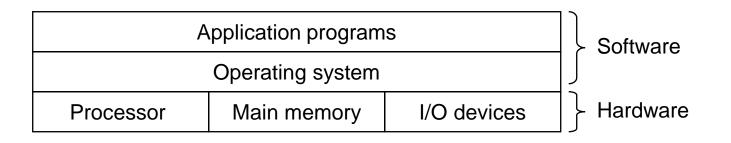
Running hello

The processor reads the hello code, executes instructions, and displays "hello..."



The Operating System (OS)

- Neither hello nor our shell interfaced with the hardware directly
 - All interactions were mediated by the *operating system*
- **Operating system:** a layer of software interposed between the application program and the hardware



- Primary goals
 - Protect resources from misuse by applications
 - Provide simple and uniform mechanisms for manipulating hardware devices
 - Manage sharing of resources between applications

A computer system is more than just HW

- A collection of intertwined hardware and software that must cooperate to achieve the end goal running applications
 - Hardware: expensive, fast, immutable
 - **Software**: cheap (comparatively), slow, flexible
 - Different tradeoffs
 - So we'll use them for different roles!

• The rest of the course will expand on this

Open Question + Break

 What part of the hello example takes the longest to run on a computer?

Open Question + Break

- What part of the hello example takes the longest to run on a computer?
 - The user typing (seconds)
 - Maybe that's cheating and we should start after they hit enter

Open Question + Break

- What part of the hello example takes the longest to run on a computer?
 - The user typing (seconds)
 - Maybe that's cheating and we should start after they hit enter
 - Almost certainly loading the program from disk (milliseconds)
 - Possibly sending text to graphics (microseconds milliseconds)
 - Definitely not executing the code (nanoseconds microseconds)

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Learning binary

• To understand how a computer really works we need to understand that data it operates on

- Computers hold data in memory as individual ones and zeros
 - These ones and zeros make up binary values
- So, we're going to need to understand binary
 - Binary will *definitely* come up again in this and other classes

Positional Numbering Systems

- The position of a *numeral* (e.g., digit) determines its contribution to the overall number
 - Makes arithmetic simple (compared to, say, roman numerals)
 - Any number has one canonical representation

- Example: base 10
 - $10456_{10} = 1*10^4 + 0*10^3 + 4*10^2 + 5*10^1 + 6*10^0$
 - Usually, we leave out the zeros:
 - $1*10^4$ + $4*10^2$ + $5*10^1$ + $6*10^0$

Other bases are also possible

- Base 60, used by the Babylonians
 - The source of 60 seconds in a minute, 60 minutes in an hour
 - And 360 degrees in a circle
- Base 20, used by the Maya and Gauls
 - Parts of this remain in French today
- Base 2, used by computers
 - Example: 10010010₂
 - Same idea as before: $1*2^7 + 1*2^4 + 1*2^1 = 128_{10} + 16_{10} + 2_{10} = 146_{10}$

Base 2 Example

- Computer Scientists use base 2 a LOT (especially in computer systems)
- Let's convert 138_{10} to base 2
- We need to decompose 138_{10} into a sum of powers of 2
 - Start with the largest power of 2 that is smaller or equal to 138_{10}
 - Subtract it, then repeat the process

$$138_{10} - 128_{10} = 10_{10}$$
$$10_{10} - 8_{10} = 2_{10}$$
$$2_{10} - 2_{10} = 0_{10}$$

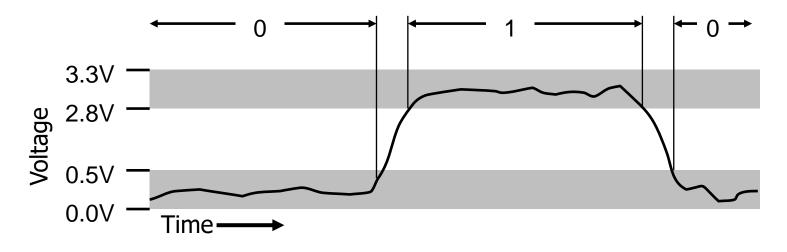
 $138_{10} = \mathbf{1} \times 128 + 0 \times 64 + 0 \times 32 + 0 \times 16 + \mathbf{1} \times 8 + 0 \times 4 + \mathbf{1} \times 2 + 0 \times 1$ $138_{10} = \mathbf{1} \times 2^7 + 0 \times 2^6 + 0 \times 2^5 + 0 \times 2^4 + \mathbf{1} \times 2^3 + 0 \times 2^2 + \mathbf{1} \times 2^1 + 0 \times 2^0$ $138_{10} = 10001010_2$

Binary practice

- Convert 101₂ to decimal
 - = $1 \times 2^2 + 0 \times 2^1 + 1 \times 2^0$
 - $\bullet = 4 + 0 + 1$
 - = 5₁₀
- Convert 4_{10} to binary: 100_2 (one less than 5)
- Convert 6_{10} to binary: 110_2 (one more than 5)

Why computers use Base 2

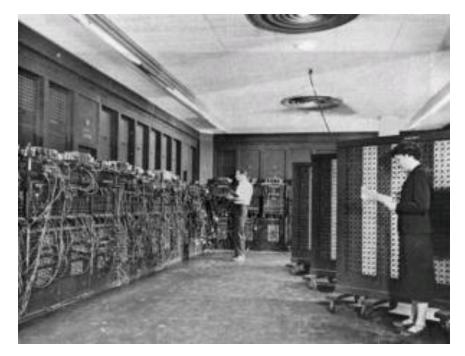
- Simple electronic implementation
 - Easy to store with bi-stable elements
 - Reliably transmitted on noisy and inaccurate wires



- Straightforward implementation of arithmetic functions
- (Pretty much) all computers use base 2

Why don't computers use Base 10?

- Because implementing it electronically is a pain
 - Hard to store
 - ENIAC (first general-purpose electronic computer) used 10 vacuum tubes / digit
 - Hard to transmit
 - Need high precision to encode 10 signal levels on single wire
 - Messy to implement digital logic functions
 - Addition, multiplication, etc.
 - (See CE203 for details)



Base 16: Hexadecimal

•

- Writing long sequences of 0s and 1s is tedious and error-prone
 - And takes up a lot of space on a page!
- So we'll often use base 16 (also called hexadecimal)

- Base 2 = 2 symbols (0, 1) Base 10 = 10 symbols (0-9) Base 16, need 16 symbols
 - Use letters A-F once we run out of decimal digits

| Hex | Decimal | Binary |
|-----|---------|--------|
| 0 | 0 | 0000 |
| 1 | 1 | 0001 |
| 2 | 2 | 0010 |
| 3 | 3 | 0011 |
| 4 | 4 | 0100 |
| 5 | 5 | 0101 |
| 6 | 6 | 0110 |
| 7 | 7 | 0111 |
| 8 | 8 | 1000 |
| 9 | 9 | 1001 |
| Α | 10 | 1010 |
| В | 11 | 1011 |
| С | 12 | 1100 |
| D | 13 | 1101 |
| E | 14 | 1110 |
| F | 15 | 1111 |

Base 16: Hexadecimal

- 16 = 2⁴, so every group of 4 bits becomes a hexadecimal digit (or *hexit*)
 - If we have a number of bits not divisible by 4, add 0s on the left (always ok, just like base 10)

"0x" prefix = it's in hex

| Hex | Decimal | Binary |
|-----|---------|--------|
| 0 | 0 | 0000 |
| 1 | 1 | 0001 |
| 2 | 2 | 0010 |
| 3 | 3 | 0011 |
| 4 | 4 | 0100 |
| 5 | 5 | 0101 |
| 6 | 6 | 0110 |
| 7 | 7 | 0111 |
| 8 | 8 | 1000 |
| 9 | 9 | 1001 |
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Bytes

- A single bit doesn't hold much information
 - Only two possible values: 0 and 1
 - So we'll typically work with larger groups of bits
- For convenience, we'll refer to groups of 8 bits as bytes
 - And usually work with multiples of 8 bits at a time
 - Conveniently, 8 bits = 2 hexits

• Some examples

"0b" prefix = it's in binary

- 1 byte: 0b01100111 = 0x67
- 2 bytes: $11000100 00101111_2 = 0xC42F$

Convert 0x42 to decimal

- Steps
 - Convert 0x42 to binary:

• Convert binary to decimal:

| Hex | Decimal | Binary | |
|-----|---------|--------|--|
| 0 | 0 | 0000 | |
| 1 | 1 | 0001 | |
| 2 | 2 | 0010 | |
| 3 | 3 | 0011 | |
| 4 | 4 | 0100 | |
| 5 | 5 | 0101 | |
| 6 | 6 | 0110 | |
| 7 | 7 | 0111 | |
| 8 | 8 | 1000 | |
| 9 | 9 | 1001 | |
| Α | 10 | 1010 | |
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Convert 0x42 to decimal

- Steps
 - Convert 0x42 to binary:
 - 0x4 -> 0b0100 0x2 -> 0b0010 0x42 -> 0b 0100 0010

• Convert binary to decimal:

Convert 0x42 to decimal

- Steps
 - Convert 0x42 to binary:
 - 0x4 -> 0b0100 0x2 -> 0b0010 0x42 -> 0b 0100 0010

- Convert binary to decimal:
 - $1*2^6 + 1*2^1 = 64 + 2 = 66$

Convert 0x42 to decimal

- Alternate method:
 - 0x42
 - = $4 \times 16^1 + 2 \times 16^0$
 - = 64 + 2
 - = 66
- But you're honestly better off converting hex to binary for now
 - It's good practice!

Big idea: bits can be used to represent anything

- Depending on the context, the bits 11000011 could mean
 - The number 195
 - The number -61
 - The number -1.1875
 - The value True
 - The character ` \-'
 - The **ret** x86 instruction

- You have to know the **context** to make sense of any bits you have!
 - People and software they write determine what the bits actually mean

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Outline

• Backup: Boolean Algebra

Boolean Algebra

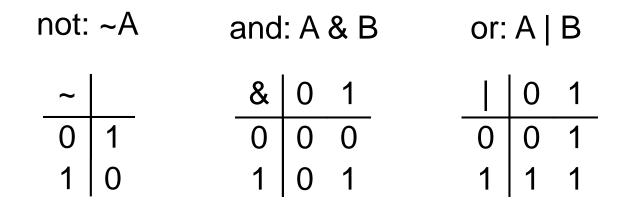
- You've programmed with and and or in earlier classes
 - Written && and || in C and C++
- Boolean algebra is a generalization of that
 - A mathematical system to represent (propositional) logic
 - 2 truth values: true = 1, false = 0
 - 3 operations: and = \mathbf{k} , or = \mathbf{I} , not (or complement) = \mathbf{v}
- Follow the rules for each operation to compute results
 - Rules are the like those you know from programming

$$(1 \mid 0) \& 0 \longrightarrow 1 \& 0 \longrightarrow 0$$

$$(1 \& 1) \& \sim (0 | 0) \rightarrow 1 \& \sim (0) \rightarrow 1 \& 1 \rightarrow 1$$

Truth Tables for Boolean Algebra

- For each possible value of each input, what is the output
 - Axes are the inputs
 - Inside of the table are the outputs



De Morgan's Laws, Exclusive Or

- Can express boolean operators in terms of the others
- De Morgan's laws: & using | and \sim , | using & and \sim
 - A & B = \sim (\sim A | \sim B)
 - A and B are true if and only if neither A nor B is false

• A | B =
$$\sim$$
(\sim A & \sim B)

- A or B are true if and only if A and B are not both false
- Can define new operators in terms of existing ones:
 - Exclusive or (xor, ^) in terms of inclusive or (|)
 - $A \land B = (\sim A \& B) | (A \& \sim B)$
 - Exactly one of A and B is true
 - $A \land B = (A | B) \& \sim (A \& B)$
 - Either A is true, or B is true, but not both
 - The two definitions are equivalent

xor: A ^ B

Generalized Boolean Algebra

• Binary bits can represent truth values: 0 = false, 1 = true

- Boolean operations can be extended to work on vectors of bits
 - Operations applied one bit at a time: *bitwise*

| 01101001 | 01101001 | 01101001 | |
|-----------------------|----------|-------------------|------------|
| <u>& 01010101</u> | 01010101 | <u>^ 01010101</u> | ~ 01010101 |
| 01000001 | 01111101 | 00111100 | 10101010 |

- All of the properties of Boolean algebra apply
 - Relationships between operations, etc.

Bit-level operations in C

- Operations &, |, ~, ^ available in C
 - Apply to any "integral" data type
 - long, int, short, char, unsigned
 - View arguments as bit vectors
 - Arguments applied bit-wise
- Examples (char data type, single byte)
 - $\sim 0 \times 00 \rightarrow 0 \times FF$

 $\sim 0000000_2 \rightarrow 1111111_2$

• $\sim 0x41 \rightarrow 0xBE$

 $\sim 01000001_2 \rightarrow 10111110_2$

• $0x69 \mid 0x55 \rightarrow 0x7D$

 $01101001_2 | 01010101_2 \rightarrow 01111101_2$

Logic operations in C – not the same!

- Logical operations ||, && and ! (Logic OR, AND & Not)
 - Contrast to bit-wise operators
 - View 0 as "False"
 - View anything nonzero as "True"
 - Always return 0 or 1 (i.e., false or true) rather than a sequence of bits
 - Early termination (if you can answer by just looking at the first argument, you are done)
- Examples (char data type)
 - $!0x41 \rightarrow 0x00$
 - $!0x00 \rightarrow 0x01$
 - $!!0x41 \rightarrow 0x01$
 - 0x59 && 0x35 → 0x01
 - 0x59 || 0x35 \rightarrow 0x01
 - p && *p (avoids null pointer access)

Watch out for && vs. & (and || vs. |) ... one of the more common slip-ups in C programming