

# Lecture 09

# Pointers and Arrays

CS213 – Intro to Computer Systems  
Branden Gena – Spring 2021

Slides adapted from:

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# Administrivia

- Homework 3 will be released tomorrow

# Today's Goals

- Understand C arrays
  - Single and multi-dimensional
- And how they translate into assembly code

# Outline

- **Pointers**
- One-dimensional Arrays
- Multi-dimensional Arrays
- Multi-level Arrays
- Dynamic arrays

# Basic Data Types

- Integers

- Stored & operated on in general (integer) registers
- Signed vs. unsigned depends on instructions used

<b>Intel</b>	<b>ASM</b>	<b>Bytes</b>	<b>C</b>
byte	<b>b</b>	1	<b>[unsigned] char</b>
word	<b>w</b>	2	<b>[unsigned] short</b>
double word	<b>d</b>	4	<b>[unsigned] int</b>
quad word	<b>q</b>	8	<b>[unsigned] long int</b>

# Floating point data

- Won't be focusing on floating point
  - Has changed much more than integer types across updates
  - Not all x86-64 machines have the same capabilities here
- Registers %xmm0 - %xmm15
  - 128-bit registers
  - On newest machines refer to as %ZMM0-%ZMM31 (512-bit registers)
- Instructions
  - addss (add scalar single-precision)
  - addsd (add scalar double-precision)
  - addpd (add packed double-precision, two doubles at once)

# More complex data types

- **Pointers and Arrays (today's lecture)**

```
int* a = &v;
```

```
int list[2] = {15, 27};
```

- **Structs and Unions (next lecture)**

```
typedef struct {  
    int a;  
    char b;  
    int* c;  
} mystruct_t;
```

# Example pointer code: calling `incr`

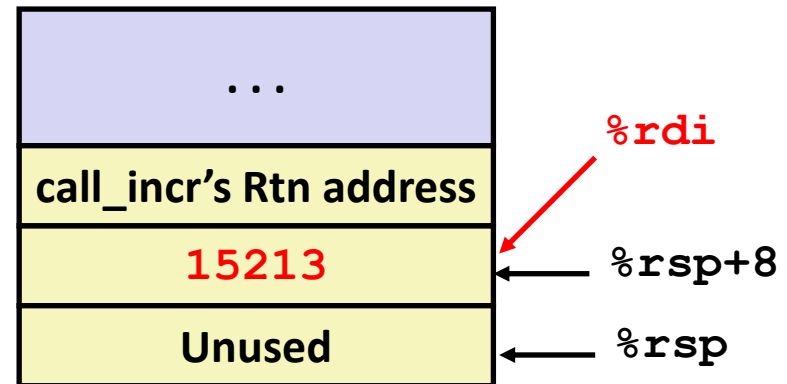
```
long call_incr() {  
    → long v1 = 15213; ↘  
    long v2 = incr(&v1, 3000);  
    return v1+v2;  
}
```

call\_incr:

```
subq    $16, %rsp  
movq    $15213, 8(%rsp)  
movq    $3000, %rsi  
leaq    8(%rsp), %rdi  
call    incr  
addq    8(%rsp), %rax  
addq    $16, %rsp  
ret
```

- Pointers are addresses
- `v1` must be stored on stack
  - Why? need to create pointer to it
- Compute pointer as `8(%rsp)`
  - Use `leaq` instruction

## Memory (stack)



Register	Use(s)
<code>%rdi</code>	<code>&amp;v1</code>
<code>%rsi</code>	3000



# Example pointer code : executing `incr`

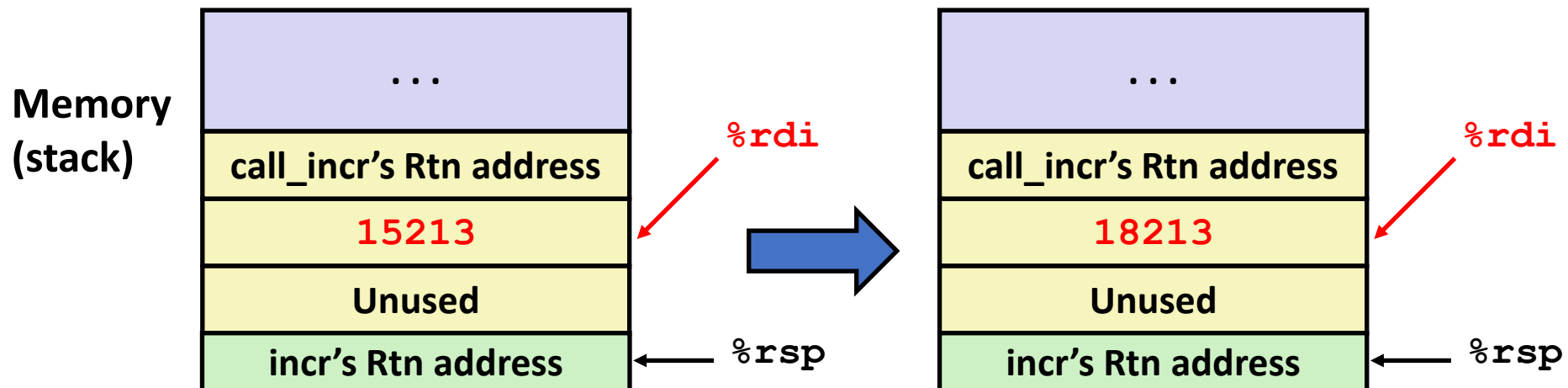
```
long incr(long *p, long val) {  
    long x = *p;  
    long y = x + val;  
    *p = y;  
    return x;  
}
```

```
incr:  
    movq    (%rdi), %rax  
    addq    %rax, %rsi  
    movq    %rsi, (%rdi)  
    ret
```

Register	Use(s)
<code>%rdi</code>	Argument <code>p</code>
<code>%rsi</code>	Argument <code>val</code> (3000)
<code>%rax</code>	...



Register	Use(s)
<code>%rdi</code>	Argument <code>p</code>
<code>%rsi</code>	18213
<code>%rax</code>	15213 (return value)



# Pointers to global variables

```
int global_var = 15;
```

```
int* myfunc(void) {  
    global_var += 2;  
    return &global_var;  
}
```

```
.text  
.globl myfunc  
.type myfunc, @function  
myfunc:  
    addl $2, 0x2f1f(%rip)  
    mov $0x404028, %eax  
    ret
```

```
.globl global_var  
.data  
.align4  
.type global_var, @object  
.size global_var, 4  
global_var:  
    .long 15
```

# Naming constants

These two are the same code.  
One just uses a name for the constant.

```
.text
.globl myfunc
.type myfunc, @function
myfunc:
addl $2, 0x2f1f(%rip)
mov $0x404028, %eax
ret
.globl global_var
.data
.align4
.type global_var, @object
.size global_var, 4
global_var:
.long 15
```

```
.text
.globl myfunc
.type myfunc, @function
myfunc:
addl $2, global_var(%rip)
mov $global_var, %eax
ret
.globl global_var
.data
.align4
.type global_var, @object
.size global_var, 4
global_var:
.long 15
```

# Outline

- Pointers
- **One-dimensional Arrays**
- Multi-dimensional Arrays
- Multi-level Arrays
- Dynamic arrays

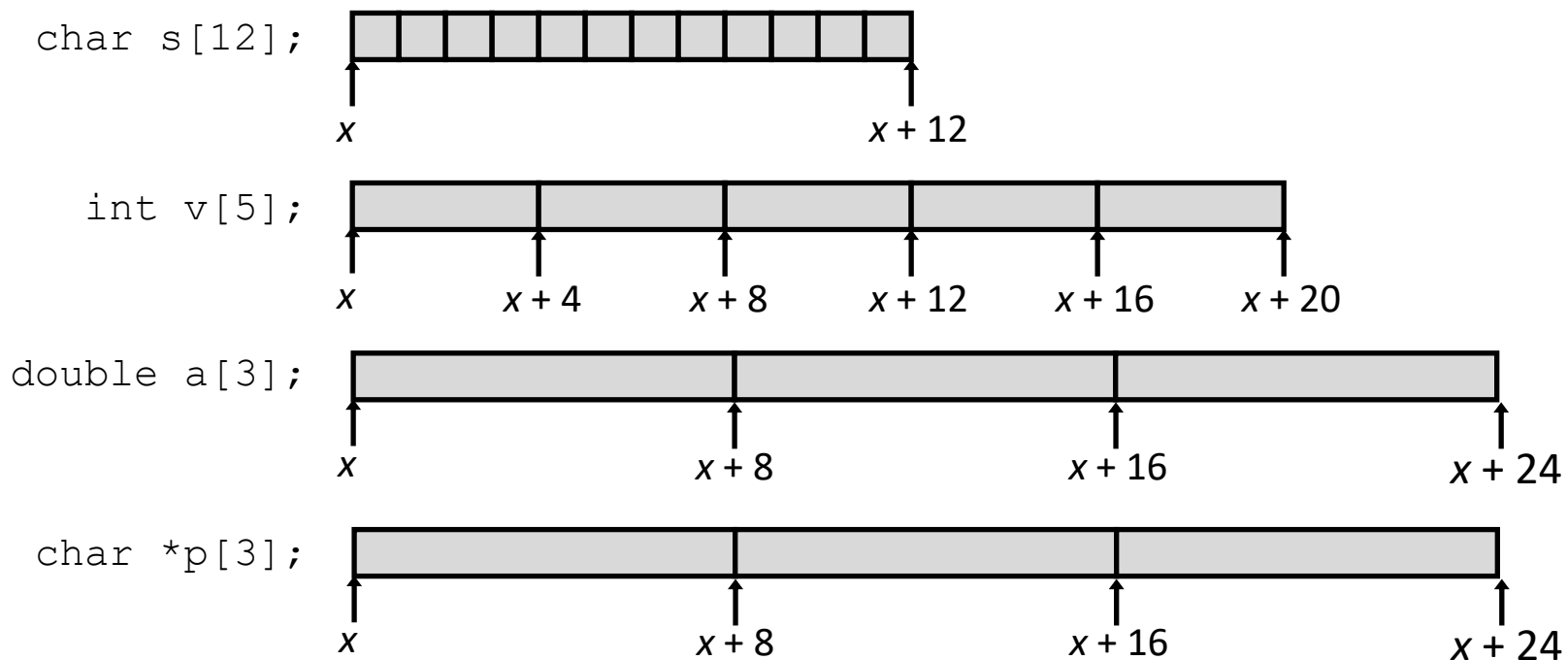
# One-Dimensional Array Allocation

- Basic Principle

`T A[L]; // e.g., int A[4];`

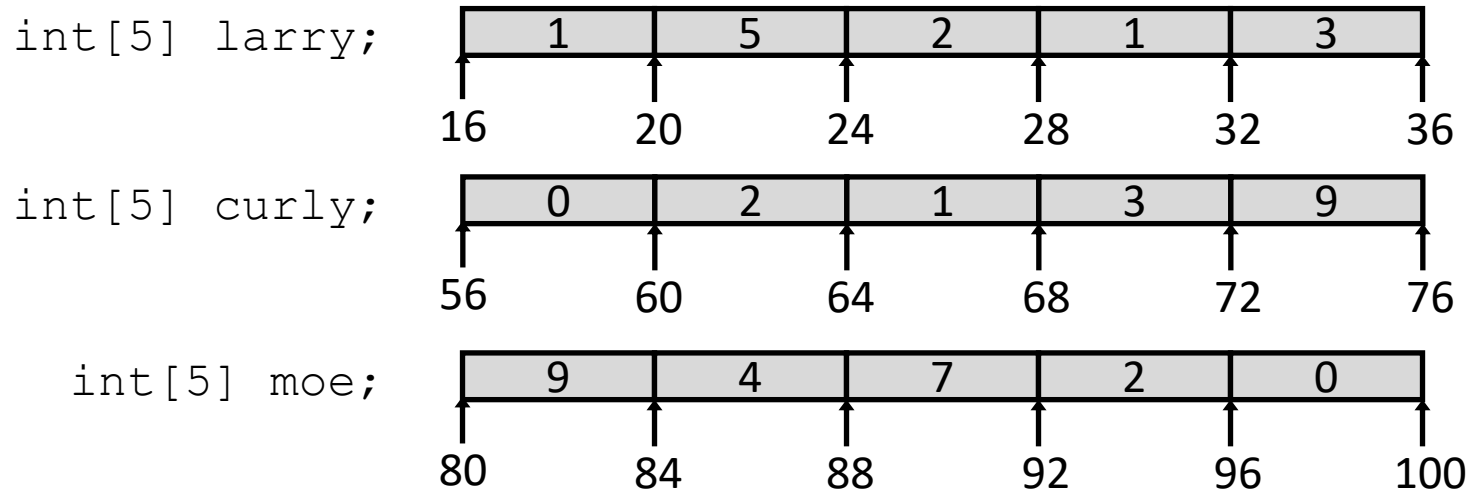
- Array of data type  $T$  and length  $L$

- Contiguously allocated region in memory of  $L * \text{sizeof}(T)$  bytes



# Placing arrays at addresses

```
int[5] larry = { 1, 5, 2, 1, 3 };  
int[5] curly = { 0, 2, 1, 3, 9 };  
int[5] moe    = { 9, 4, 7, 2, 0 };
```



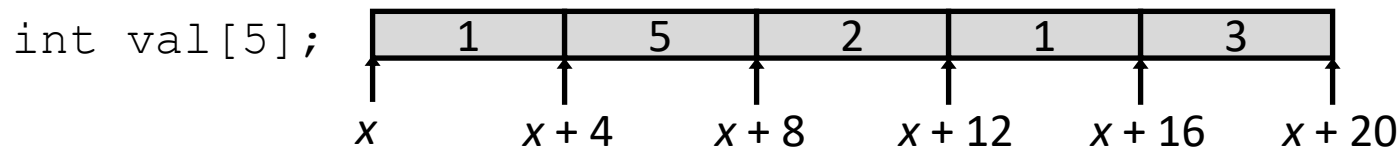
- Each array is allocated in contiguous 20 byte blocks
  - But no guarantee that `curly[]` will be right after `larry[]`, etc.

# Array Access and Pointer Arithmetic

- Basic Principle

$T$  **A**[ $L$ ];

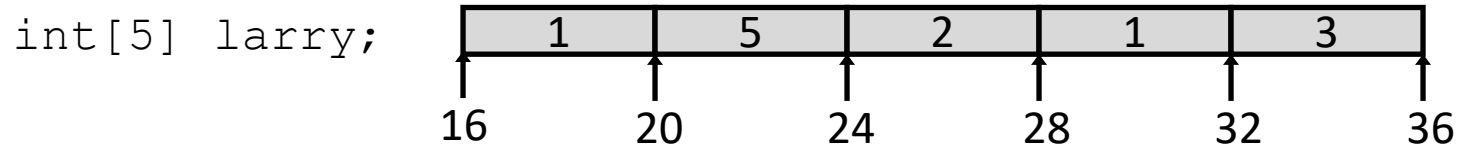
- Identifier **A** can be used as a pointer to array element 0: **A** is of type  $T^*$
- **Warning:** in C, arrays count # of elements, in assembly count # bytes!



- Reference

	Type	Value	
<code>val[4]</code>	<code>int</code>	3	
<code>val</code>	<code>int *</code>	$x$	
<code>val+1</code>	<code>int *</code>	$x+4$	
<code>&amp;val[2]</code>	<code>int *</code>	$x+8$	
<code>val[5]</code>	<b><code>int</code></b>	<b>??</b>	<b>No array bounds checking!!!</b>
<code>*(val+1)</code>	<code>int</code>	5	
<code>val + i</code>	<code>int *</code>	$x+4i$	

# One-Dimensional Array Accessing Example



```
int get_digit(int[5] larry, size_t digit)
{
    return larry[digit];
}
```

```
get_digit:
    # %rdi = larry
    # %rsi = digit
    movl (%rdi,%rsi,4),%rax    # z[digit]
    retq
```

`%rdi` -> starting address of array

`%rsi` -> array index

- Desired digit at `%rdi + 4*%rsi`
- Use memory addressing!  
(`%rdi,%rsi,4`)
- Now we see why x86 memory operands are `D(Rb, Ri, s)`
  - Scale 1, 2, 4, or 8



# One-Dimensional Array Loop Example

```
void zincr(int *z) {
    size_t i;
    for (i = 0; i < 4; i++)
        z[i]++;
}
```

```
zincr:
    # %rdi = z
    movl    $0, %eax          # i = 0
    jmp     .L3              # goto middle
.L4:                          # loop:
    → addl   $1, (%rdi,%rax,4) # z[i]++
    addq   $1, %rax          # i++
.L3:                          # middle:
    cmpq   $4, %rax          # i:4
    jbe    .L4              # if i<=4, goto loop
    retq
```

# Quiz + Break

`z -> %rdi`

`i -> %rax`

`addl $1, (%rdi,%rax,4) # z[i]++ (int z[])`

- What changes if `z` is instead an array of:
  - short
  - char
  - bool
  - char\*
  - unsigned int

# Quiz + Break

`z -> %rdi`

`i -> %rax`

`addl $1, (%rdi,%rax,4) # z[i]++ (int z[])`

- What changes if `z` is instead an array of:
  - short `addl $1, (%rdi,%rax,2)`
  - char
  - bool
  - char\*
  - unsigned int

# Quiz + Break

`z -> %rdi`

`i -> %rax`

`addl $1, (%rdi,%rax,4) # z[i]++ (int z[])`

- What changes if `z` is instead an array of:
  - short `addl $1, (%rdi,%rax,2)`
  - char `addl $1, (%rdi,%rax,1)`
  - bool
  - char\*
  - unsigned int

# Quiz + Break

`z -> %rdi`

`i -> %rax`

`addl $1, (%rdi,%rax,4) # z[i]++ (int z[])`

- What changes if `z` is instead an array of:
  - short `addl $1, (%rdi,%rax,2)`
  - char `addl $1, (%rdi,%rax,1)`
  - bool `addl $1, (%rdi,%rax,1)`
  - char\*
  - unsigned int

# Quiz + Break

`z -> %rdi`

`i -> %rax`

`addl $1, (%rdi,%rax,4) # z[i]++ (int z[])`

- What changes if `z` is instead an array of:
  - `short` `addl $1, (%rdi,%rax,2)`
  - `char` `addl $1, (%rdi,%rax,1)`
  - `bool` `addl $1, (%rdi,%rax,1)`
  - `char*` `addl $1, (%rdi,%rax,8)`
  - `unsigned int`

# Quiz + Break

`z -> %rdi`

`i -> %rax`

`addl $1, (%rdi,%rax,4) # z[i]++ (int z[])`

- What changes if `z` is instead an array of:
  - `short` `addl $1, (%rdi,%rax,2)`
  - `char` `addl $1, (%rdi,%rax,1)`
  - `bool` `addl $1, (%rdi,%rax,1)`
  - `char*` `addl $1, (%rdi,%rax,8)`
  - `unsigned int` Nothing. Still 4 bytes. `add` works the same on sign/unsigned

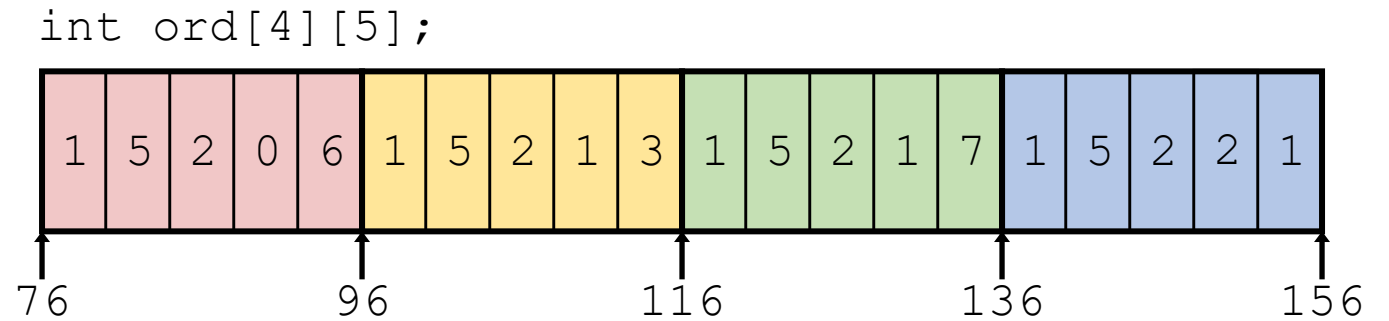
# Outline

- Pointers
- One-dimensional Arrays
- **Multi-dimensional Arrays**
- Multi-level Arrays
- Dynamic arrays



# Multidimensional (Nested) Array Example

```
int ord[4][5] =  
    /* 4 rows, 5 cols */  
    {{1, 5, 2, 0, 6},  
     {1, 5, 2, 1, 3},  
     {1, 5, 2, 1, 7},  
     {1, 5, 2, 2, 1}};
```



- Let's decipher "int ord[4][5]"

`int ord[4][5]`: `ord` is an array of **4** elements, allocated contiguously

`int ord[4][5]`: Each element is an array of **5** `int`'s, allocated contiguously

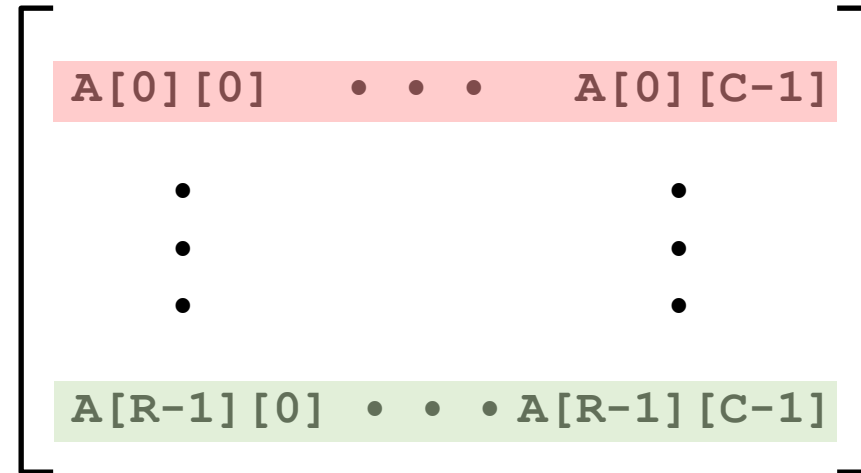
- "Row-Major" ordering of all elements guaranteed
  - Entire row (all columns in it) will be placed in memory before the next row starts

# Multidimensional (Nested) Arrays

- Declaration

$T \mathbf{A}[R][C];$

- 2D array of data type  $T$
- $R$  rows,  $C$  columns
- Type  $T$  element requires  $K$  bytes

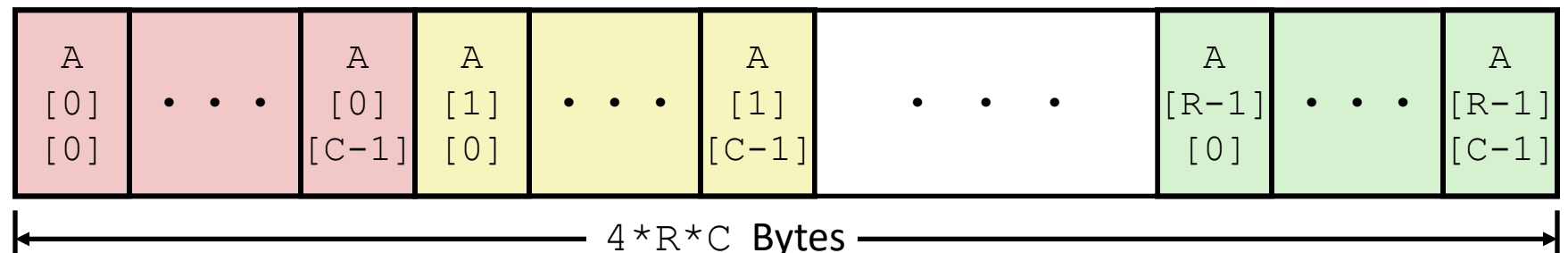


- Types

- What is  $A$ ?  $T[R][C]$
- What is  $A[i]$ ?  $T[C]$
- What is  $A[i][j]$ ?  $T$

- Arrangement

- Row-Major Ordering

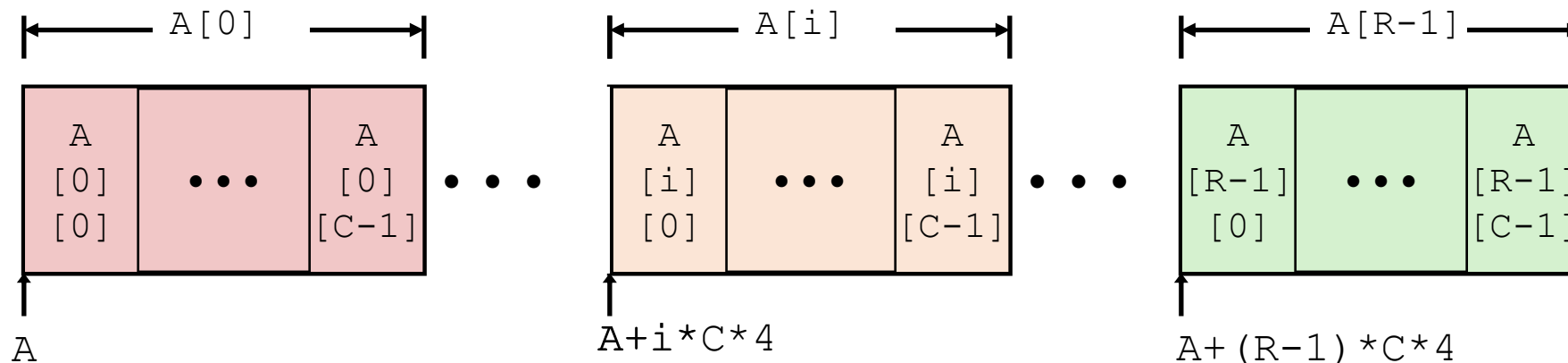


`int A[R][C];`

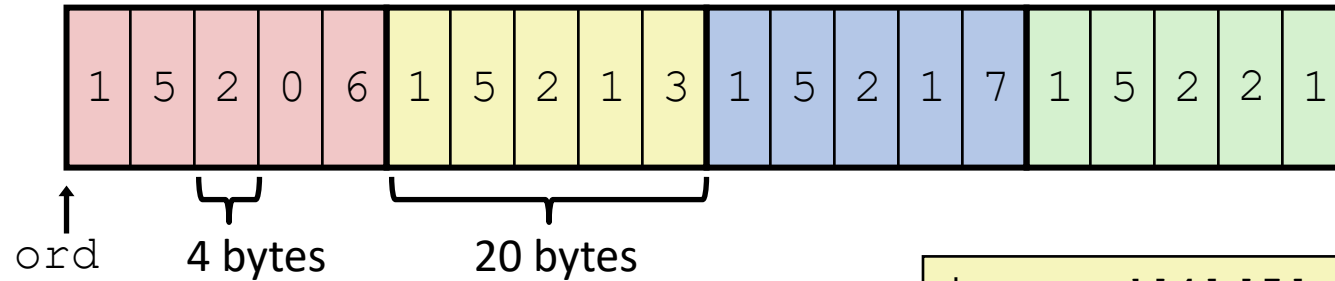
# Nested Array Row Access

- To figure out how to get the element we want
  - Let's first figure out how to get the row we want (its starting address)
- Row Vectors
  - $\mathbf{A}[i]$  (row) is array of  $C$  elements
  - Each element of type  $T$  requires  $K$  bytes
  - Starting address  $\mathbf{A} + i * (C * K)$

```
int A[R][C];
```



# Nested Array Row Access Code

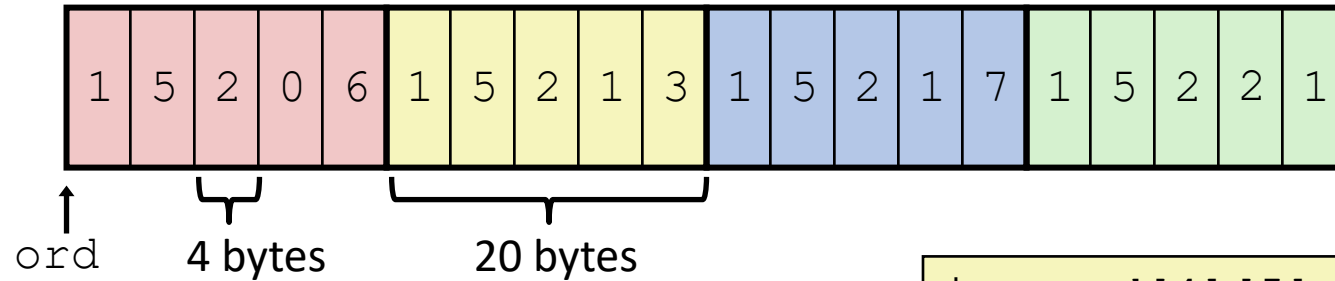


```
int *get_ord_row(size_t index)
{
    return ord[index];
}
```

```
int ord[4][5] =
    {{1, 5, 2, 0, 6},
     {1, 5, 2, 1, 3 },
     {1, 5, 2, 1, 7 },
     {1, 5, 2, 2, 1 }};
```

```
# %rdi = index
leaq (%rdi,%rdi,4),%rax    # %rax = 5 * index
leaq ord(,%rax,4),%rax    # %rax = ord + 4*(5*index)
```

# Nested Array Row Access Code



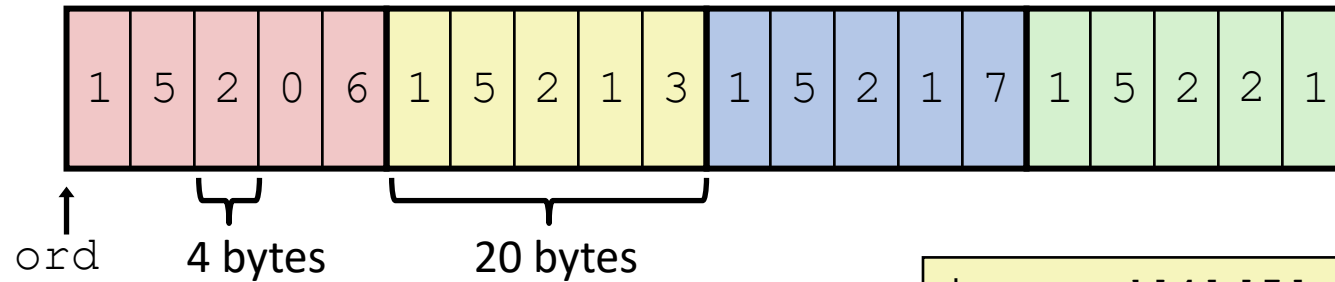
```
int *get_ord_row(size_t index)
{
    return ord[index];
}
```

```
int ord[4][5] =
    {{1, 5, 2, 0, 6},
     {1, 5, 2, 1, 3 },
     {1, 5, 2, 1, 7 },
     {1, 5, 2, 2, 1 }};
```

```
# %rdi = index
leaq (%rdi,%rdi,4),%rax    # %rax = 5 * index
leaq ord(,%rax,4),%rax    # %rax = ord + 4*(5*index)
```

- What's that displacement?
  - Constant address
  - `ord` is a global. Always in a location known at compile-time. So constant address!

# Nested Array Row Access Code



```
int *get_ord_row(size_t index)
{
    return ord[index];
}
```

```
int ord[4][5] =
    {{1, 5, 2, 0, 6},
     {1, 5, 2, 1, 3 },
     {1, 5, 2, 1, 7 },
     {1, 5, 2, 2, 1 }};
```

```
# %rdi = index
leaq (%rdi,%rdi,4),%rax    # %rax = 5 * index
leaq ord(,%rax,4),%rax    # %rax = ord + 4*(5*index)
```

- Row Vector

- **ord[index]** is array of 5 **int**'s
- Starting address **ord + 20\*index**

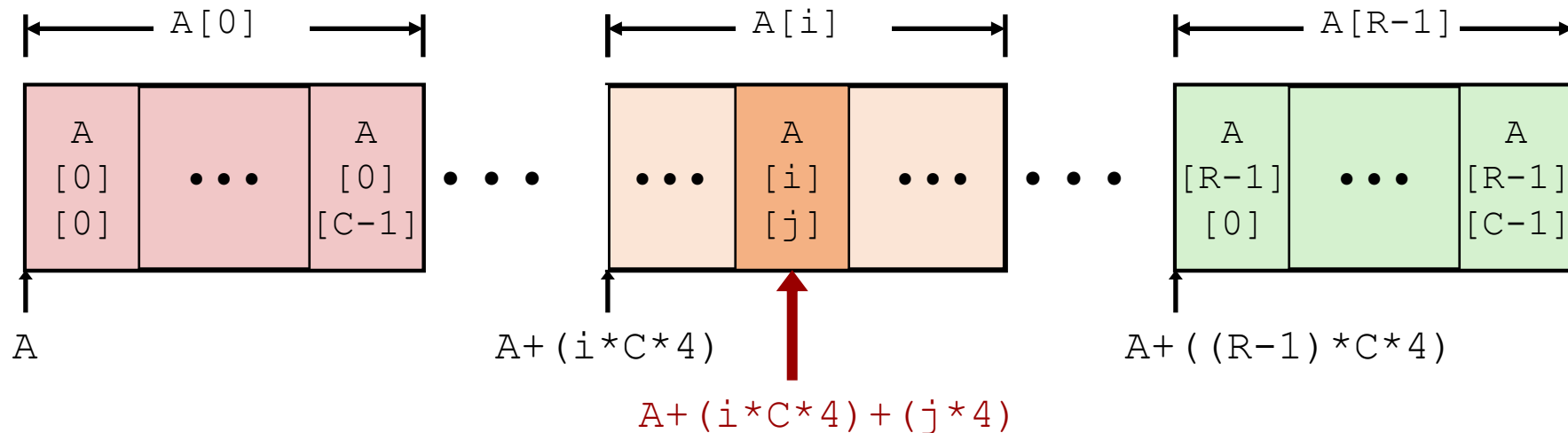
- Assembly Code

- Computes and returns address
- **ord + 4\*(index+(4\*index))**

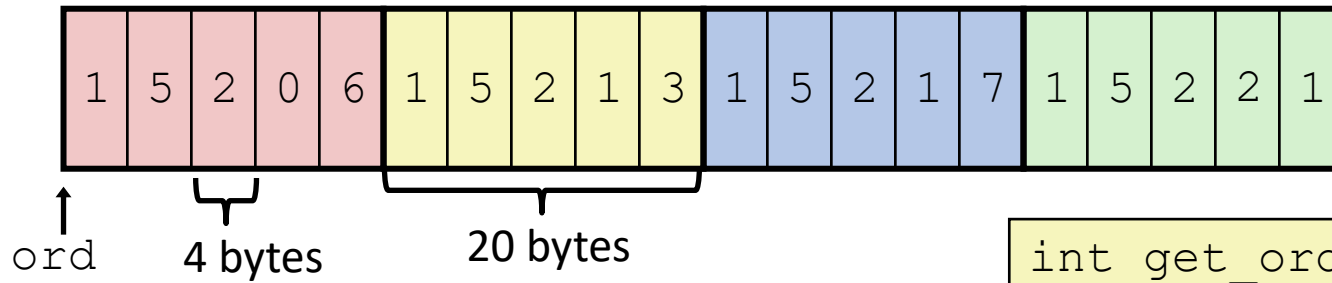
# Nested Array Element Access

- Now, let's find the *element* that we want
- Array Elements
  - $A[i][j]$  is element of type  $T$ , which requires  $K$  bytes
  - Address  $A + i * (C * K) + j * K = A + (i * C + j) * K$

```
int A[R][C];
```



# Nested Array Element Access Code



```
int get_ord_digit(size_t index, size_t digit)
{
    return ord[index][digit];
}
```

```
# %rdi = index
leaq  (%rdi,%rdi,4), %rax    # 5*index
addq  %rax, %rsi            # 5*index + digit
movl  ord(,%rsi,4), %eax    # M[ord + 4*(5*index+digit)]
```

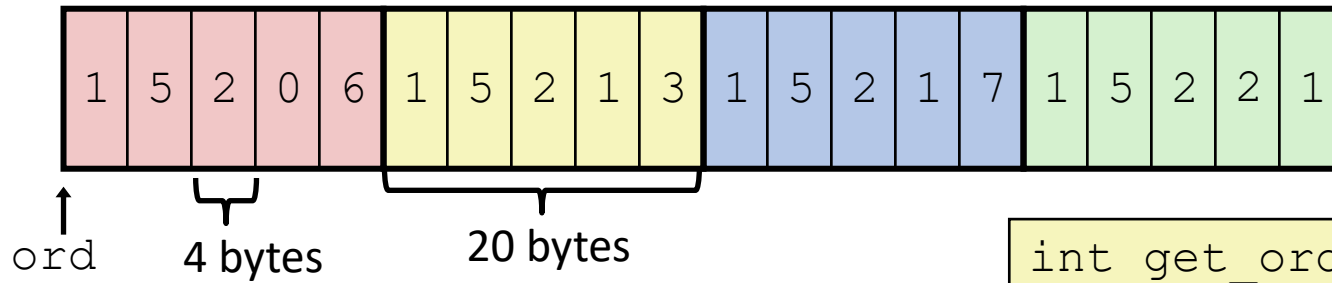
- Array Elements

- `ord[index][digit]` is type `int`

- Address:  $\text{ord} + 20 \cdot \text{index} + 4 \cdot \text{digit} = \text{ord} + 4 \cdot (5 \cdot \text{index} + \text{digit})$



# Nested Array Element Access Code



```
int get_ord_digit(size_t index, size_t digit)
{
    return ord[index][digit];
}
```

```
# %rdi = index
leaq  (%rdi,%rdi,4), %rax    # 5*index
addq  %rax, %rsi            # 5*index + digit
movl  ord(,%rsi,4), %eax    # M[ord + 4*(5*index+digit)]
```

- Array Elements

- `ord[index][digit]` is type `int`

- Address:  $\text{ord} + 20 \cdot \text{index} + 4 \cdot \text{digit} = \text{ord} + 4 \cdot (5 \cdot \text{index} + \text{digit})$

- QUIZ: what is the address of `ord[2][4]`? `ord+56`

# Break + Practice

- Find the addresses (assume array starts at address 0)

- $A + (i * C * K) + (j * K)$

- `int A[16][16];`      `A[1][3]`

- `char B[16][16];`      `B[10][7]`

- `char* B[10][10];`      `B[0][2]`

# Break + Practice

- Find the addresses (assume array starts at address 0)

- $A + (i * C * K) + (j * K)$

- `int A[16][16];`

`A[1][3]`

- $A + (i * C * K) + (j * K) =$

$0 + (1 * 16 * 4) + (3 * 4) = 64 + 12 = 76$

- `char B[16][16];`

`B[10][7]`

- `char* B[10][10];`

`B[0][2]`

# Break + Practice

- Find the addresses (assume array starts at address 0)

- $A + (i * C * K) + (j * K)$

- `int A[16][16];`

`A[1][3]`

- $A + (i * C * K) + (j * K) = 0 + (1 * 16 * 4) + (3 * 4) = 64 + 12 = 76$

- `char B[16][16];`

`B[10][7]`

- $A + (i * C * K) + (j * K) = 0 + (10 * 16 * 1) + (7 * 1) = 160 + 7 = 167$

- `char* B[10][10];`

`B[0][2]`

# Break + Practice

- Find the addresses (assume array starts at address 0)

- $A + (i * C * K) + (j * K)$

- `int A[16][16];`

`A[1][3]`

- $A + (i * C * K) + (j * K) = 0 + (1 * 16 * 4) + (3 * 4) = 64 + 12 = 76$

- `char B[16][16];`

`B[10][7]`

- $A + (i * C * K) + (j * K) = 0 + (10 * 16 * 1) + (7 * 1) = 160 + 7 = 167$

- `char* B[10][10];`

`B[0][2]`

- $A + (i * C * K) + (j * K) = 0 + (0 * 10 * 8) + (2 * 8) = 16$

# Outline

- Pointers
- One-dimensional Arrays
- Multi-dimensional Arrays
- **Multi-level Arrays**
- Dynamic arrays

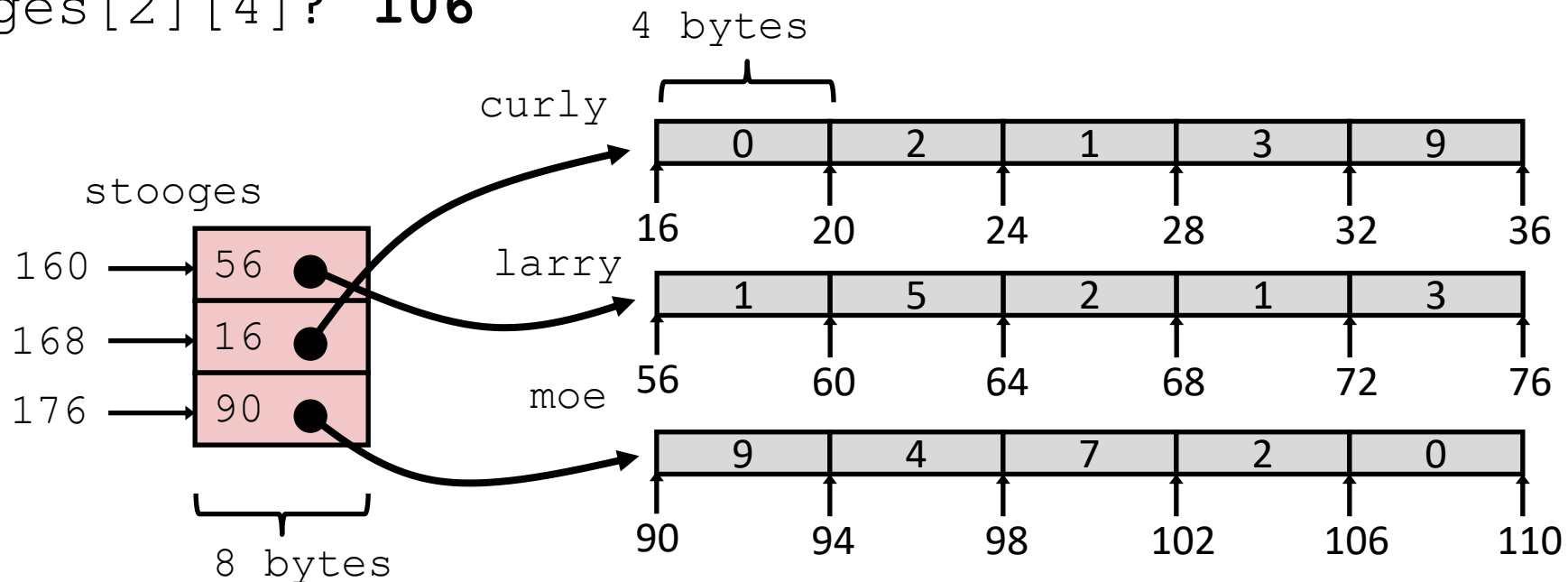
# Multi-Level Array Example

```
int larry [5] = { 1, 5, 2, 1, 3 };  
int curly [5] = { 0, 2, 1, 3, 9 };  
int moe [5]   = { 9, 4, 7, 2, 0 };
```

```
int *stooges[3] =  
    {larry, curly, moe};
```

- Variable `stooges` denotes array of 3 elements
- Each element is a pointer (8 bytes)
- Each pointer points to array of `ints`
- `stooges` is of type `int* []`
- `stooges` is of type `int**`

QUIZ: What is the address of `stooges[2][4]`? **106**



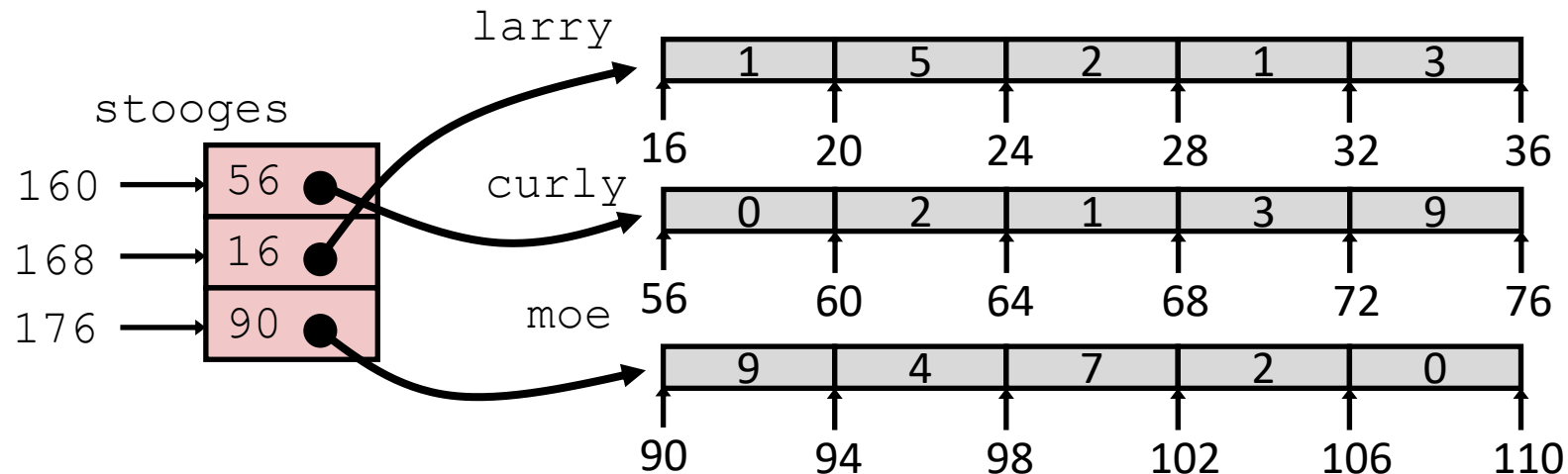
# Multi-Level Array Element Access

```
int get_stooge_digit
(size_t index, size_t digit){
    return stooges[index][digit];
}
```

- Must do two memory reads
  - First get pointer to row array
  - Then access element within array

```
salq    $2, %rsi          # 4*digit
addq    stooges(,%rdi,8), %rsi # p = stooges[8*index] + 4*digit
movl    (%rsi), %eax      # return *p
ret
```

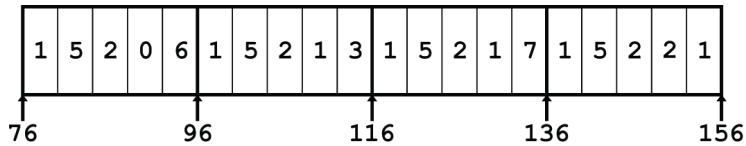
Element access  $\text{Mem}[\text{Mem}[\text{stooges} + 8 * \text{index}] + 4 * \text{digit}]$



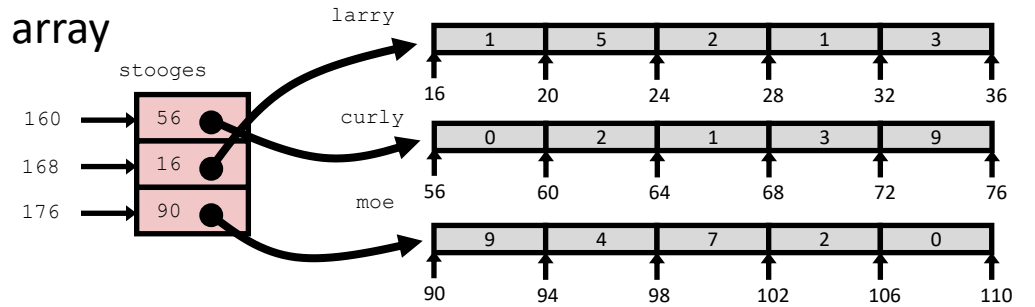


# Nested vs. Multi-Level Array Element Accesses

Nested array



Multi-level array



```
int ord [4][5];
```

```
int get_ord_digit  
  (size_t index, size_t digit){  
  return ord[index][digit];  
}
```

```
int larry[5], curly[5], moe[5];  
int *stooges[3] = {larry, curly, moe};
```

```
int get_stooge_digit  
  (size_t index, size_t digit){  
  return stooges[index][digit];  
}
```

Accesses look similar in C, but address computations are very different:

ord is of type `int*`

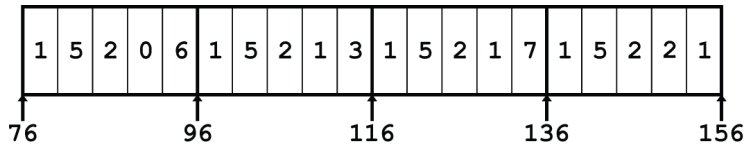
stooges is of type `int**`

`Mem[ord + (20*index) + (4*digit)]`

`Mem[Mem[stooges + (8*index)] + (4*digit)]`

# Nested versus Multi-Level Arrays

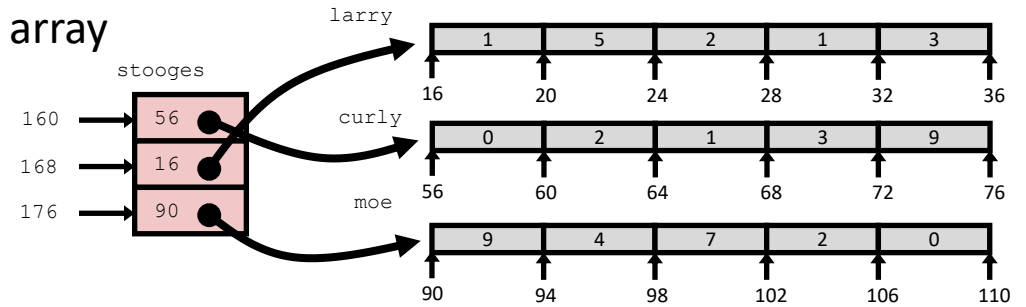
Nested array



$$\text{Mem}[\text{ord} + (20 * \text{index}) + (4 * \text{digit})]$$

- Strengths
  - Fast element access
    - Single memory access
  - Efficient memory usage
    - Stored in contiguous memory
- Limitations
  - Requires fixed size rows
  - Large memory usage
    - All rows need to be allocated

Multi-level array



$$\text{Mem}[\text{Mem}[\text{stooges} + (8 * \text{index})] + (4 * \text{digit})]$$

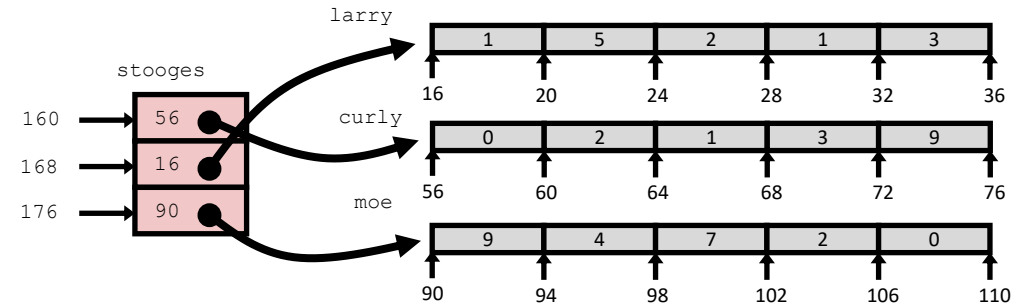
- Strengths
  - Rows may be of different size
  - Rows could even be different types
    - First array would store `void*`
- Limitations
  - Slow element access
    - Two memory references
  - Memory fragmentation
    - Many small chunks allocated

# Outline

- Pointers
- One-dimensional Arrays
- Multi-dimensional Arrays
- Multi-level Arrays
- **Dynamic arrays**

# Dynamic Multi-dimensional arrays – multi-level

- Multi-level is one way to make them



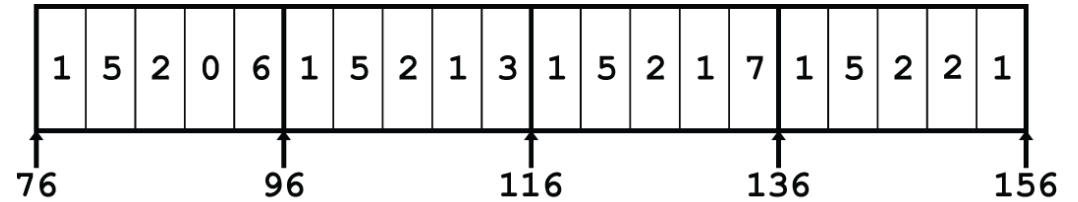
```
int** array_2d = (int**)malloc(rows * sizeof(int*));
```

```
for (int i=0; i<rows; i++) {  
    array_2d[i] = (int*)malloc(cols * sizeof(int));  
}
```

```
array_2d[2][4] = 0;
```

# Dynamic multi-dimensional arrays - nested

- Nested works as well
  - Handle nested manually
  - Make sure you get it right!



```
int* array_2d = (int*)malloc(rows * cols * sizeof(int));  
array_2d[2*cols + 4] = 0; // array_2d[2][4]
```

# Nested arrays – static versus dynamic

```
void testarray(void) {  
    volatile int A[16][16];  
    A[2][4] = 0;  
  
    volatile int* B =  
        (int*)malloc(16*16*sizeof(int));  
    B[2*16 + 4] = 0;  
}
```

```
testarray():  
    sub    $0x408,%rsp  
    movl   $0x0,0x90(%rsp)  
    mov    $0x400,%edi  
    call  400480 <malloc@plt>  
    movl   $0x0,0x90(%rax)  
    add    $0x408,%rsp  
    ret
```

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    A[2][4] = 0;  
  
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- Dynamic arrays